



OWL The Web Ontology Language

part I: overview

Pavel Klinov





In the next 90 mins

Intro

- why OWL?
- relationship to RDF(S) and logics (DLs)

Application areas and tools

Basics

- entities, expressions, axioms
- dealing with data values
- non-logical part: annotations, imports, and versions



ulm university universität UUUM So why Semantic Web needs OWL?

First, we've got RDF

- A simple graph language to express facts (LD)
- A simple data model + low-level data integration tools



So why Semantic Web needs OWL?

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First, we've got RDF

- A simple graph language to express facts (LD)
- A simple data model + low-level data integration tools No schema? But we have RDFS!
- A lightweight schema, good for simple vocabularies
- Some simple inferencing (transitivity of rdfs:subClassOf)





Schemas are often just too weak

- Can say: :hasWife rdfs:domain :Woman (rdfs:range :Woman)
- Cannot say: :Peter :hasWife some :Woman (only :Woman)





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So we need a language(s) that:

- Provides adequate balance between expressivity and computational complexity
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So we need a language(s) that:

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That language is called OWL 2





Application areas

Vocabulary-centric applications

- manage complex terminologies (in a machine-processable way)
- share terminology across applications

Data-centric applications

- lightweight reasoning over tons of data
- data integration





Terminology management

Example: medical informatics

- Terminologies are huge
 - ICD: ~100K medical codes
 - SNOMED CT: >300K classes
- Applications are different
 - medical diagnosis tools
 - electronic medical records
 - learning and statistical analysis tools

Scalable schema reasoning is key

• e.g. for quality assurance

All must agree on the meaning of the terms





Data-centric apps

Mostly about querying loads of data

- w.r.t. some (simple) schema
- on top of RDF (or SQL) database







Data integration

Data sources are often heterogeneous

- RDF data
- relational data
- spreadsheets









Data integration

Data sources are often heterogeneous

• RDF data







One size does not fit all!







Reasoners

- DL: Pellet, FaCT++, HermiT, RACER
- Lightweight: CEL/jCEL/ELK/Quest

Semantic databases

- Virtuoso, Stardog, OWLIM, Oracle 11
- not always fully implement profiles
- APIs: OWL API, RDF-based APIs (Jena, Sesame, etc.)
- Data integration (PDQ)
- Matchers, editors, debuggers, visualizers...





Extended RDF or logic?

OWL as RDF extension

- Every OWL ontology can be expressed as RDF graph (the other way is trickier)
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- OWL as a logic with a Web-friendly syntax
- OWL ontology is a DL knowledge base
- with a DL semantics





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- Every OWL ontology can be expressed as RDF graph (the other way is trickier)
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- OWL ontology is a DL knowledge base
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The views are compatible to a certain extent

We adopt the second view in this lecture





Schema vs Data

Think RDB

- schema defines structure (tables, keys, attributes)
- data specifies facts





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OWL

- schema (TBox) statements describe the domain
- data (ABox) statements express facts (like RDF)
- both are called axioms
- TBox + ABox is called ontology





Modeling example

Family

- parent, children
- cousins, aunts, uncles, nephews, etc.
- pets

Need to model to define terms unambiguously

- to manage data
- to make apps understand the data
- to make sure different apps agree on terms





A simple example (TBox, ABox)

TBox: conceptual modeling

- a parent is a mother or a father
- father and mother are disjoint concepts
- every person must have one parent of each kins
- your parents' parents are your grandparents





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ABox: a specific family

- Peter is a father, Lois is a mother
- Peter and Lois are parents of Chris, Meg, and Stewie
- Pewterschmidts are parents of Lois







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 - explicitly true: Peter is a father
 - implicitly true: Chris, Meg, and Stewie have grandparents
 - false: Lois is a father
 - unknown: Chris is a parent





