

Unifying Lazy and Strict Computations

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```
theory ULSC
```

```
imports Main
```

```
begin
```

```
class mult = times
```

```
begin
```

```
notation
```

```
  times (infixl · 70) and
```

```
  times (infixl ; 70)
```

```
end
```

```
class neg = uminus
```

```
begin
```

```
no-notation
```

```
  uminus (− - [81] 80)
```

```
notation
```

```
  uminus (− - [80] 80)
```

```
end
```

context order

begin

definition *isotone* :: ('a ⇒ 'a) ⇒ bool
where *isotone* f ↔ (∀ x y . x ≤ y → f(x) ≤ f(y))

definition *is-fixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-fixpoint* f x ↔ f(x) = x
definition *is-prefixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-prefixpoint* f x ↔ f(x) ≤ x
definition *is-postfixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-postfixpoint* f x ↔ f(x) ≥ x
definition *is-least-fixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-least-fixpoint* f x ↔ f(x) = x ∧ (∀ y . f(y) = y → x ≤ y)
definition *is-greatest-fixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-greatest-fixpoint* f x ↔ f(x) = x ∧ (∀ y . f(y) = y → x ≥ y)
definition *is-least-prefixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-least-prefixpoint* f x ↔ f(x) ≤ x ∧ (∀ y . f(y) ≤ y → x ≤ y)
definition *is-greatest-postfixpoint* :: ('a ⇒ 'a) ⇒ 'a ⇒ bool **where** *is-greatest-postfixpoint* f x ↔ f(x) ≥ x ∧ (∀ y . f(y) ≥ y → x ≥ y)
definition *has-fixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-fixpoint* f ↔ (∃ x . *is-fixpoint* f x)
definition *has-prefixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-prefixpoint* f ↔ (∃ x . *is-prefixpoint* f x)
definition *has-postfixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-postfixpoint* f ↔ (∃ x . *is-postfixpoint* f x)
definition *has-least-fixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-least-fixpoint* f ↔ (∃ x . *is-least-fixpoint* f x)
definition *has-greatest-fixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-greatest-fixpoint* f ↔ (∃ x . *is-greatest-fixpoint* f x)
definition *has-least-prefixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-least-prefixpoint* f ↔ (∃ x . *is-least-prefixpoint* f x)
definition *has-greatest-postfixpoint* :: ('a ⇒ 'a) ⇒ bool **where** *has-greatest-postfixpoint* f ↔ (∃ x . *is-greatest-postfixpoint* f x)
definition *the-least-fixpoint* :: ('a ⇒ 'a) ⇒ 'a (μ - [201] 200) **where** μ f = (THE x . *is-least-fixpoint* f x)
definition *the-greatest-fixpoint* :: ('a ⇒ 'a) ⇒ 'a (ν - [201] 200) **where** ν f = (THE x . *is-greatest-fixpoint* f x)
definition *the-least-prefixpoint* :: ('a ⇒ 'a) ⇒ 'a (pμ - [201] 200) **where** pμ f = (THE x . *is-least-prefixpoint* f x)
definition *the-greatest-postfixpoint* :: ('a ⇒ 'a) ⇒ 'a (pν - [201] 200) **where** pν f = (THE x . *is-greatest-postfixpoint* f x)

lemma *least-fixpoint-unique*: *has-least-fixpoint* f → (∃!x . *is-least-fixpoint* f x)
by (smt antisym *has-least-fixpoint-def is-least-fixpoint-def*)

lemma *greatest-fixpoint-unique*: *has-greatest-fixpoint* f → (∃!x . *is-greatest-fixpoint* f x)
by (smt antisym *has-greatest-fixpoint-def is-greatest-fixpoint-def*)

lemma *least-prefixpoint-unique*: *has-least-prefixpoint* f → (∃!x . *is-least-prefixpoint* f x)
by (smt antisym *has-least-prefixpoint-def is-least-prefixpoint-def*)

lemma *greatest-postfixpoint-unique*: *has-greatest-postfixpoint* f → (∃!x . *is-greatest-postfixpoint* f x)
by (smt antisym *has-greatest-postfixpoint-def is-greatest-postfixpoint-def*)

lemma *least-fixpoint*: *has-least-fixpoint* f → *is-least-fixpoint* f (μ f)

proof

assume *has-least-fixpoint* f
hence *is-least-fixpoint* f (THE x . *is-least-fixpoint* f x)
by (smt *least-fixpoint-unique theI'*)
thus *is-least-fixpoint* f (μ f)
by (simp add: *is-least-fixpoint-def the-least-fixpoint-def*)

qed

lemma *greatest-fixpoint*: $has\text{-}greatest\text{-}fixpoint\ f \rightarrow is\text{-}greatest\text{-}fixpoint\ f\ (\nu\ f)$

proof

assume *has-greatest-fixpoint* f

hence *is-greatest-fixpoint* f (*THE* x . *is-greatest-fixpoint* $f\ x$)

by (*smt greatest-fixpoint-unique theI'*)

thus *is-greatest-fixpoint* $f\ (\nu\ f)$

by (*simp add: is-greatest-fixpoint-def the-greatest-fixpoint-def*)

qed

lemma *least-prefixpoint*: $has\text{-}least\text{-}prefixpoint\ f \rightarrow is\text{-}least\text{-}prefixpoint\ f\ (p\mu\ f)$

proof

assume *has-least-prefixpoint* f

hence *is-least-prefixpoint* f (*THE* x . *is-least-prefixpoint* $f\ x$)

by (*smt least-prefixpoint-unique theI'*)

thus *is-least-prefixpoint* $f\ (p\mu\ f)$

by (*simp add: is-least-prefixpoint-def the-least-prefixpoint-def*)

qed

lemma *greatest-postfixpoint*: $has\text{-}greatest\text{-}postfixpoint\ f \rightarrow is\text{-}greatest\text{-}postfixpoint\ f\ (p\nu\ f)$

proof

assume *has-greatest-postfixpoint* f

hence *is-greatest-postfixpoint* f (*THE* x . *is-greatest-postfixpoint* $f\ x$)

by (*smt greatest-postfixpoint-unique theI'*)

thus *is-greatest-postfixpoint* $f\ (p\nu\ f)$

by (*simp add: is-greatest-postfixpoint-def the-greatest-postfixpoint-def*)

qed

lemma *least-fixpoint-same*: $is\text{-}least\text{-}fixpoint\ f\ x \rightarrow x = \mu\ f$

by (*metis least-fixpoint least-fixpoint-unique has-least-fixpoint-def*)

lemma *greatest-fixpoint-same*: $is\text{-}greatest\text{-}fixpoint\ f\ x \rightarrow x = \nu\ f$

by (*metis greatest-fixpoint greatest-fixpoint-unique has-greatest-fixpoint-def*)

lemma *least-prefixpoint-same*: $is\text{-}least\text{-}prefixpoint\ f\ x \rightarrow x = p\mu\ f$

by (*metis least-prefixpoint least-prefixpoint-unique has-least-prefixpoint-def*)

lemma *greatest-postfixpoint-same*: $is\text{-}greatest\text{-}postfixpoint\ f\ x \rightarrow x = p\nu\ f$

by (*metis greatest-postfixpoint greatest-postfixpoint-unique has-greatest-postfixpoint-def*)

lemma *least-fixpoint-char*: $is\text{-}least\text{-}fixpoint\ f\ x \leftrightarrow has\text{-}least\text{-}fixpoint\ f \wedge x = \mu\ f$

by (*metis least-fixpoint-same has-least-fixpoint-def*)

lemma *least-prefixpoint-char*: $is\text{-}least\text{-}prefixpoint\ f\ x \leftrightarrow has\text{-}least\text{-}prefixpoint\ f \wedge x = p\mu\ f$

by (*metis least-prefixpoint-same has-least-prefixpoint-def*)

lemma *greatest-fixpoint-char*: $is\text{-}greatest\text{-}fixpoint\ f\ x \leftrightarrow has\text{-}greatest\text{-}fixpoint\ f \wedge x = \nu\ f$

by (*metis greatest-fixpoint-same has-greatest-fixpoint-def*)

lemma *greatest-postfixpoint-char*: $is\text{-}greatest\text{-}postfixpoint\ f\ x \leftrightarrow has\text{-}greatest\text{-}postfixpoint\ f \wedge x = p\nu\ f$
by (*metis greatest-postfixpoint-same has-greatest-postfixpoint-def*)

lemma *least-prefixpoint-fixpoint*: $has\text{-}least\text{-}prefixpoint\ f \wedge isotone\ f \rightarrow is\text{-}least\text{-}fixpoint\ f\ (p\mu\ f)$
by (*smt eq-iff is-least-fixpoint-def is-least-prefixpoint-def isotone-def least-prefixpoint*)

