

On the Decidability of HTN Planning with Task Insertion

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Abstract

We give a simplified propositional HTN formalization inspired by the formalization of Erol et al. [2] and show

- the plan existence problem is still undecidable despite the simplifications
- HTN planning with insertion (hybrid planning) is decidable; from the proof of decidability, we obtain an upper complexity bound of **EXPSpace** for the plan existence problem for propositional hybrid planning

Possible Modifications

- Decomposition:
 - given a task network $tn = (T, \prec)$, use method $(t, tn') \in M$ to replace $t \in T$ by tn' . Then adjust ordering constraints.
- Insertion:
 - insert primitive tasks from O
 - insert ordering constraints

Decidability Criteria

Imposing certain restrictions on planning problems can make HTN planning decidable:

- Criterion 1:** Decomposition tree is acyclic
Intuition: search space becomes finite.
- Criterion 2:** All methods are totally ordered (cf. SHOP system [5])
Intuition: solution corresponds to an intersection of a regular grammar with one that is context-free (decidable problem).
- Criterion 3:** Methods contain at most one compound task (regular)
Intuition: the combinations of possible states before and after the abstract task are finite.
- Criterion 4:** Allow task insertion
Intuition: insertion makes cyclic method applications superfluous \rightarrow minimal solution lengths are bounded like in classical planning.

Definition (Task Network)

A task network $tn = (T, \prec)$ is a partially ordered sequence of tasks:

- T is a finite and non-empty set of tasks
- $\prec \subseteq T \times T$ is a strict partial order on T

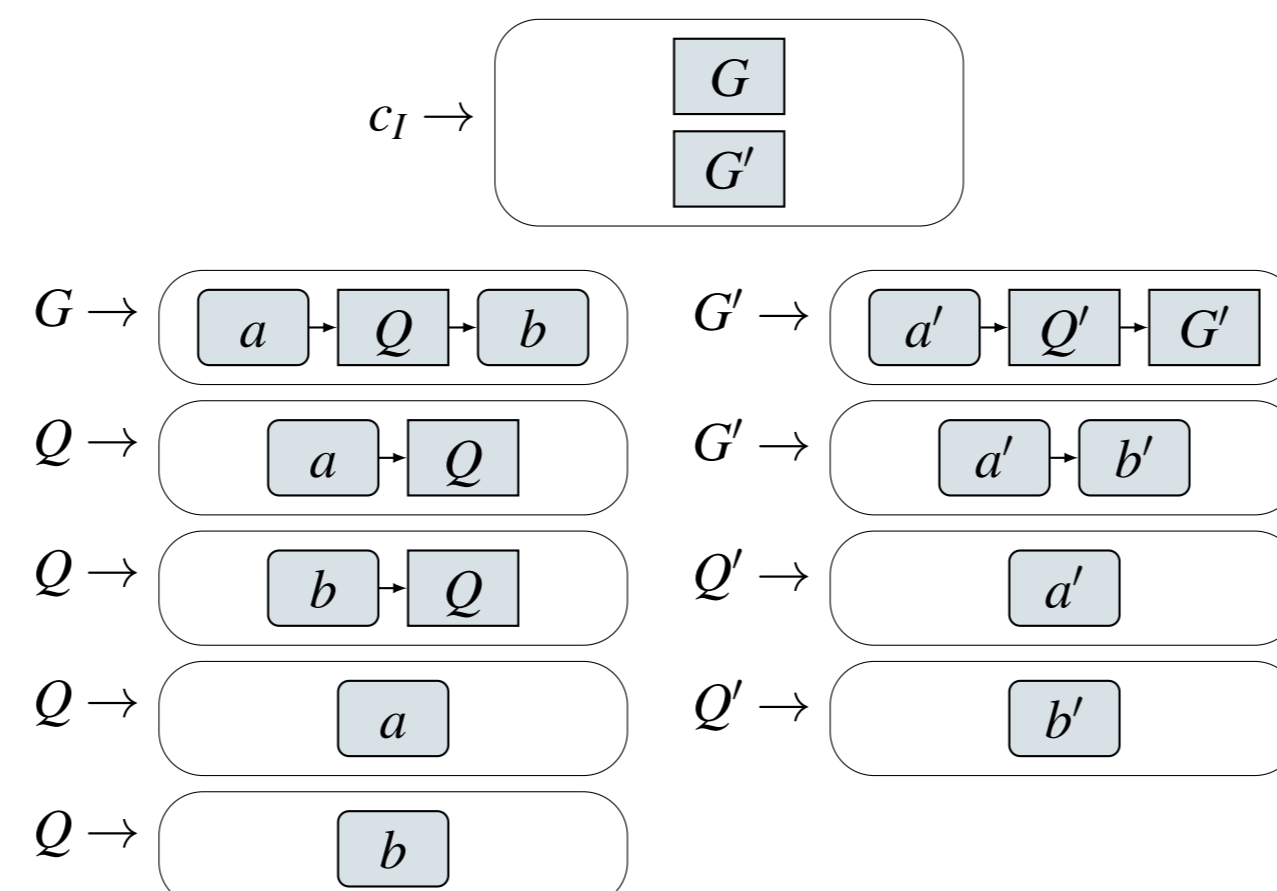
Definition (Planning Problem)