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 - explicitly true: Peter is a father
 - implicitly true: Chris, Meg, and Stewie have grandparents
 - false: Lois is a father
 - unknown: Chris is a parent
- no unique name assumption: Chris, Meg, and Stewie could all denote the same person





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- modeling is declarative, describes what's true
- no procedural semantics (triggers, slots, etc.)
- doesn't specify how to infer what's true





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- modeling is declarative, describes what's true
- no procedural semantics (triggers, slots, etc.)
- doesn't specify how to infer what's true
- OWL is not a schema language
 - can't impose syntactic constraints on documents (e.g. like in XML Schema)
 - example: can't require that parent axioms are syntactically present





On OWL syntaxes

There are many:

- **RDF**-native: RDF/XML, Turtle, N3, etc. all describe triples
- OWL-native: OWL/XML, Functional, Manchester all describe axioms
- This lecture uses the Functional Syntax
 - avoid OWL axiom to RDF triples mapping
 - avoid XML verbosity





Axioms, Entities, and Expressions





Entities

Main building blocks: classes, properties, individuals (all denoted with IRIs)

• Individuals: specific objects

:Peter, :Lois, etc.

- Classes (concepts): sets of individuals
 :Family, :Parent
- Properties (roles): sets of pairs of individuals
 :marriedTo, :childOf





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Entities need to be declared in OWL 2 DL

Declaration (ObjectProperty (:hasParent))





Class expressions (CE)

Classes with a IRI are called named or atomic

:Person, :Parent, ...

owl:Thing (\top) and owl:Nothing (\bot) are predefined

Can be combined into class expressions

- expressions don't have IRIs
- still interpreted as sets
- propositional and non-propositional





Property expressions

Named properties

- identified with IRI
- owl:topObjectProperty and owl:bottomObjectProperty
- object properties and data properties

Property expressions

- no IRIs
- also interpreted as relations





CE: boolean constructors

OWL 2 DL is a propositionally complete language intersection: ObjectIntersectionOf(:Woman :Parent) union: ObjectUnionOf(:Mother :Father) complement: ObjectComplementOf(:Parent : Mother)







CE: restrictions on properties

Existentials:

ObjectSomeValuesFrom(:hasChild :Man)






CE: restrictions on properties

Existentials:

ObjectSomeValuesFrom(:hasChild :Man)







CE: restrictions on properties

Universals:

ObjectAllValuesFrom(:hasChild :Woman)







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ObjectAllValuesFrom(:hasChild :Woman)







Nominals classes

Sometimes you just want to enumerate things ObjectOneOf(:Chris :Meg :Stewie)







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Sometimes you just want to enumerate things ObjectOneOf(:Chris :Meg :Stewie)



Does it mean that the class

- contains exactly 3 objects?
- at least 3? at most 3?





Self-restrictions

Can define a class of objects related to itself ObjectHasSelf(:likes)







Cardinality restrictions

ObjectMinCardinalityFrom (2:hasChild owl:Thing) ObjectMaxCardinalityFrom (2:hasParent:Parent)





Property restrictions

Inverse properties

ObjectInverseOf(:hasChild)

interpreted as inverse relations







Property restrictions

Property chains

ObjectPropertyChain(:hasParent :hasSibling)

interpreted as compositions of relations







Axioms

TBox axioms

- relationships between classes (e.g. inclusion)
- relationships between properties

ABox axioms

- class membership
- property membership
- individual equality/inequality





Class inclusions

- SubClassOf(:Woman :Person)
- SubClassOf(
 - :Grandfather
 - ObjectIntersectionOf(:Man :Parent))





Class equivalence

EquivalentClasses(

:Mother

ObjectIntersectionOf(:Woman :Parent))

- all mothers are women and parents
- vice versa





Class disjointness

DisjointClasses(:Father :Mother)

no instance of A is an instance of B (and vice versa)





Property axioms

Property inclusions

simple: SubObjectPropertyOf(:hasWife :hasSpouse)
chains:

SubObjectPropertyOf(

ObjectPropertyChain(

:hasParent :hasParent) :hasGrandparent)