lemma *pmu-mu*: $has\text{-}least\text{-}prefixpoint\ f \wedge isotone\ f \rightarrow p\mu\ f = \mu\ f$
by (*smt has-least-fixpoint-def is-least-fixpoint-def least-fixpoint-unique least-prefixpoint-fixpoint least-fixpoint*)

definition *lifted-less-eq* :: $('a \Rightarrow 'a) \Rightarrow ('a \Rightarrow 'a) \Rightarrow bool\ ((-\leq -)\ [51, 51]\ 50)$
where $f \leq g \leftrightarrow (\forall x. f(x) \leq g(x))$

lemma *lifted-reflexive*: $f = g \rightarrow f \leq g$
by (*metis lifted-less-eq-def order-refl*)

lemma *lifted-transitive*: $f \leq g \wedge g \leq h \rightarrow f \leq h$
by (*smt lifted-less-eq-def order-trans*)

lemma *lifted-antisymmetric*: $f \leq g \wedge g \leq f \rightarrow f = g$
by (*metis antisym ext lifted-less-eq-def*)

lemma *pmu-isotone*: $has\text{-}least\text{-}prefixpoint\ f \wedge has\text{-}least\text{-}prefixpoint\ g \wedge f \leq g \rightarrow p\mu\ f \leq p\mu\ g$
by (*smt is-least-prefixpoint-def least-prefixpoint lifted-less-eq-def order-trans*)

lemma *mu-isotone*: $has\text{-}least\text{-}prefixpoint\ f \wedge has\text{-}least\text{-}prefixpoint\ g \wedge isotone\ f \wedge isotone\ g \wedge f \leq g \rightarrow \mu\ f \leq \mu\ g$
by (*metis pmu-isotone pmu-mu*)

lemma *greatest-postfixpoint-fixpoint*: $has\text{-}greatest\text{-}postfixpoint\ f \wedge isotone\ f \rightarrow is\text{-}greatest\text{-}fixpoint\ f\ (p\nu\ f)$
by (*smt eq-iff is-greatest-fixpoint-def is-greatest-postfixpoint-def isotone-def greatest-postfixpoint*)

lemma *pnu-nu*: $has\text{-}greatest\text{-}postfixpoint\ f \wedge isotone\ f \rightarrow p\nu\ f = \nu\ f$
by (*smt has-greatest-fixpoint-def is-greatest-fixpoint-def greatest-fixpoint-unique greatest-postfixpoint-fixpoint greatest-fixpoint*)

lemma *pnu-isotone*: $has\text{-}greatest\text{-}postfixpoint\ f \wedge has\text{-}greatest\text{-}postfixpoint\ g \wedge f \leq g \rightarrow p\nu\ f \leq p\nu\ g$
by (*smt is-greatest-postfixpoint-def lifted-less-eq-def order-trans greatest-postfixpoint*)

lemma *nu-isotone*: $has\text{-}greatest\text{-}postfixpoint\ f \wedge has\text{-}greatest\text{-}postfixpoint\ g \wedge isotone\ f \wedge isotone\ g \wedge f \leq g \rightarrow \nu\ f \leq \nu\ g$
by (*metis pnu-isotone pnu-nu*)

lemma *mu-square*: $isotone\ f \wedge has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}least\text{-}fixpoint\ (f \circ f) \rightarrow \mu\ f = \mu\ (f \circ f)$
by (*metis antisym is-least-fixpoint-def isotone-def least-fixpoint-char least-fixpoint-unique o-apply*)

lemma *nu-square*: $isotone\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ (f \circ f) \rightarrow \nu\ f = \nu\ (f \circ f)$
by (*metis antisym is-greatest-fixpoint-def isotone-def greatest-fixpoint-char greatest-fixpoint-unique o-apply*)

lemma *mu-roll*: $isotone\ g \wedge has\text{-}least\text{-}fixpoint\ (f \circ g) \wedge has\text{-}least\text{-}fixpoint\ (g \circ f) \rightarrow \mu\ (g \circ f) = g(\mu\ (f \circ g))$

apply (*rule impI, rule antisym*)
apply (*smt is-least-fixpoint-def least-fixpoint o-apply*)
by (*smt is-least-fixpoint-def isotone-def least-fixpoint o-apply*)

lemma *nu-roll: isotone g \wedge has-greatest-fixpoint (f \circ g) \wedge has-greatest-fixpoint (g \circ f) \rightarrow ν (g \circ f) = g(ν (f \circ g))*

apply (*rule impI, rule antisym*)
apply (*smt is-greatest-fixpoint-def greatest-fixpoint isotone-def o-apply*)
by (*smt is-greatest-fixpoint-def greatest-fixpoint o-apply*)

lemma *mu-below-nu: has-least-fixpoint f \wedge has-greatest-fixpoint f \rightarrow μ f \leq ν f*

by (*metis is-greatest-fixpoint-def is-least-fixpoint-def least-fixpoint greatest-fixpoint*)

lemma *pmu-below-pnu-fix: has-fixpoint f \wedge has-least-prefixpoint f \wedge has-greatest-postfixpoint f \rightarrow $p\mu$ f \leq $p\nu$ f*

by (*smt has-fixpoint-def is-fixpoint-def is-greatest-postfixpoint-def is-least-prefixpoint-def le-less order-trans least-prefixpoint greatest-postfixpoint*)

lemma *pmu-below-pnu-iso: isotone f \wedge has-least-prefixpoint f \wedge has-greatest-postfixpoint f \rightarrow $p\mu$ f \leq $p\nu$ f*

by (*metis has-fixpoint-def is-fixpoint-def is-least-fixpoint-def least-prefixpoint-fixpoint pmu-below-pnu-fix*)

definition *galois :: ('a \Rightarrow 'a) \Rightarrow ('a \Rightarrow 'a) \Rightarrow bool*

where *galois l u \leftrightarrow ($\forall x y . l(x) \leq y \leftrightarrow x \leq u(y)$)*

lemma *galois-char: galois l u \leftrightarrow ($\forall x . x \leq u(l(x))$) \wedge ($\forall x . l(u(x)) \leq x$) \wedge isotone l \wedge isotone u*

apply (*rule iffI*)
apply (*smt galois-def isotone-def order-refl order-trans*)
by (*metis galois-def isotone-def order-trans*)

lemma *galois-closure: galois l u \rightarrow l(x) = l(u(l(x))) \wedge u(x) = u(l(u(x)))*

by (*smt antisym galois-char isotone-def*)

lemma *mu-fusion-1: galois l u \wedge isotone h \wedge has-least-prefixpoint g \wedge has-least-fixpoint h \wedge l(g(u(μ h))) \leq h(l(u(μ h))) \rightarrow l($p\mu$ g) \leq μ h*

proof

assume 1: *galois l u \wedge isotone h \wedge has-least-prefixpoint g \wedge has-least-fixpoint h \wedge l(g(u(μ h))) \leq h(l(u(μ h)))*

hence *l(u(μ h)) \leq μ h*

by (*metis galois-char*)

hence *g(u(μ h)) \leq u(μ h) using 1*

by (*smt galois-def least-fixpoint-same least-fixpoint-unique is-least-fixpoint-def isotone-def order-trans*)

thus *l($p\mu$ g) \leq μ h using 1*

by (*metis galois-def least-prefixpoint is-least-prefixpoint-def*)

qed

lemma *mu-fusion-2: galois l u \wedge isotone h \wedge has-least-prefixpoint g \wedge has-least-fixpoint h \wedge l \circ g \leq h \circ l \rightarrow l($p\mu$ g) \leq μ h*

by (*metis lifted-less-eq-def mu-fusion-1 o-apply*)

lemma *mu-fusion-equal-1: galois l u \wedge isotone g \wedge isotone h \wedge has-least-prefixpoint g \wedge has-least-fixpoint h \wedge l(g(u(μ h))) \leq h(l(u(μ h))) \wedge l(g($p\mu$ g)) = h(l($p\mu$ g)) \rightarrow μ h = l($p\mu$ g) \wedge μ h = l(μ g)*

by (*metis antisym least-fixpoint least-prefixpoint-fixpoint is-least-fixpoint-def mu-fusion-1 pmu-mu*)

lemma mu-fusion-equal-2: *galois l u ∧ isotone h ∧ has-least-prefixpoint g ∧ has-least-prefixpoint h ∧ $l(g(u(\mu h))) \leq h(l(u(\mu h))) \wedge l(g(p\mu g)) = h(l(p\mu g)) \rightarrow p\mu h = l(p\mu g) \wedge \mu h = l(p\mu g)$*

by (*smt antisym galois-char least-fixpoint-char least-prefixpoint least-prefixpoint-fixpoint is-least-prefixpoint-def isotone-def mu-fusion-1*)

lemma mu-fusion-equal-3: *galois l u ∧ isotone g ∧ isotone h ∧ has-least-prefixpoint g ∧ has-least-fixpoint h ∧ $l \circ g = h \circ l \rightarrow \mu h = l(p\mu g) \wedge \mu h = l(\mu g)$*

by (*smt mu-fusion-equal-1 o-apply order-refl*)

lemma mu-fusion-equal-4: *galois l u ∧ isotone h ∧ has-least-prefixpoint g ∧ has-least-prefixpoint h ∧ $l \circ g = h \circ l \rightarrow p\mu h = l(p\mu g) \wedge \mu h = l(p\mu g)$*

by (*smt mu-fusion-equal-2 o-apply order-refl*)

lemma nu-fusion-1: *galois l u ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-fixpoint h ∧ $h(u(l(\nu h))) \leq u(g(l(\nu h))) \rightarrow \nu h \leq u(p\nu g)$*