A planning problem is a 6-tuple $P = (V, C, O, M, c_I, s_I)$ and

- V is a finite set of state variables
- C is a finite set of compound tasks
- O is a finite set of primitive tasks, for $o \in O$, $(prec(o), add(o), del(o)) \in 2^V \times 2^V \times 2^V$ is an operator
- $M \subseteq C \times TN$ is a finite set of decomposition methods
- $c_I \in C$ is the initial task
- $s_I \in 2^V$ is the initial state

Note that the part of the planning problem that is usually called the domain (tasks and methods) is given with the problem.

Figure 2: decomposition methods

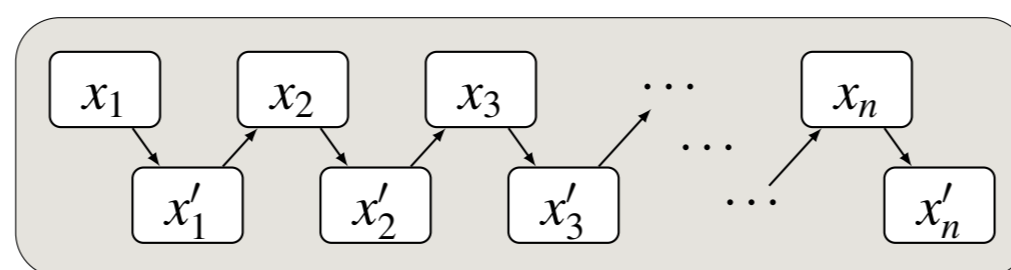


Here, the decomposition methods describe two context-free grammars (CFGs); their languages are $L(G) = a(a|b)^+b$ and $L(G') = (a'(a'|b'))^+a'b'$.

Solution Criterion

- A task network tn is an HTN solution iff:
 - tn is obtained via decomposition
 - there is an executable linearization of tn 's tasks
- A task network tn is a hybrid solution iff:
 - tn is obtained by decomposition followed by insertion
 - there is an executable linearization of tn 's tasks

Figure 3: structure of solutions



Theorem

The plan existence problem is undecidable for HTN planning.

Proof idea (by Erol et al. [4]):

- the following question is undecidable: Given two CFGs, do their languages produce a common word?
- observe that the production rules of CFGs can be simulated by decomposition methods
- given two CFGs, construct a planning problem which has an HTN solution iff the languages of the two grammars have a non-empty intersection

