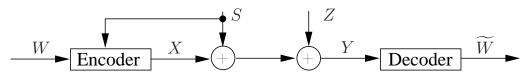


Exercises for Applied Information Theory

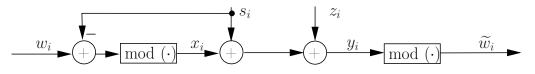
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Task 8.1 (Tomlinson-Harashima Precoding)

Consider the channel Y = X + S + Z depicted below, where both S and Z are noise terms and S is known to the transmitter.



Now, the channel input X can be encoded to optimize the transmission of a binary message $W \in \{-1, +1\}$ to the receiver. For this purpose, we will use Tomlinson-Harashima precoding and assume Z = 0 at all times. The encoding and decoding for the Tomlinson-Harashima scheme is done as follows:



Thus, the channel input is chosen such that

$$x_i = (w_i - s_i) \bmod q,$$

while the output of the decoder is chosen such that

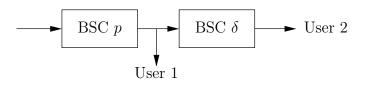
$$\widetilde{w}_i = y_i \mod q.$$

- a) How is q selected in the Tomlinson-Harashima precoding scheme?
- b) Given the input sequence $\mathbf{w} = [-1, -1, +1, -1, +1, +1, -1]$ and the known interference sequence $\mathbf{s} = [3.2, -2.4, 0.3, -1.8, -4.2, -1, 1.3, 1.5]$, calculate the transmit sequence \mathbf{x} , the receive sequence \mathbf{y} and the decoded sequence $\tilde{\mathbf{w}}$
- c) What is the average transmit power of the sequence \mathbf{x} ? Compare this to the transmit power without the modulo operation, i.e., $x_i = w_i s_i$.



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Task 8.2 (Degraded Broadcast Channel, BSC)



We consider the degraded broadcast channel consisting of a cascade of two BSCs with error probabilities p and δ , respectively.

In the broadcast case, a superposition code shall be used

$$\underline{c} = \underline{c_1} + \underline{c_2}$$

For long codewords, the addition of \underline{c}_1 can be interpreted as an additional BSC with error rate $\gamma = P(1)$.

- a) Calculate the capacity C_1 for user 1 and C_2 for user 2 for the case, that only user 1 or user 2 use the channel exclusively.
- b) Show that the rate region that is obtained by applying the superposition code is given by

$$R_1 \leq h(\gamma(1-p) + (1-\gamma)p) - h(p)$$

$$R_2 \leq 1 - h(\gamma(1-\epsilon) + (1-\gamma)\epsilon) \text{ with } \epsilon = (1-p)\delta + p(1-\delta)$$

Hint: Assume that user 1 can decode \underline{c}_2 correctly, see script section 6.3