proof

assume 1: *galois l u ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-fixpoint h ∧ $h(u(l(\nu h))) \leq u(g(l(\nu h)))$*

hence $\nu h \leq u(l(\nu h))$

by (*metis galois-char*)

hence $l(\nu h) \leq g(l(\nu h))$ **using** 1

by (*smt galois-def greatest-fixpoint-same greatest-fixpoint-unique is-greatest-fixpoint-def isotone-def order-trans*)

thus $\nu h \leq u(p\nu g)$ **using** 1

by (*metis galois-def greatest-postfixpoint is-greatest-postfixpoint-def*)

qed

lemma nu-fusion-2: *galois l u ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-fixpoint h ∧ $h \circ u \leq u \circ g \rightarrow \nu h \leq u(p\nu g)$*

by (*metis lifted-less-eq-def nu-fusion-1 o-apply*)

lemma nu-fusion-equal-1: *galois l u ∧ isotone g ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-fixpoint h ∧ $h(u(l(\nu h))) \leq u(g(l(\nu h))) \wedge h(u(p\nu g)) = u(g(p\nu g)) \rightarrow \nu h = u(p\nu g) \wedge \nu h = u(\nu g)$*

by (*metis antisym greatest-fixpoint greatest-postfixpoint-fixpoint is-greatest-fixpoint-def nu-fusion-1 pnu-nu*)

lemma nu-fusion-equal-2: *galois l u ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-postfixpoint h ∧ $h(u(l(\nu h))) \leq u(g(l(\nu h))) \wedge h(u(p\nu g)) = u(g(p\nu g)) \rightarrow p\nu h = u(p\nu g) \wedge \nu h = u(p\nu g)$*

by (*smt antisym galois-char greatest-fixpoint-char greatest-postfixpoint greatest-postfixpoint-fixpoint is-greatest-postfixpoint-def isotone-def nu-fusion-1*)

lemma nu-fusion-equal-3: *galois l u ∧ isotone g ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-fixpoint h ∧ $h \circ u = u \circ g \rightarrow \nu h = u(p\nu g) \wedge \nu h = u(\nu g)$*

by (*metis nu-fusion-equal-1 o-apply order-refl*)

lemma nu-fusion-equal-4: *galois l u ∧ isotone h ∧ has-greatest-postfixpoint g ∧ has-greatest-postfixpoint h ∧ $h \circ u = u \circ g \rightarrow p\nu h = u(p\nu g) \wedge \nu h = u(p\nu g)$*

by (*smt nu-fusion-equal-2 o-apply order-refl*)

lemma mu-exchange-1: *galois l u ∧ isotone g ∧ isotone h ∧ has-least-prefixpoint (l ∘ h) ∧ has-least-prefixpoint (h ∘ g) ∧ has-least-fixpoint (g ∘ h) ∧ $l \circ h \circ g \leq g \circ h \circ l \rightarrow \mu(l \circ h) \leq \mu(g \circ h)$*

by (*smt galois-char is-least-prefixpoint-def isotone-def least-fixpoint-char least-prefixpoint least-prefixpoint-fixpoint mu-fusion-2 mu-roll o-apply o-assoc*)

lemma mu-exchange-2: *galois l u ∧ isotone g ∧ isotone h ∧ has-least-prefixpoint (l ∘ h) ∧ has-least-prefixpoint (h ∘ l) ∧ has-least-prefixpoint (h ∘ g) ∧ has-least-fixpoint (g ∘ h) ∧ has-least-fixpoint (h ∘ g) ∧ $l \circ h \circ g \leq g \circ h \circ l \rightarrow \mu(h \circ l) \leq \mu(h \circ g)$*

by (*smt galois-char isotone-def least-fixpoint-char least-prefixpoint-fixpoint mu-exchange-1 mu-roll o-apply*)

lemma mu-exchange-equal: *galois l u ∧ galois k t ∧ isotone h ∧ has-least-prefixpoint (l ∘ h) ∧ has-least-prefixpoint (h ∘ l) ∧ has-least-prefixpoint (k ∘ h) ∧ has-least-prefixpoint (h ∘ k) ∧*

$l \circ h \circ k = k \circ h \circ l \rightarrow \mu(l \circ h) = \mu(k \circ h) \wedge \mu(h \circ l) = \mu(h \circ k)$

by (*smt antisym galois-char isotone-def least-fixpoint-char least-prefixpoint-fixpoint lifted-reflexive mu-exchange-1 mu-exchange-2 o-apply*)

lemma *nu-exchange-1*: $\text{galois } l \ u \wedge \text{isotone } g \wedge \text{isotone } h \wedge \text{has-greatest-postfixpoint } (u \circ h) \wedge \text{has-greatest-postfixpoint } (h \circ g) \wedge \text{has-greatest-fixpoint } (g \circ h) \wedge g \circ h \circ u \leq \leq u \circ h \circ g \rightarrow \nu(g \circ h) \leq \nu(u \circ h)$

by (*smt galois-char is-greatest-postfixpoint-def isotone-def greatest-fixpoint-char greatest-postfixpoint greatest-postfixpoint-fixpoint nu-fusion-2 nu-roll o-apply o-assoc*)

lemma *nu-exchange-2*: $\text{galois } l \ u \wedge \text{isotone } g \wedge \text{isotone } h \wedge \text{has-greatest-postfixpoint } (u \circ h) \wedge \text{has-greatest-postfixpoint } (h \circ u) \wedge \text{has-greatest-postfixpoint } (h \circ g) \wedge \text{has-greatest-fixpoint } (g \circ h) \wedge \text{has-greatest-fixpoint } (h \circ g) \wedge g \circ h \circ u \leq \leq u \circ h \circ g \rightarrow \nu(h \circ g) \leq \nu(h \circ u)$

by (*smt galois-char isotone-def greatest-fixpoint-char greatest-postfixpoint-fixpoint nu-exchange-1 nu-roll o-apply*)

lemma *nu-exchange-equal*: $\text{galois } l \ u \wedge \text{galois } k \ t \wedge \text{isotone } h \wedge \text{has-greatest-postfixpoint } (u \circ h) \wedge \text{has-greatest-postfixpoint } (h \circ u) \wedge \text{has-greatest-postfixpoint } (t \circ h) \wedge \text{has-greatest-postfixpoint } (h \circ t) \wedge u \circ h \circ t = t \circ h \circ u \rightarrow \nu(u \circ h) = \nu(t \circ h) \wedge \nu(h \circ u) = \nu(h \circ t)$

by (*smt antisym galois-char isotone-def greatest-fixpoint-char greatest-postfixpoint-fixpoint lifted-reflexive nu-exchange-1 nu-exchange-2 o-apply*)

lemma *mu-commute-fixpoint-1*: $\text{isotone } f \wedge \text{has-least-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow \text{is-fixpoint } f \ (\mu(f \circ g))$

by (*metis is-fixpoint-def mu-roll*)

lemma *mu-commute-fixpoint-2*: $\text{isotone } g \wedge \text{has-least-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow \text{is-fixpoint } g \ (\mu(f \circ g))$

by (*metis is-fixpoint-def mu-roll*)

lemma *mu-commute-least-fixpoint*: $\text{isotone } f \wedge \text{isotone } g \wedge \text{has-least-fixpoint } f \wedge \text{has-least-fixpoint } g \wedge \text{has-least-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow (\mu(f \circ g) = \mu f \rightarrow \mu g \leq \mu f)$

by (*metis is-least-fixpoint-def least-fixpoint-same least-fixpoint-unique mu-roll*)

lemma *nu-commute-fixpoint-1*: $\text{isotone } f \wedge \text{has-greatest-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow \text{is-fixpoint } f \ (\nu(f \circ g))$

by (*metis is-fixpoint-def nu-roll*)

lemma *nu-commute-fixpoint-2*: $\text{isotone } g \wedge \text{has-greatest-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow \text{is-fixpoint } g \ (\nu(f \circ g))$

by (*metis is-fixpoint-def nu-roll*)

lemma *nu-commute-greatest-fixpoint*: $\text{isotone } f \wedge \text{isotone } g \wedge \text{has-greatest-fixpoint } f \wedge \text{has-greatest-fixpoint } g \wedge \text{has-greatest-fixpoint } (f \circ g) \wedge f \circ g = g \circ f \rightarrow (\nu(f \circ g) = \nu f \rightarrow \nu f \leq \nu g)$

by (*smt is-greatest-fixpoint-def greatest-fixpoint-same greatest-fixpoint-unique nu-roll*)

lemma *mu-diagonal-1*: $\text{isotone } (\lambda x . f \ x \ x) \wedge (\forall x . \text{isotone } (\lambda y . f \ x \ y)) \wedge \text{isotone } (\lambda x . \mu(\lambda y . f \ x \ y)) \wedge (\forall x . \text{has-least-fixpoint } (\lambda y . f \ x \ y)) \wedge \text{has-least-prefixpoint } (\lambda x . \mu(\lambda y . f \ x \ y)) \rightarrow \mu(\lambda x . f \ x \ x) = \mu(\lambda x . \mu(\lambda y . f \ x \ y))$