Figure 4: removing cycles from a decomposition tree

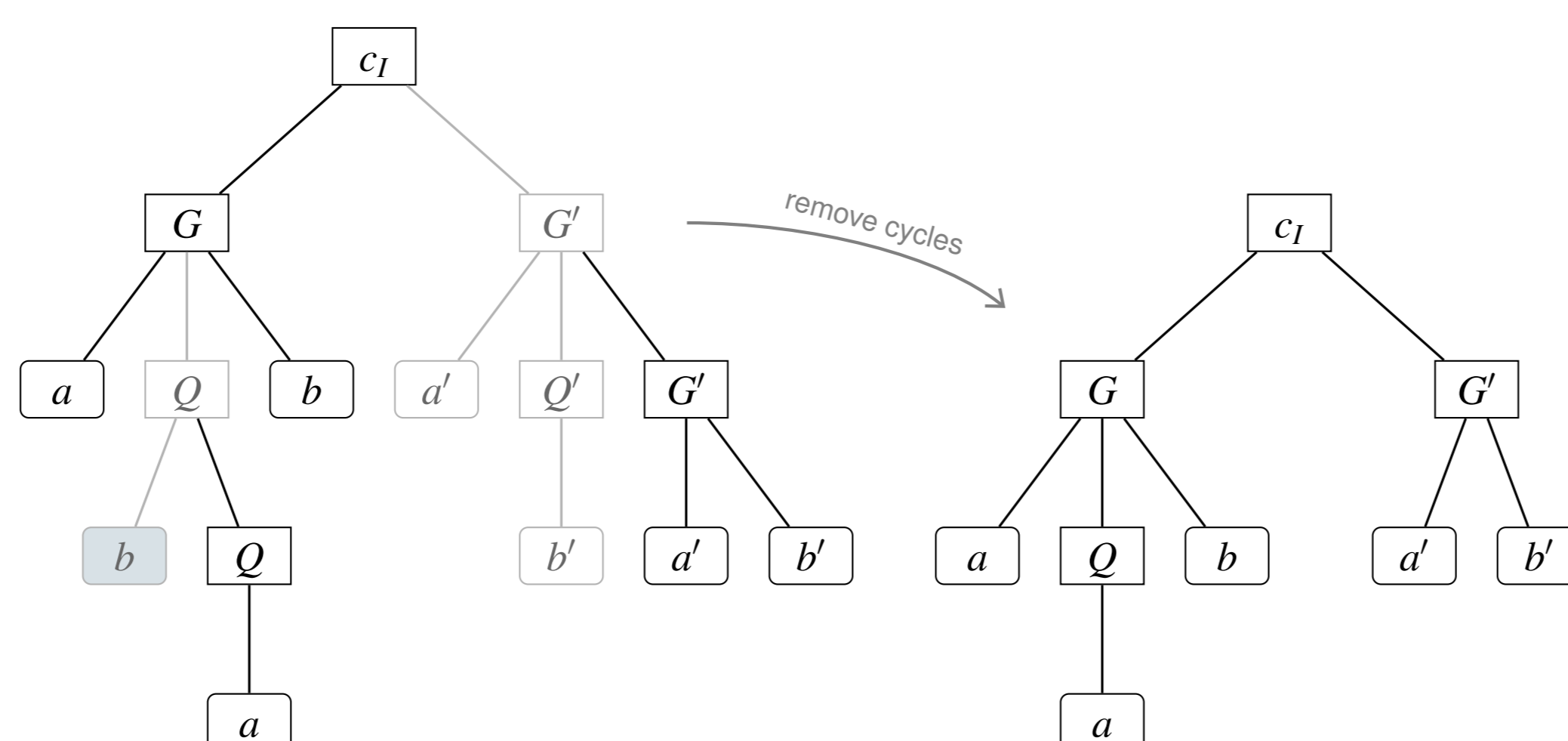
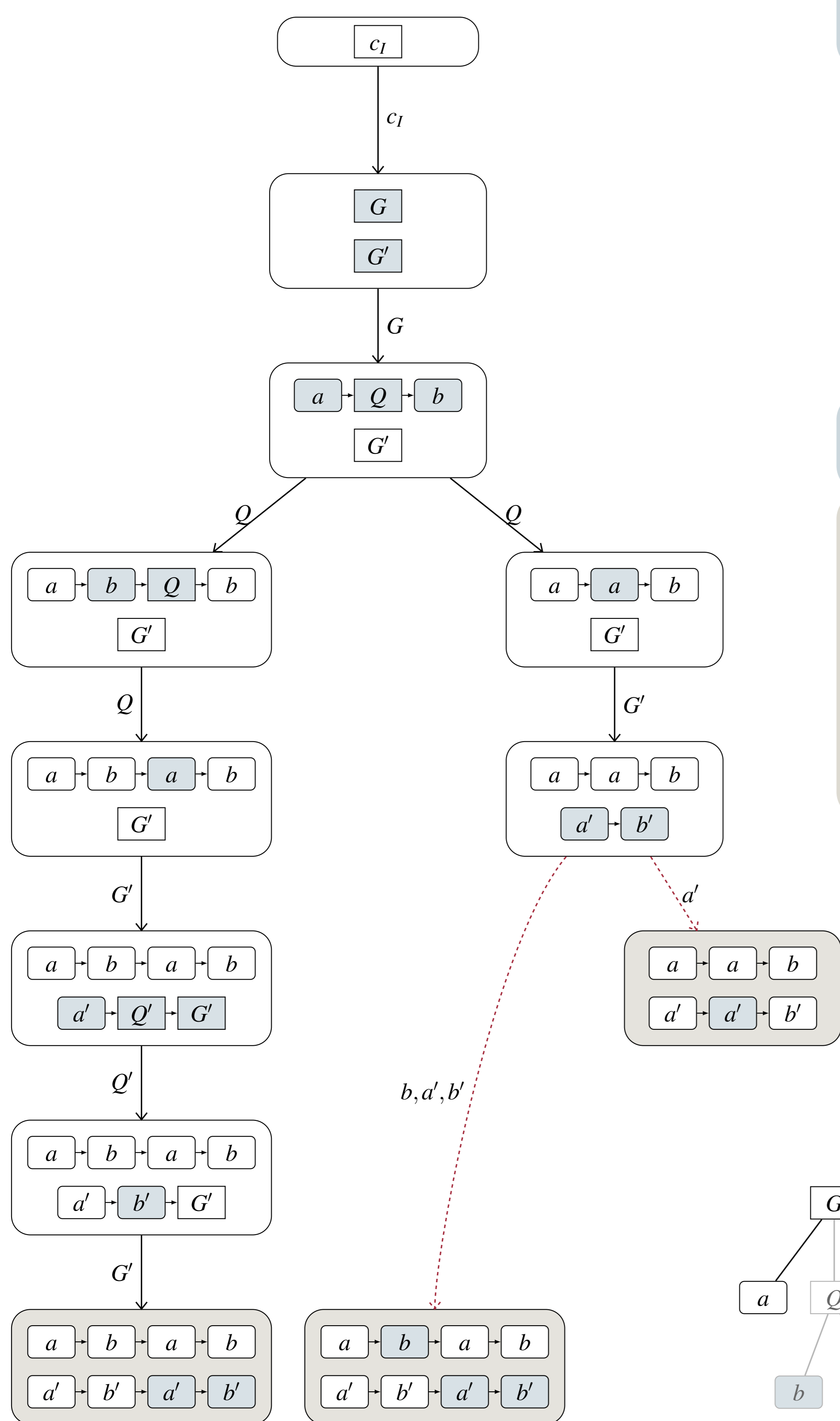