Property axioms

FunctionalObjectProperty(:hasMother) InverseFunctionalObjectProperty(:motherOf) ReflexiveObjectProperty(:likes) IrreflexiveObjectProperty(:hates) TransitiveObjectProperty(:partOf) SymmetrycObjectProperty(:hasSpouse) AsymmetricObjectProperty(:hasChild)





The feature set isn't minimal

Existentials and universals

- ObjectSomeValuesFrom(:hasChild :Person)
- ObjectAllValuesFrom(:hasChild :Person)

Class equivalence and disjointness (trivial) Transitivity?





The feature set isn't minimal

Existentials and universals

- ObjectSomeValuesFrom(:hasChild :Person)
- ObjectAllValuesFrom(:hasChild :Person)
- Class equivalence and disjointness (trivial) Transitivity?
- Or even ABoxes?
 - SubClassOf(ObjectOneOf(:Stewie) :Person)
 - ClassAssertion(:Person :Stewie)





Where are we?

Parts we've covered

- entities, class expressions
- object properties

Next

- data types and data properties
- very similar to classes and object properties!
- Later: non-logical part
- imports
- versions
- annotations





OWL and data values

- OWL is a two-sorted language
 - The abstract domain
 - classes, properties, named objects
 ObjectPropertyAssertion(:fatherOf :Peter :Meg)





OWL and data values

OWL is a two-sorted language

- The abstract domain
 - classes, properties, named objects
 ObjectPropertyAssertion(:fatherOf :Peter :Meg)
- The concrete (or data) domain
 - strings, numbers, dates, etc.

DataPropertyAssertion(:hasAge :Meg "17"^^xsd:integer)





Abstract and data domains

Abstract domain: Δ

- non-empty and arbitrary
- finite or infinite

Data domain Δ

- a superset of standard value sets (e.g., integers)
- fixed!

The domains are disjoint





The abstract world of logic

For developing theories about the world

- modelers often cautious and pedantical
- Open World Assumption, no Unique Name Assumption
- instances are defined only by axioms





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Makes sense because

- usually better to under-model than to over-model
 - under-modeling loses entailments
 - over-modeling introduces errors
- gives extra flexibility





The concrete, data world

For re-using standard data theories

- have excellent theories about numbers, etc.
- know how to use them, how to compute with them
- don't need custom, ill-made integer ontologies!
- don't need UNA





The concrete, data world

For re-using standard data theories

- have excellent theories about numbers, etc.
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- don't need custom, ill-made integer ontologies!
- don't need UNA

Datatypes fix what we know about, e.g., integers

- "4"^^xsd:integer and "6"^^integer aren't equal
- because all names have fixed meaning
 - somewhat like owl:'Thing
 - except that the concrete domain is always the same





Connecting the worlds

Data properties

- map abstract individuals to concrete data values
- DataSomeValueFrom(:hasWeight

:Peter "100"^^xsd:integer)

Semantics

• interpreted as subsets of $\Delta x \Delta d$





Data axioms

Axioms (mostly as for object properties)

- equivalence, inclusion, disjointness
- domains and ranges
- assertions
- functionality

But

- no chains (even transitivity)
- no inverses, reflexivity, symmetry
- can't go the other way or break the separation





More on fixed semantics

Example:

DataPropertyAssertion(:hasAge

```
:Meg "17"^^xsd:integer)
```

```
DataPropertyAssertion(:hasAge
```

```
:Meg "16"^^xsd:integer)
```

FunctionalDataProperty(:hasAge)

This is inconsistent

Try to formalize this logically!

- Remember UNA
- DifferentIndividuals(17 16)?