by (*smt is-least-fixpoint-def is-least-prefixpoint-def least-fixpoint-same least-fixpoint-unique least-prefixpoint least-prefixpoint-fixpoint*)

lemma *mu-diagonal-2*: $(\forall x . \text{isotone } (\lambda y . f \ x \ y) \wedge \text{isotone } (\lambda y . f \ y \ x) \wedge \text{has-least-prefixpoint } (\lambda y . f \ x \ y)) \wedge \text{has-least-prefixpoint } (\lambda x . \mu(\lambda y . f \ x \ y)) \rightarrow \mu(\lambda x . f \ x \ x) = \mu(\lambda x . \mu(\lambda y . f \ x \ y))$

by (*smt is-least-fixpoint-def is-least-prefixpoint-def isotone-def least-fixpoint-same least-prefixpoint least-prefixpoint-fixpoint*)

lemma *nu-diagonal-1*: $\text{isotone } (\lambda x . f \ x \ x) \wedge (\forall x . \text{isotone } (\lambda y . f \ x \ y)) \wedge \text{isotone } (\lambda x . \nu(\lambda y . f \ x \ y)) \wedge (\forall x . \text{has-greatest-fixpoint } (\lambda y . f \ x \ y)) \wedge \text{has-greatest-postfixpoint } (\lambda x . \nu(\lambda y . f \ x \ y)) \rightarrow \nu(\lambda x . f \ x \ x) = \nu(\lambda x . \nu(\lambda y . f \ x \ y))$

by (*smt is-greatest-fixpoint-def is-greatest-postfixpoint-def greatest-fixpoint-same greatest-fixpoint-unique greatest-postfixpoint greatest-postfixpoint-fixpoint*)

lemma *nu-diagonal-2*: $(\forall x . \text{isotone } (\lambda y . f x y) \wedge \text{isotone } (\lambda y . f y x) \wedge \text{has-greatest-postfixpoint } (\lambda y . f x y)) \wedge \text{has-greatest-postfixpoint } (\lambda x . \nu(\lambda y . f x y)) \rightarrow \nu(\lambda x . f x x) = \nu(\lambda x . \nu(\lambda y . f x y))$

using [[*smt-timeout = 120*]] **by** (*smt is-greatest-fixpoint-def is-greatest-postfixpoint-def isotone-def greatest-fixpoint-same greatest-postfixpoint greatest-postfixpoint-fixpoint*)

end

class *semiring* = *mult* + *one* + *ord* + *plus* + *zero* +
assumes *add-associative* : $(x + y) + z = x + (y + z)$
assumes *add-commutative* : $x + y = y + x$
assumes *add-idempotent* : $x + x = x$
assumes *add-left-zero* : $0 + x = x$
assumes *mult-associative* : $(x ; y) ; z = x ; (y ; z)$
assumes *mult-left-one* : $1 ; x = x$
assumes *mult-right-one* : $x ; 1 = x$
assumes *mult-left-dist-add* : $x ; (y + z) = x ; y + x ; z$
assumes *mult-right-dist-add* : $(x + y) ; z = x ; z + y ; z$
assumes *mult-left-zero* : $0 ; x = 0$
assumes *less-eq-def* : $x \leq y \leftrightarrow x + y = y$
assumes *less-def* : $x < y \leftrightarrow x \leq y \wedge \neg (y \leq x)$

begin

subclass *order*

by (*unfold-locales (metis less-def, metis add-idempotent less-eq-def, metis add-associative less-eq-def, metis add-commutative less-eq-def)*)

lemma *add-left-isotone*: $x \leq y \rightarrow x + z \leq y + z$

by (*smt add-associative add-commutative add-idempotent less-eq-def*)

lemma *add-right-isotone*: $x \leq y \rightarrow z + x \leq z + y$

by (*metis add-commutative add-left-isotone*)

lemma *add-isotone*: $w \leq y \wedge x \leq z \rightarrow w + x \leq y + z$

by (*smt add-associative add-commutative less-eq-def*)

lemma *add-left-upper-bound*: $x \leq x + y$

by (*metis add-associative add-idempotent less-eq-def*)

lemma *add-right-upper-bound*: $y \leq x + y$

by (*metis add-commutative add-left-upper-bound*)

lemma *add-least-upper-bound*: $x \leq z \wedge y \leq z \leftrightarrow x + y \leq z$

by (*smt add-associative add-commutative add-left-upper-bound less-eq-def*)

lemma *add-left-divisibility*: $x \leq y \leftrightarrow (\exists z . x + z = y)$

by (*metis add-left-upper-bound less-eq-def*)

lemma *add-right-divisibility*: $x \leq y \leftrightarrow (\exists z . z + x = y)$

by (*metis add-commutative add-left-divisibility*)

lemma *add-right-zero*: $x + 0 = x$

by (*metis add-commutative add-left-zero*)

lemma *zero-least*: $0 \leq x$

by (*metis add-left-upper-bound add-left-zero*)

lemma *mult-left-isotone*: $x \leq y \rightarrow x ; z \leq y ; z$

by (*metis less-eq-def mult-right-dist-add*)

lemma *mult-right-isotone*: $x \leq y \rightarrow z ; x \leq z ; y$

by (*metis less-eq-def mult-left-dist-add*)

lemma *mult-isotone*: $w \leq y \wedge x \leq z \rightarrow w ; x \leq y ; z$

by (*smt mult-left-isotone mult-right-isotone order-trans*)

lemma *mult-left-subdist-add-left*: $x ; y \leq x ; (y + z)$

by (*metis add-left-upper-bound mult-left-dist-add*)

lemma *mult-left-subdist-add-right*: $x ; z \leq x ; (y + z)$

by (*metis add-right-upper-bound mult-left-dist-add*)

lemma *mult-right-subdist-add-left*: $x ; z \leq (x + y) ; z$

by (*metis add-left-upper-bound mult-right-dist-add*)

lemma *mult-right-subdist-add-right*: $y ; z \leq (x + y) ; z$

by (*metis add-right-upper-bound mult-right-dist-add*)

lemma *case-split-left*: $1 \leq w + z \wedge w ; x \leq y \wedge z ; x \leq y \rightarrow x \leq y$

by (*smt add-least-upper-bound mult-left-isotone mult-left-one mult-right-dist-add order-trans*)

lemma *case-split-left-equal*: $w + z = 1 \wedge w ; x = w ; y \wedge z ; x = z ; y \rightarrow x = y$

by (*metis mult-left-one mult-right-dist-add*)

lemma *case-split-right*: $1 \leq w + z \wedge x ; w \leq y \wedge x ; z \leq y \rightarrow x \leq y$

by (*smt add-least-upper-bound mult-right-isotone mult-right-one mult-left-dist-add order-trans*)

lemma *case-split-right-equal*: $w + z = 1 \wedge x ; w = y ; w \wedge x ; z = y ; z \rightarrow x = y$

by (*metis mult-left-dist-add mult-right-one*)

lemma *zero-right-mult-decreasing*: $x ; 0 \leq x$

by (*metis add-right-zero mult-left-subdist-add-right mult-right-one*)

lemma *add-same-context*: $x \leq y + z \wedge y \leq x + z \rightarrow x + z = y + z$
by (*smt add-associative add-commutative less-eq-def*)

end

class *semiring-T* = *semiring* +
fixes *T* :: 'a (\top)
assumes *add-left-top*: $T + x = T$

begin

lemma *add-right-top*: $x + T = T$
by (*metis add-commutative add-left-top*)

lemma *top-greatest*: $x \leq T$
by (*metis add-left-top add-right-upper-bound*)

lemma *top-left-mult-increasing*: $x \leq T ; x$
by (*metis mult-left-isotone mult-left-one top-greatest*)

lemma *top-right-mult-increasing*: $x \leq x ; T$
by (*metis mult-right-isotone mult-right-one top-greatest*)

lemma *top-mult-top*: $T ; T = T$
by (*metis add-right-divisibility add-right-top top-right-mult-increasing*)

definition *vector* :: 'a \Rightarrow bool
where *vector* $x \leftrightarrow x = x ; T$

lemma *vector-zero*: *vector* 0
by (*metis mult-left-zero vector-def*)

lemma *vector-top*: *vector* T
by (*metis top-mult-top vector-def*)

lemma *vector-add-closed*: *vector* $x \wedge$ *vector* $y \rightarrow$ *vector* $(x + y)$
by (*metis mult-right-dist-add vector-def*)

lemma *vector-left-mult-closed*: *vector* $y \rightarrow$ *vector* $(x ; y)$
by (*metis mult-associative vector-def*)

end

class *itering-0* = *semiring* +
fixes *circ* :: 'a \Rightarrow 'a ($^{-\circ}$ [100] 100)

assumes *circ-mult*: $(x ; y)^\circ = 1 + x ; (y ; x)^\circ ; y$