Figure 1: a search space fragment



Theorem

The plan existence problem is decidable for hybrid planning.

Proof Idea:

- establish an upper bound on the size of shortest solutions
- enumerate all short task networks and check whether they are a solution
- if no solution has been found, then we know that no solution exists at all

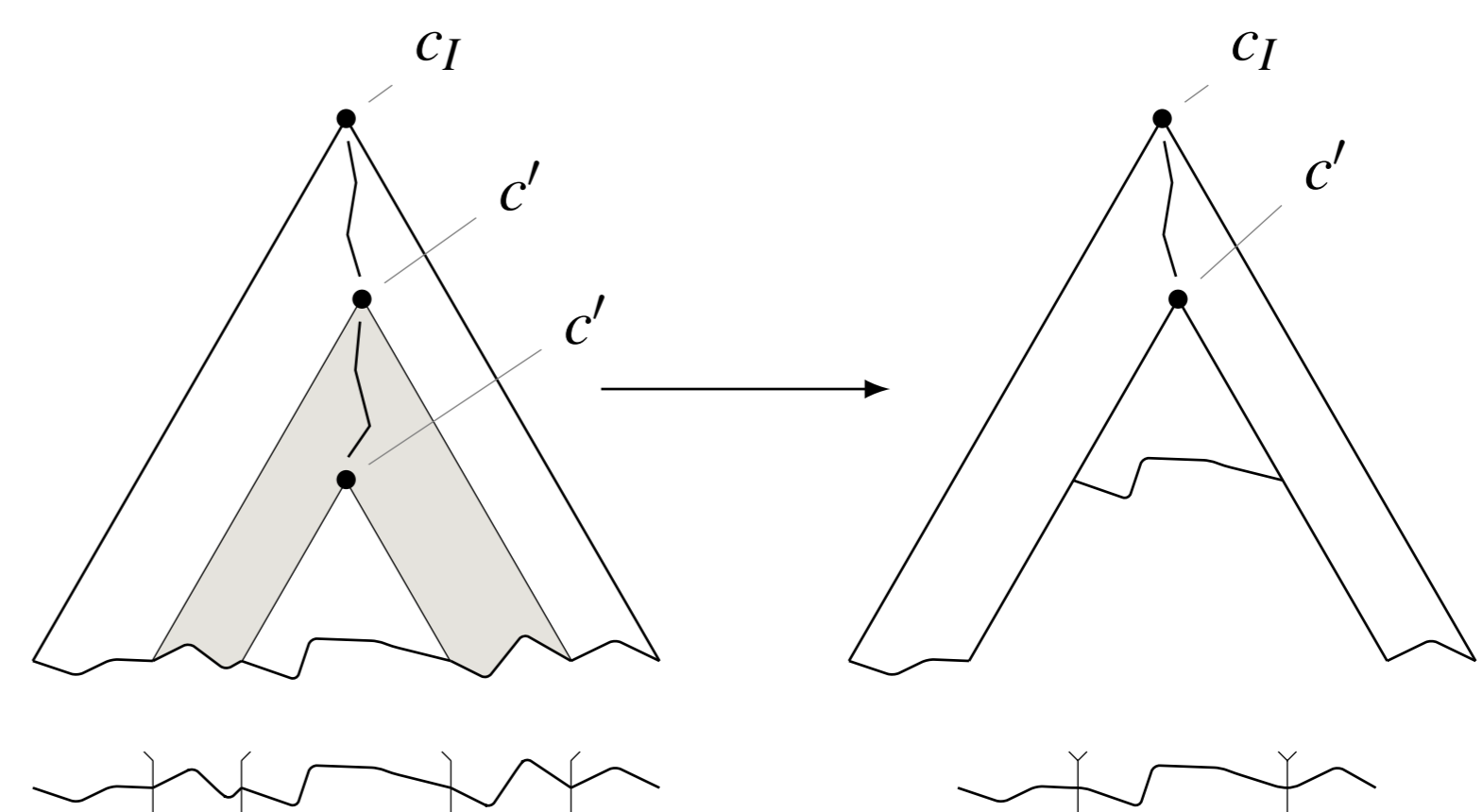
Proof — Bounding Decomposition

Any hybrid solution to P can be constructed using at most b^c decompositions where b is the number of tasks inside the largest method. (maximum branching factor of the decomposition tree)

We apply the idea of the pumping lemma for context-free grammars to task decomposition:

- remove all cycles from the decomposition tree
 - replace the removed elements using task and ordering insertion
- Right after decomposition the intermediate task network contains at most b^c tasks because the depth of the generating tree is limited by the number of compound tasks $|C|$.

Figure 5: pumping down decompositions



Proof — Bounding Task Insertion

Given a task network tn with n tasks, we have to insert at most $n2^{|V|}$ additional tasks to turn it into a solution, if this is possible at all.

- to obtain the bound, fix a (totally ordered) solution that contains tn ; thus it can be obtained from tn via insertion
- remove all task sequences that produce loops in the state space and that do not contain tasks from tn
- we obtain a solution which contains at most $n(2^{|V|} + 1)$ tasks

Proof — Bounding Solutions

For every planning problem P , there exists a hybrid solution with at most $b^c(2^{|V|} + 1)$ tasks if such a solution exists at all.

The bound follows from the bounds on task decomposition and the bound on task insertion.

Conclusion & Discussion

- allowing task insertion makes HTN planning decidable resulting in the following complexity classes [1, 3, 4]:

Computational Complexity			
classical	hybrid	HTN	
lifted	EXPSpace-complete	EXPSpace-hard	RE
propositional	PSPACE-complete	PSPACE-hard \cap EXPSpace	RE

- one can think of HTN planning with insertion having a different meaning than HTN planning
 - HTN planning: specifies, what **must not** be in a plan
 - Hybrid planning: specifies, what **has to be** in a plan

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