More on datatypes

Datatype: a kind of data values (integers, strings)

- IRI
- lexical space: "str", "1"^^xsd:integer, "01"^^xsd:integer
- value space: "str", 1
- facet spaces: pairs (F, v), mapped to a subset of VS
 - F: constraining facet
 - v: constraining value
 - (xsd:minExclusive , "30"^^xsd:integer)

Datatype map: a particular set of datatypes

- for a language
- for a particular tool (reasoner)





The OWL 2 datatype map

XSD datatypes

- decimals, integers (and subtypes)
- xsd:float
- xsd:double
- strings (subtypes of rdf:PlainLiteral)

Nuances:

- decimals and integers are subtypes of owl:real
- which is pairwise disjoint with xsd:float and xsd:double
 DataPropertyRange(a:hasAge xsd:integer)
 DataPropertyAssertion(a:hasAge a:Meg "17"^^xsd:double)





Data ranges

Abstract world analogue: class expressions

Can define custom datatypes based on standard ones DataRange

- Datatype (like xsd:integer)
- DataUnionOf, DataIntersectionOf, DataComplementOf
 DataUnionOf(xsd:string xsd:integer)
- DataOneOf

DataOneOf("1"^^xsd:integer "2.5"^^xsd:double)

DatatypeRestriction





Datatype restrictions

Can constrain a datatype using facets DatatypeRestriction(DT F1 V1 ... F2 Vn)

Example:

DatatypeRestriction(xsd:integer

xsd:minInclusive "5"^^xsd:integer xsd:maxExclusive "10"^^xsd:integer)

contains only 5, 6, 7, 8, 9

facets are combined conjunctively





Datatype definitions

Can assign names to custom (restricted) datatypes DatatypeDefinition(DT DR) Example:

DatatypeDefinition(

:email

DatatypeRestriction(xsd:string xsd:pattern "..."))

Now can use :email in data axioms:

DataPropertyRange(:hasEmail :email)





Identifying abstract individuals

What if we need to identify objects by their "attributes"?

For object property values

 use inverse functional properties InverseFunctionalProperty(:hasName)





Identifying abstract individuals

What if we need to identify objects by their "attributes"?

For object property values

use inverse functional properties
 InverseFunctionalProperty(:hasName)

Problems:

- global inverse functionality often undesirable (name's only unique within the Griffin family)
- how about data properties?
 no inverse functional data properties





Keys

HasKey(CE (OPE1 ... OPEm) (DPE1 ... DPEn)) This says that:

- if two named individuals of CE coincide on ...
- ... values of all object properties ...
- ... and values of all data properties, then
- the individuals are identical

Example:

- HasKey(:GriffinFamily(:hasName)())
- HasKey(owl:Thing () (:hasTaxId))





Where are we?

Covered the logical part

- abstract part (class expressions, object properties)
- data part (datatypes, data ranges, data properties)
- axioms

Next: non-logical part

- imports
- versions
- annotations





Imports

Ontologies are meant to be reusable OWL supports knowledge reuse via importing

Ontology(<http://fox.com/familyguy>
Import(<http://example.org/families.owl>))





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Particularly important in HCLS, biology, etc.

- pros: reuse other people efforts
- cons: can be too heavyweight solution: modularity (on Friday)





Versions

Ontologies are identified with a IRI

but also may have a version IRI to distinguish versions

Ontology(<http://fox.com/familyguy>

<http://fox.com/familyguy/2.0>

Why?

- ontologies are like public APIs (for your or shared data)
- changing your ontology may break others





Annotations

Not all content has to be logical

Meta-information

- author info
- axiom labels, comments
- provision

Modeling these on the logical level is unnecessary

- aren't statements about the domain
- statements about statements about the domain!

OWL 2 provides annotation support for these





Annotations

Subjects: ontologies or entities

Assertion: <annotationProperty, annotationValue> Values: IRIs, literals, or individuals Examples:

- AnnotationAssertion(rdfs:label a:Peter "Represents the main character from Family Guy")
- Ontology(<http://fox.com/familyguy>

Annotation(rdfs:label "A Family Guy ontology") Often useful for i18n





End of the basics

questions?