assumes *circ-add*: $(x + y)^\circ = (x^\circ ; y)^\circ ; x^\circ$

begin

lemma *circ-isotone*: $x \leq y \rightarrow x^\circ \leq y^\circ$
by (*metis add-left-divisibility circ-add circ-mult mult-left-one mult-right-subdist-add-left*)

lemma *circ-slide*: $x ; (y ; x)^\circ = (x ; y)^\circ ; x$
by (*smt circ-mult mult-associative mult-left-dist-add mult-left-one mult-right-dist-add mult-right-one*)

lemma *circ-right-unfold*: $1 + x^\circ ; x = x^\circ$
by (*metis circ-mult mult-left-one mult-right-one*)

lemma *circ-left-unfold*: $1 + x ; x^\circ = x^\circ$
by (*metis circ-mult circ-slide mult-associative mult-right-one*)

lemma *circ-zero*: $0^\circ = 1$
by (*metis add-right-zero circ-left-unfold mult-left-zero*)

lemma *circ-increasing*: $x \leq x^\circ$
by (*metis add-right-upper-bound circ-right-unfold mult-left-one mult-right-subdist-add-left order-trans*)

lemma *circ-transitive-equal*: $x^\circ ; x^\circ = x^\circ$
by (*metis add-idempotent circ-add circ-mult mult-associative mult-left-one mult-right-one*)

lemma *circ-circ-circ*: $x^{\circ\circ} = x^{\circ\circ}$
by (*metis add-idempotent circ-add circ-increasing circ-transitive-equal less-eq-def*)

lemma *circ-one*: $1^\circ = 1^\circ$
by (*metis circ-circ-circ circ-zero*)

lemma *circ-add-1*: $(x + y)^\circ = x^\circ ; (y ; x)^\circ$
by (*metis circ-add circ-slide*)

lemma *circ-reflexive*: $1 \leq x^\circ$
by (*metis add-left-divisibility circ-right-unfold*)

lemma *circ-add-2*: $(x + y)^\circ \leq (x^\circ ; y)^\circ$
by (*metis add-least-upper-bound circ-increasing circ-isotone circ-reflexive mult-isotone mult-left-one mult-right-one*)

lemma *mult-zero-circ*: $(x ; 0)^\circ = 1 + x ; 0$
by (*metis circ-mult mult-associative mult-left-zero*)

lemma *mult-zero-add-circ*: $(x + y ; 0)^\circ = x^\circ ; (y ; 0)^\circ$
by (*metis circ-add-1 mult-associative mult-left-zero*)

lemma *mult-zero-add-circ-2*: $(x + y ; 0)^\circ = x^\circ + x^\circ ; y ; 0$
by (*metis mult-associative mult-left-dist-add mult-right-one mult-zero-add-circ mult-zero-circ*)

lemma *circ-plus-same*: $x^\circ ; x = x ; x^\circ$
by (*metis circ-slide mult-left-one mult-right-one*)

lemma *circ-plus-one*: $x^\circ = 1 + x^\circ$
by (*metis less-eq-def circ-reflexive*)

lemma *circ-rtc-2*: $1 + x + x^\circ ; x^\circ = x^\circ$
by (*metis add-associative circ-increasing circ-plus-one circ-transitive-equal less-eq-def*)

lemma *circ-unfold-sum*: $(x + y)^\circ = x^\circ + x^\circ ; y ; (x + y)^\circ$
by (*smt circ-add-1 circ-mult mult-associative mult-left-dist-add mult-right-one*)

lemma *circ-loop-fixpoint*: $y ; (y^\circ ; z) + z = y^\circ ; z$
by (*metis add-commutative circ-left-unfold mult-associative mult-left-one mult-right-dist-add*)

lemma *left-plus-below-circ*: $x ; x^\circ \leq x^\circ$
by (*metis add-right-upper-bound circ-left-unfold*)

lemma *right-plus-below-circ*: $x^\circ ; x \leq x^\circ$
by (*metis circ-plus-same left-plus-below-circ*)

lemma *circ-add-upper-bound*: $x \leq z^\circ \wedge y \leq z^\circ \rightarrow x + y \leq z^\circ$
by (*metis add-least-upper-bound*)

lemma *circ-mult-upper-bound*: $x \leq z^\circ \wedge y \leq z^\circ \rightarrow x ; y \leq z^\circ$
by (*metis mult-isotone circ-transitive-equal*)

lemma *circ-sub-dist*: $x^\circ \leq (x + y)^\circ$
by (*metis add-left-upper-bound circ-isotone*)

lemma *circ-sub-dist-1*: $x \leq (x + y)^\circ$
by (*metis add-least-upper-bound circ-increasing*)

lemma *circ-sub-dist-2*: $x ; y \leq (x + y)^\circ$
by (*metis add-commutative circ-mult-upper-bound circ-sub-dist-1*)

lemma *circ-sub-dist-3*: $x^\circ ; y^\circ \leq (x + y)^\circ$
by (*metis add-commutative circ-mult-upper-bound circ-sub-dist*)

lemma *circ-sup-one-left-unfold*: $1 \leq x \rightarrow x ; x^\circ = x^\circ$
by (*metis antisym less-eq-def mult-left-one mult-right-subdist-add-left left-plus-below-circ*)

lemma *circ-sup-one-right-unfold*: $1 \leq x \rightarrow x^\circ ; x = x^\circ$
by (*metis antisym less-eq-def mult-left-subdist-add-left mult-right-one right-plus-below-circ*)

lemma *circ-decompose-4*: $(x^\circ ; y^\circ)^\circ = x^\circ ; (y^\circ ; x^\circ)^\circ$
by (*smt circ-add circ-increasing circ-plus-one circ-slide circ-transitive-equal less-eq-def mult-associative mult-left-subdist-add-left mult-right-one*)

lemma *circ-decompose-5*: $(x^\circ ; y^\circ)^\circ = (y^\circ ; x^\circ)^\circ$
by (*metis add-associative add-commutative circ-add circ-decompose-4 circ-slide circ-zero mult-right-one*)

lemma *circ-decompose-6*: $x^\circ ; (y ; x^\circ)^\circ = y^\circ ; (x ; y^\circ)^\circ$
by (*metis add-commutative circ-add-1*)

lemma *circ-decompose-7*: $(x + y)^\circ = x^\circ ; y^\circ ; (x + y)^\circ$
by (*metis add-commutative circ-add circ-slide circ-transitive-equal mult-associative*)

lemma *circ-decompose-8*: $(x + y)^\circ = (x + y)^\circ ; x^\circ ; y^\circ$
by (*metis antisym eq-refl mult-associative mult-isotone mult-right-one circ-mult-upper-bound circ-reflexive circ-sub-dist-3*)

lemma *circ-decompose-9*: $(x^\circ ; y^\circ)^\circ = x^\circ ; y^\circ ; (x^\circ ; y^\circ)^\circ$
by (*metis circ-decompose-4 mult-associative*)

lemma *circ-decompose-10*: $(x^\circ ; y^\circ)^\circ = (x^\circ ; y^\circ)^\circ ; x^\circ ; y^\circ$
by (*metis circ-decompose-9 circ-plus-same mult-associative*)

lemma *circ-add-mult-zero*: $x^\circ ; y = (x + y ; 0)^\circ ; y$
by (*smt mult-associative mult-left-zero mult-right-one circ-add circ-slide circ-zero*)

lemma *circ-back-loop-fixpoint*: $(z ; y^\circ) ; y + z = z ; y^\circ$
by (*metis add-commutative mult-associative mult-left-dist-add mult-right-one circ-plus-same circ-left-unfold*)

lemma *circ-loop-is-fixpoint*: *is-fixpoint* $(\lambda x . y ; x + z) (y^\circ ; z)$
by (*metis circ-loop-fixpoint is-fixpoint-def*)

lemma *circ-back-loop-is-fixpoint*: *is-fixpoint* $(\lambda x . x ; y + z) (z ; y^\circ)$
by (*metis circ-back-loop-fixpoint is-fixpoint-def*)

lemma *circ-circ-add*: $(1 + x)^\circ = x^{\circ\circ}$
by (*metis add-commutative circ-add-1 circ-decompose-4 circ-zero mult-right-one*)

lemma *circ-circ-mult-sub*: $x^\circ ; 1^\circ \leq x^{\circ\circ}$
by (*metis circ-increasing circ-isotone circ-mult-upper-bound circ-reflexive*)

lemma *right-plus-circ*: $(x^\circ ; x)^\circ = x^\circ$
by (*metis add-idempotent circ-add circ-mult circ-right-unfold circ-slide*)

lemma *left-plus-circ*: $(x ; x^\circ)^\circ = x^\circ$
by (*metis circ-plus-same right-plus-circ*)

lemma *circ-add-sub-add-one*: $x ; x^\circ ; (x + y) \leq x ; x^\circ ; (1 + y)$

by (smt add-least-upper-bound add-left-upper-bound add-right-upper-bound circ-plus-same mult-associative mult-isotone mult-left-dist-add mult-right-one right-plus-below-circ)

lemma *circ-elimination*: $x ; y = 0 \rightarrow x ; y^\circ \leq x$

by (metis add-left-zero circ-back-loop-fixpoint circ-plus-same le-less mult-associative mult-left-zero)

lemma *circ-square*: $(x ; x)^\circ \leq x^\circ$

