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Solution to Task 8.1 (Tomlinson-Harashima Precoding)

- a) In Tomlinson-Harashima precoding the input alphabet is extended, such that the points are equidistant. Here, we have $W \in \{-1, +1\}$, i.e., the extension is such that -1 is represented by $\ldots, -5, -1, 3, 7, \ldots$ while +1 is represented by $\ldots, -7, -3, 1, 5, \ldots$. Thus, the interval we are repeating ranges from -2 to 2. As q has to be chosen as the length of this interval we have q = 4.
- b) The output range of the modulo functions in THP is chosen as $\left[-\frac{q}{2}, \frac{q}{2}\right)$, so that we obtain the following sequences

Thus, the sequence $\widetilde{\mathbf{w}}$ is equivalent to the input sequence \mathbf{w} .

c) The average transmit power of the transmit vector x is given by

$$\frac{1}{8}\sum_{i=1}^8 |x_i|^2 = 1.3638$$

If the modulo operation is omitted, the transmit vector would be

$$\mathbf{w} - \mathbf{s} = [-4.2, 1.4, 0.7, 0.8, 5.2, 2.0, -0.3, -2.5],$$

which has an average power of

$$\frac{1}{8}\sum_{i=1}^{8}|w_i - s_i|^2 = 7.2638.$$

Thus, the modulo operation reduces the transmit power enormously.



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Solution to Task 8.2 (Degraded Broadcast Channel, BSC)

a) A cascade of two BSCs can be transformed into a single BSC:



with $\epsilon = (1 - p)\delta + p(1 - \delta)$. Thus $C_1 = 1 - h(p)$ and $C_2 = 1 - h(\epsilon)$.

b) User 2:

We treat \underline{c}_1 as an additional error on \underline{c}_2 (see script, section 6.3), i.e. an additional BSC with error probability $P(1) = \gamma$. This means we again have a cascade of two BSCs with error probabilities γ and ϵ , which gives us a BSC with error probability $(1 - \gamma)\epsilon + \gamma(1 - \epsilon)$. $\Rightarrow R_{2,BC} \leq 1 - h((1 - \gamma)\epsilon + \gamma(1 - \epsilon))$

User 1: For superpositioning, the channel input can be written as

$$X = X_1 + X_2,$$

where X_i is the bit that shall be transmitted to user *i* and $f_{X_i}(1) = \gamma$. As user 1 has a better reception than user 2, it can decode \underline{c}_2 and thus calculate

$$\widetilde{Y}_1 = Y_1 - X_2 = X_1 + Z,$$

where Y_1 is the received signal of user 1 and Z is the possible bit error from the first BSC with $f_Z(1) = p$.

The entropy of \widetilde{Y}_1 is calculated as

$$f_{\widetilde{Y}_{1}}(1) = f_{X_{1}}(1)f_{Z}(0) + f_{X_{1}}(0)f_{Z}(1) = (1-p)\gamma + (1-\gamma)p =: q,$$

$$f_{\widetilde{Y}_{1}}(0) = 1 - q,$$

$$H(\widetilde{Y}_{1}) = h(q).$$



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And together with

$$H(\widetilde{Y}_1|X_1) = H(Z) = h(p)$$

we get

$$R_{1,BC} \le I(X_1, \widetilde{Y}_1) = H(\widetilde{Y}_1) - H(\widetilde{Y}_1|X_1) = h(q) - h(p) = h(\gamma(1-p) + (1-\gamma)p) - h(p).$$