



# Error Correction for Physical Unclonable Functions Using Generalized Concatenated Codes

Sven Müelich\*, Sven Puchinger\*, Martin Bossert\*, Matthias Hiller<sup>⋄</sup>, Georg Sigl<sup>⋄</sup>

\*Institute of Communications Engineering, Ulm University, Germany

Institute for Security in Information Technology, TU Munich, Germany

ACCT Svetlogorsk, September 7-13, 2014

# Outline



- Motivation
- Physical Unclonable Functions (PUFs)
- 3 Example Code Construction
- 4 Conclusion

## Motivation



## Challenges when implementing a cryptosystem:

- Secure key generation
  - Random, unique and unpredictable keys
  - Satisfying these properties is hard to achieve
- Secure key storage
  - Key bits in a non-volatile memory
  - Adversaries can gain physical access to (protected) memories

Physical Unclonable Functions (PUFs) can be used to realize secure key generation and secure key storage

# Physical Unclonable Functions (PUFs)



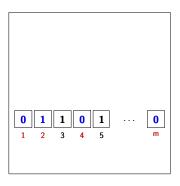
#### What is a PUF?

- Physical entity with challenge-response behavior
- Properties:
  - Uniqueness
  - Reproducibility



# Physical Unclonable Functions (PUFs)





## **Example:** SRAM PUFs

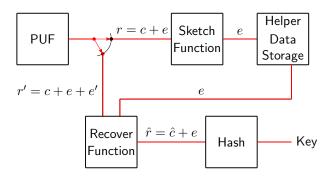
- device with memory cells
- random initialization when powering on
- randomness static over lifetime
- Challenge: Subset of memory cells
- Response: Values in selected memory cells

# Physical Unclonable Functions (PUFs)



## Why coding theory?

 Responses are not perfectly reproducible and hence cannot be used as key directly



# **Example Code construction**



#### Challenge: Find good codes for Secure Sketches

#### **Constraints:**

- Time and area consumption
- Binary codes
- Dimension ≥ key length
- Codelength as small as possible

# **Example Code construction**



# Existing scheme given in [Maes2012]<sup>1</sup>:

- Binary Symmetric Channel with p = 0.14
- Generate 128 bit key with block error probability  $P_{\rm err}=10^{-9}$
- $\bullet$  Concatenation of (318, 174, 35) BCH code and (7, 1, 7) code

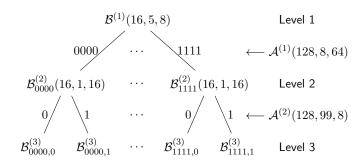
## What is our goal?

- Generate 128 bit key with block error probability  $P_{\rm err} < 10^{-9}$
- Improve existing scheme in
  - Codelength
  - Block error probability
  - Simple implementation

<sup>&</sup>lt;sup>1</sup>R. Maes, A. Herrewege, I. Verbauwhede, "PUFKY: A Fully Functional PUF-Based Cryptographic Key Generator", CHES, 2012

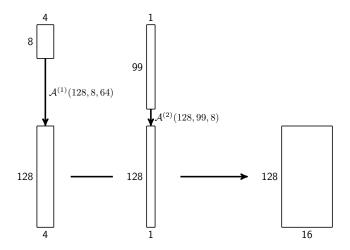


#### Partitioning:



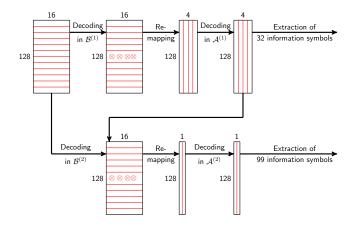


## **Encoding:**





## **Decoding:**





#### Used decoding methods:

- Generalized Concatenated Codes (GC Codes)
- RM Error Erasure Decoding
- Generalized Minimum Distance (GMD) Decoding

# Conclusion



## How good is our code construction?

Code	$P_{err}$	Length	Largest Field
[Maes2012]	$\approx 10^{-9}$	2226	$\mathbb{F}_{2^8}$ (BCH)
New	$\approx 5.37 \cdot 10^{-10}$	2048	$\mathbb{F}_2$

Thank you for your attention.