by (metis circ-increasing circ-isotone left-plus-circ mult-right-isotone)

lemma *circ-mult-sub-add*: $(x ; y)^\circ \leq (x + y)^\circ$

by (metis add-left-upper-bound add-right-upper-bound circ-isotone circ-square mult-isotone order-trans)

end

class *itering-1* = *itering-0* +

assumes *circ-simulate*: $z ; x \leq y ; z \rightarrow z ; x^\circ \leq y^\circ ; z$

begin

lemma *circ-circ-mult*: $1^\circ ; x^\circ = x^{\circ\circ}$

by (metis antisym circ-circ-add circ-reflexive circ-simulate circ-sub-dist-3 circ-sup-one-left-unfold circ-transitive-equal mult-left-one order-refl)

lemma *sub-mult-one-circ*: $x ; 1^\circ \leq 1^\circ ; x$

by (metis circ-simulate mult-left-one mult-right-one order-refl)

end

class *itering-2* = *itering-1* +

assumes *circ-simulate-right*: $z ; x \leq y ; z + w \rightarrow z ; x^\circ \leq y^\circ ; (z + w ; x^\circ)$

assumes *circ-simulate-left*: $x ; z \leq z ; y + w \rightarrow x^\circ ; z \leq (z + x^\circ ; w) ; y^\circ$

begin

lemma *circ-simulate-right-1*: $z ; x \leq y ; z \rightarrow z ; x^\circ \leq y^\circ ; z$

by (metis add-right-zero circ-simulate-right mult-left-zero)

lemma *circ-simulate-left-1*: $x ; z \leq z ; y \rightarrow x^\circ ; z \leq z ; y^\circ + x^\circ ; 0$

by (smt add-right-zero circ-simulate-left mult-associative mult-left-zero mult-right-dist-add)

lemma *circ-simulate-1*: $y ; x \leq x ; y \rightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$

by (smt add-associative add-right-zero circ-loop-fixpoint circ-simulate circ-simulate-left-1 mult-associative mult-left-zero mult-zero-add-circ-2)

lemma *circ-separate-1*: $y ; x \leq x ; y \rightarrow (x + y)^\circ = x^\circ ; y^\circ$

proof –

have $y ; x \leq x ; y \rightarrow y^\circ ; x ; y^\circ \leq x ; y^\circ + y^\circ ; 0$

by (smt circ-simulate-left-1 circ-transitive-equal mult-associative mult-left-isotone mult-left-zero mult-right-dist-add)

thus *?thesis*

by (smt add-commutative circ-add-1 circ-simulate-right circ-sub-dist-3 less-eq-def mult-associative mult-left-zero zero-right-mult-decreasing)

qed

lemma *circ-circ-mult-1*: $x^\circ ; 1^\circ = x^\circ$

by (*metis add-commutative circ-circ-add circ-separate-1 mult-left-one mult-right-one order-refl*)

lemma *atomicity-refinement*: $s = s ; q \wedge x = q ; x \wedge q ; b = 0 \wedge r ; b \leq b ; r \wedge r ; l \leq l ; r \wedge x ; l \leq l ; x \wedge b ; l \leq l ; b \wedge q ; l \leq l ; q \wedge r^\circ ; q \leq q ; r^\circ \wedge q \leq 1 \rightarrow s ; (x + b + r + l)^\circ ; q \leq s ; (x ; b^\circ ; q + r + l)^\circ$

proof

assume *I*: $s = s ; q \wedge x = q ; x \wedge q ; b = 0 \wedge r ; b \leq b ; r \wedge r ; l \leq l ; r \wedge x ; l \leq l ; x \wedge b ; l \leq l ; b \wedge q ; l \leq l ; q \wedge r^\circ ; q \leq q ; r^\circ \wedge q \leq 1$

hence $s ; (x + b + r + l)^\circ ; q = s ; l^\circ ; (x + b + r)^\circ ; q$

using [[*smt-timeout = 120*]] **by** (*smt add-commutative add-least-upper-bound circ-separate-1 mult-associative mult-left-subdist-add-right mult-right-dist-add order-trans*)

also have $\dots = s ; l^\circ ; b^\circ ; r^\circ ; q ; (x ; b^\circ ; r^\circ ; q)^\circ$ **using** *1*

by (*smt add-associative add-commutative circ-add-1 circ-separate-1 circ-slide mult-associative*)

also have $\dots \leq s ; l^\circ ; b^\circ ; r^\circ ; q ; (x ; b^\circ ; q ; r^\circ)^\circ$ **using** *1*

by (*metis circ-isotone mult-associative mult-right-isotone*)

also have $\dots \leq s ; q ; l^\circ ; b^\circ ; r^\circ ; (x ; b^\circ ; q ; r^\circ)^\circ$ **using** *1*

by (*metis mult-left-isotone mult-right-isotone mult-right-one*)

also have $\dots \leq s ; l^\circ ; q ; b^\circ ; r^\circ ; (x ; b^\circ ; q ; r^\circ)^\circ$ **using** *1*

by (*metis circ-simulate mult-associative mult-left-isotone mult-right-isotone*)

also have $\dots \leq s ; l^\circ ; r^\circ ; (x ; b^\circ ; q ; r^\circ)^\circ$ **using** *1*

by (*metis add-left-zero circ-back-loop-fixpoint circ-plus-same mult-associative mult-left-zero mult-left-isotone mult-right-isotone mult-right-one*)

also have $\dots \leq s ; (x ; b^\circ ; q + r + l)^\circ$ **using** *1*

by (*metis add-commutative circ-add-1 circ-sub-dist-3 mult-associative mult-right-isotone*)

finally show $s ; (x + b + r + l)^\circ ; q \leq s ; (x ; b^\circ ; q + r + l)^\circ$.

qed

end

class *itering-3 = itering-2 +*

assumes *circ-simulate-right-plus*: $z ; x \leq y ; y^\circ ; z + w \rightarrow z ; x^\circ \leq y^\circ ; (z + w ; x^\circ)$

assumes *circ-simulate-left-plus*: $x ; z \leq z ; y^\circ + w \rightarrow x^\circ ; z \leq (z + x^\circ ; w) ; y^\circ$

begin

lemma *circ-simulate-right-plus-1*: $z ; x \leq y ; y^\circ ; z \rightarrow z ; x^\circ \leq y^\circ ; z$

by (*metis add-right-zero circ-simulate-right-plus mult-left-zero*)

lemma *circ-simulate-left-plus-1*: $x ; z \leq z ; y^\circ \rightarrow x^\circ ; z \leq z ; y^\circ + x^\circ ; 0$

by (*smt add-right-zero circ-simulate-left-plus mult-associative mult-left-zero mult-right-dist-add*)

lemma *circ-simulate-2*: $y ; x^\circ \leq x^\circ ; y^\circ \leftrightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$

by (*rule iffI*,

smt add-associative add-right-zero circ-loop-fixpoint circ-simulate-left-plus-1 mult-associative mult-left-zero mult-zero-add-circ-2,

metis circ-increasing mult-left-isotone order-trans)

lemma *circ-simulate-absorb*: $y ; x \leq x \rightarrow y^\circ ; x \leq x + y^\circ ; 0$

by (*metis circ-simulate-left-plus-1 circ-zero mult-right-one*)

lemma *circ-simulate-3*: $y ; x^\circ \leq x^\circ \rightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$
by (*metis add-least-upper-bound circ-reflexive circ-simulate-2 less-eq-def mult-right-isotone mult-right-one*)

lemma *circ-square-2*: $(x ; x)^\circ ; (x + 1) \leq x^\circ$
by (*metis add-least-upper-bound circ-increasing circ-mult-upper-bound circ-reflexive circ-simulate-right-plus-1 mult-left-one mult-right-isotone mult-right-one*)

lemma *circ-separate-mult-1*: $y ; x \leq x ; y \rightarrow (x ; y)^\circ \leq x^\circ ; y^\circ$
by (*metis circ-mult-sub-add circ-separate-1*)

lemma *circ-extra-circ*: $(y ; x^\circ)^\circ = (y ; y^\circ ; x^\circ)^\circ$

proof –

have $y ; y^\circ ; x^\circ \leq y ; x^\circ ; (y ; x^\circ)^\circ$

by (*metis add-commutative circ-add-1 circ-sub-dist-3 mult-associative mult-right-isotone*)

hence $(y ; y^\circ ; x^\circ)^\circ \leq (y ; x^\circ)^\circ$

by (*metis circ-simulate-right-plus-1 mult-left-one mult-right-one*)

thus *?thesis*

by (*metis antisym circ-back-loop-fixpoint circ-isotone mult-right-subdist-add-right*)

qed

lemma *circ-separate-unfold*: $(y ; x^\circ)^\circ = y^\circ + y^\circ ; y ; x ; x^\circ ; (y ; x^\circ)^\circ$
by (*smt add-commutative circ-add circ-left-unfold circ-loop-fixpoint mult-associative mult-left-dist-add mult-right-one*)

lemma *separation*: $y ; x \leq x ; y^\circ \rightarrow (x + y)^\circ = x^\circ ; y^\circ$

proof –

have $y ; x \leq x ; y^\circ \rightarrow y^\circ ; x ; y^\circ \leq x ; y^\circ + y^\circ ; 0$

by (*smt circ-simulate-left-plus-1 circ-transitive-equal mult-associative mult-left-isotone mult-left-zero mult-right-dist-add*)

thus *?thesis*

by (*smt add-commutative circ-add-1 circ-simulate-right circ-sub-dist-3 less-eq-def mult-associative mult-left-zero zero-right-mult-decreasing*)

qed

lemma *circ-simulate-4*: $y ; x \leq x ; x^\circ ; (1 + y) \rightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$

proof –

have $y ; x \leq x ; x^\circ ; (1 + y) \rightarrow (1 + y) ; x \leq x ; x^\circ ; (1 + y)$

by (*smt add-associative add-commutative add-left-upper-bound circ-back-loop-fixpoint less-eq-def mult-left-dist-add mult-left-one mult-right-dist-add mult-right-one*)

thus *?thesis*

by (*smt add-least-upper-bound circ-increasing circ-reflexive circ-simulate-2 circ-simulate-right-plus-1 mult-right-isotone mult-right-subdist-add-right order-trans*)

qed

lemma *circ-simulate-5*: $y ; x \leq x ; x^\circ ; (x + y) \rightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$

by (*metis circ-add-sub-add-one circ-simulate-4 order-trans*)

lemma *circ-simulate-6*: $y ; x \leq x ; (x + y) \rightarrow y^\circ ; x^\circ \leq x^\circ ; y^\circ$

by (*metis add-commutative circ-back-loop-fixpoint circ-simulate-5 mult-right-subdist-add-left order-trans*)

lemma *circ-separate-4*: $y ; x \leq x ; x^\circ ; (1 + y) \rightarrow (x + y)^\circ = x^\circ ; y^\circ$

proof

assume 1: $y ; x \leq x ; x^\circ ; (1 + y)$
hence $y ; x ; x^\circ \leq x ; x^\circ + x ; x^\circ ; y ; x^\circ$
by (*smt circ-transitive-equal less-eq-def mult-associative mult-left-dist-add mult-right-dist-add mult-right-one*)
also have $\dots \leq x ; x^\circ + x ; x^\circ ; x^\circ ; y^\circ$ **using** 1
by (*metis add-right-isotone circ-simulate-2 circ-simulate-4 mult-associative mult-right-isotone*)
finally have $y ; x ; x^\circ \leq x ; x^\circ ; y^\circ$
by (*metis circ-reflexive circ-transitive-equal less-eq-def mult-associative mult-right-isotone mult-right-one*)
hence $y^\circ ; (y^\circ ; x)^\circ \leq x^\circ ; (y^\circ + y^\circ ; 0 ; (y^\circ ; x)^\circ)$
by (*smt add-right-upper-bound circ-back-loop-fixpoint circ-simulate-left-plus-1 circ-simulate-right-plus circ-transitive-equal mult-associative order-trans*)
thus $(x + y)^\circ = x^\circ ; y^\circ$
by (*smt add-commutative antisym circ-add-1 circ-slide circ-sub-dist-3 circ-transitive-equal less-eq-def mult-associative mult-left-zero mult-right-subdist-add-right zero-right-mult-decreasing*)
qed

lemma *circ-separate-5*: $y ; x \leq x ; x^\circ ; (x + y) \rightarrow (x + y)^\circ = x^\circ ; y^\circ$
by (*metis circ-add-sub-add-one circ-separate-4 order-trans*)

lemma *circ-separate-6*: $y ; x \leq x ; (x + y) \rightarrow (x + y)^\circ = x^\circ ; y^\circ$
by (*metis add-commutative circ-back-loop-fixpoint circ-separate-5 mult-right-subdist-add-left order-trans*)

end

class *itering-T* = *semiring-T* + *itering-3*

begin

lemma *circ-top*: $T^\circ = T$
by (*metis add-right-top antisym circ-left-unfold mult-left-subdist-add-left mult-right-one top-greatest*)

lemma *circ-right-top*: $x^\circ ; T = T$
by (*metis add-right-top circ-loop-fixpoint*)

lemma *circ-left-top*: $T ; x^\circ = T$
by (*metis add-right-top circ-add circ-right-top circ-top*)

lemma *mult-top-circ*: $(x ; T)^\circ = 1 + x ; T$
by (*metis circ-left-top circ-mult mult-associative*)

end

class *itering-L* = *itering-3* +
fixes $L :: 'a$
assumes *L-def*: $L = 1^\circ ; 0$

begin

lemma *one-circ-split*: $1^\circ = L + 1$

by (*metis L-def add-commutative antisym circ-add-upper-bound circ-reflexive circ-simulate-absorb mult-right-one order-refl zero-right-mult-decreasing*)

lemma *one-circ-mult-split*: $1^\circ ; x = L + x$

by (*metis L-def add-commutative circ-loop-fixpoint mult-associative mult-left-zero mult-zero-circ one-circ-split*)

lemma *one-circ-circ-split*: $1^{\circ\circ} = L + 1$

by (*metis circ-one one-circ-split*)

lemma *one-circ-mult-split-2*: $1^\circ ; x = x ; 1^\circ + L$

by (*smt L-def add-associative add-commutative circ-left-unfold less-eq-def mult-left-subdist-add-left mult-right-one one-circ-mult-split sub-mult-one-circ*)

lemma *sub-mult-one-circ-split*: $x ; 1^\circ \leq x + L$

by (*metis add-commutative one-circ-mult-split sub-mult-one-circ*)

lemma *sub-mult-one-circ-split-2*: $x ; 1^\circ \leq x + 1^\circ$

by (*metis L-def add-right-isotone order-trans sub-mult-one-circ-split zero-right-mult-decreasing*)

lemma *L-split*: $x ; L \leq x ; 0 + L$

by (*smt L-def mult-associative mult-left-isotone mult-right-dist-add sub-mult-one-circ-split-2*)

lemma *L-left-zero*: $L ; x = L$

by (*metis L-def mult-associative mult-left-zero*)

lemma *one-circ-L*: $1^\circ ; L = L$

by (*metis add-idempotent one-circ-mult-split*)

lemma *circ-circ-split*: $x^{\circ\circ} = L + x^\circ$

by (*metis circ-circ-mult one-circ-mult-split*)

lemma *circ-left-induct-mult-L*: $L \leq x \rightarrow x ; y \leq x \rightarrow x ; y^\circ \leq x$

by (*metis circ-one circ-simulate less-eq-def one-circ-mult-split*)

lemma *circ-left-induct-mult-iff-L*: $L \leq x \rightarrow x ; y \leq x \leftrightarrow x ; y^\circ \leq x$

by (*smt add-least-upper-bound circ-back-loop-fixpoint circ-left-induct-mult-L less-eq-def*)

lemma *circ-left-induct-L*: $L \leq x \rightarrow x ; y + z \leq x \rightarrow z ; y^\circ \leq x$

by (*metis add-least-upper-bound circ-left-induct-mult-L less-eq-def mult-right-dist-add*)

lemma *mult-L-circ*: $(x ; L)^\circ = 1 + x ; L$

by (*metis L-left-zero circ-mult mult-associative*)

lemma *mult-L-circ-mult*: $(x ; L)^\circ ; y = y + x ; L$

by (*metis L-left-zero mult-L-circ mult-associative mult-left-one mult-right-dist-add*)

lemma *circ-L*: $L^\circ = L + 1$

by (*metis L-left-zero add-commutative circ-left-unfold*)

lemma *L-below-one-circ*: $L \leq 1^\circ$

by (*metis add-left-upper-bound one-circ-split*)

lemma *circ-add-6*: $L + (x + y)^\circ = (x^\circ ; y^\circ)^\circ$

by (*metis add-associative add-commutative circ-add-1 circ-circ-add circ-circ-split circ-decompose-4*)

end

class *kleene-algebra* = *semiring* +

fixes *star* :: 'a \Rightarrow 'a (-* [100] 100)

assumes *star-left-unfold* : $1 + y ; y^* \leq y^*$

assumes *star-left-induct* : $z + y ; x \leq x \rightarrow y^* ; z \leq x$

assumes *star-right-induct*: $z + x ; y \leq x \rightarrow z ; y^* \leq x$

begin

lemma *star-left-unfold-equal*: $1 + x ; x^* = x^*$

by (*smt add-least-upper-bound antisym-conv mult-right-isotone mult-right-one order-refl order-trans star-left-induct star-left-unfold*)

lemma *star-unfold-slide*: $(x ; y)^* = 1 + x ; (y ; x)^* ; y$

by (*smt antisym mult-associative mult-left-dist-add mult-left-one mult-right-dist-add mult-right-isotone mult-right-one order-refl star-left-induct star-left-unfold-equal*)

lemma *star-decompose*: $(x + y)^* = (x^* ; y)^* ; x^*$

apply (*rule antisym*)

apply (*smt add-least-upper-bound add-left-upper-bound add-right-upper-bound mult-associative mult-left-dist-add mult-left-one mult-right-dist-add mult-right-one star-left-induct star-left-unfold-equal*)

by (*smt add-least-upper-bound less-eq-def mult-left-subdist-add-left mult-right-one star-left-induct star-left-unfold star-right-induct*)

lemma *star-simulation-right*: $z ; x \leq y ; z \rightarrow z ; x^* \leq y^* ; z$

by (*smt add-least-upper-bound less-eq-def mult-associative mult-left-dist-add mult-left-one mult-right-subdist-add-left star-left-induct star-left-unfold star-right-induct*)

end

sublocale *kleene-algebra* < *star!*: *itering-1* **where** *circ* = *star*

by (*unfold-locales, metis star-unfold-slide, metis star-decompose, metis star-simulation-right*)

context *kleene-algebra*

begin

Most lemmas in this class are taken from Georg Struth's theories.

lemma *star-sub-one*: $x \leq 1 \rightarrow x^* = 1$

by (*metis add-least-upper-bound eq-iff mult-left-one star.circ-reflexive star-right-induct*)

lemma *star-one*: $1^* = 1$

by (*metis eq-iff star-sub-one*)

lemma *star-left-induct-mult*: $x ; y \leq y \rightarrow x^* ; y \leq y$

by (*metis add-commutative less-eq-def order-refl star-left-induct*)

lemma *star-left-induct-mult-iff*: $x ; y \leq y \leftrightarrow x^* ; y \leq y$

by (*smt mult-associative mult-left-isotone mult-left-one mult-right-isotone order-trans star-left-induct-mult star.circ-reflexive star.left-plus-below-circ*)

lemma *star-involutive*: $x^* = x^{**}$

by (*smt antisym less-eq-def mult-left-subdist-add-left mult-right-one star-left-induct star.circ-plus-one star.left-plus-below-circ star.circ-transitive-equal*)

lemma *star-sup-one*: $(1 + x)^* = x^*$

by (*smt add-commutative less-eq-def mult-left-dist-add mult-left-one mult-right-dist-add mult-right-one star-involutive star.circ-plus-one star.circ-sub-dist star.left-plus-below-circ*)

lemma *star-left-induct-equal*: $z + x ; y = y \rightarrow x^* ; z \leq y$

by (*metis order-refl star-left-induct*)

lemma *star-left-induct-mult-equal*: $x ; y = y \rightarrow x^* ; y \leq y$

by (*metis order-refl star-left-induct-mult*)

lemma *star-star-upper-bound*: $x^* \leq z^* \rightarrow x^{**} \leq z^*$

by (*metis star-involutive*)

lemma *star-simulation-left*: $x ; z \leq z ; y \rightarrow x^* ; z \leq z ; y^*$

by (*smt add-commutative add-least-upper-bound mult-right-dist-add less-eq-def mult-associative mult-right-one star.left-plus-below-circ star.circ-increasing star-left-induct star-involutive star.circ-isotone star.circ-reflexive mult-left-subdist-add-left*)

lemma *quasicomm-1*: $y ; x \leq x ; (x + y)^* \leftrightarrow y^* ; x \leq x ; (x + y)^*$

by (*smt mult-isotone order-refl order-trans star.circ-increasing star-involutive star-simulation-left*)

lemma *star-rtc-3*: $1 + x + y ; y = y \rightarrow x^* \leq y$

by (*metis add-least-upper-bound less-eq-def mult-left-subdist-add-left mult-right-one star-left-induct-mult-iff star.circ-sub-dist*)

lemma *star-decompose-1*: $(x + y)^* = (x^* ; y^*)^*$

by (*smt add-least-upper-bound antisym mult-isotone mult-left-one mult-right-one star.circ-increasing star-involutive star.circ-isotone star.circ-reflexive star.circ-sub-dist-3*)

lemma *star-sum*: $(x + y)^* = (x^* + y^*)^*$

by (*metis star-decompose-1 star-involutive*)

lemma *star-decompose-3*: $(x^* ; y^*)^* = x^* ; (y ; x^*)^*$

by (*metis star-decompose-1 star.circ-add-1*)

lemma *star-right-induct-mult*: $y ; x \leq y \rightarrow y ; x^* \leq y$

by (*metis add-commutative less-eq-def order-refl star-right-induct*)

lemma *star-right-induct-mult-iff*: $y ; x \leq y \leftrightarrow y ; x^* \leq y$

by (*metis mult-right-isotone order-trans star.circ-increasing star-right-induct-mult*)

lemma *star-simulation-right-equal*: $z ; x = y ; z \rightarrow z ; x^* = y^* ; z$
by (*metis eq-iff star-simulation-left star-simulation-right*)

lemma *star-simulation-star*: $x ; y \leq y ; x \rightarrow x^* ; y^* \leq y^* ; x^*$
by (*metis star-simulation-left star-simulation-right*)

lemma *star-right-induct-equal*: $z + y ; x = y \rightarrow z ; x^* \leq y$
by (*metis order-refl star-right-induct*)

lemma *star-right-induct-mult-equal*: $y ; x = y \rightarrow y ; x^* \leq y$
by (*metis order-refl star-right-induct-mult*)

lemma *star-loop-least-fixpoint*: $y ; x + z = x \rightarrow y^* ; z \leq x$
by (*metis add-commutative star-left-induct-equal*)

lemma *star-back-loop-least-fixpoint*: $x ; y + z = x \rightarrow z ; y^* \leq x$
by (*metis add-commutative star-right-induct-equal*)

lemma *star-loop-is-least-fixpoint*: *is-least-fixpoint* $(\lambda x . y ; x + z)$ $(y^* ; z)$
by (*smt is-least-fixpoint-def star.circ-loop-fixpoint star-loop-least-fixpoint*)

lemma *star-loop-mu*: $\mu (\lambda x . y ; x + z) = y^* ; z$
by (*metis least-fixpoint-same star-loop-is-least-fixpoint*)

lemma *star-back-loop-is-least-fixpoint*: *is-least-fixpoint* $(\lambda x . x ; y + z)$ $(z ; y^*)$
by (*smt is-least-fixpoint-def star.circ-back-loop-fixpoint star-back-loop-least-fixpoint*)

lemma *star-back-loop-mu*: $\mu (\lambda x . x ; y + z) = z ; y^*$
by (*metis least-fixpoint-same star-back-loop-is-least-fixpoint*)

lemma *star-square*: $x^* = (1 + x) ; (x ; x)^*$

proof –

let $?f = \lambda y . y ; x + 1$

have 1: *isotone* $?f$

by (*smt add-left-isotone isotone-def mult-left-isotone*)

have 2: $?f \circ ?f = (\lambda y . y ; (x ; x) + (1 + x))$

by (*simp add: add-associative add-commutative mult-associative mult-left-one mult-right-dist-add o-def*)

thus *?thesis* using 1

by (*metis mu-square mult-left-one star-back-loop-mu has-least-fixpoint-def star-back-loop-is-least-fixpoint*)

qed

lemma *star-square-2*: $x^* = (x ; x)^* ; (x + 1)$

by (*smt add-commutative mult-left-dist-add mult-right-dist-add mult-right-one star.circ-decompose-5 star-involutive star-one star.circ-slide star-square*)

lemma *star-circ-simulate-right-plus*: $z ; x \leq y ; y^* ; z + w \rightarrow z ; x^* \leq y^* ; (z + w ; x^*)$

proof

assume $l: z ; x \leq y ; y^* ; z + w$
have $y^* ; (z + w ; x^*) ; x \leq y^* ; z ; x + y^* ; w ; x^*$
by (*smt add-left-upper-bound add-right-isotone mult-associative mult-left-dist-add mult-right-dist-add star.circ-back-loop-fixpoint*)
also have $\dots \leq y^* ; y ; z + y^* ; w + y^* ; w ; x^*$ **using** l
by (*smt add-left-isotone less-eq-def mult-associative mult-left-dist-add star.circ-plus-same star.circ-transitive-equal*)
also have $\dots \leq y^* ; (z + w ; x^*)$
by (*smt add-associative add-idempotent add-left-isotone mult-associative mult-left-dist-add mult-right-one mult-right-subdist-add-right star.circ-plus-one star.circ-plus-same star.circ-transitive-equal star.circ-unfold-sum*)
finally show $z ; x^* \leq y^* ; (z + w ; x^*)$
by (*smt add-least-upper-bound add-right-divisibility star.circ-loop-fixpoint star-right-induct*)
qed

lemma *star-circ-simulate-left-plus*: $x ; z \leq z ; y^* + w \rightarrow x^* ; z \leq (z + x^* ; w) ; y^*$

proof

assume $l: x ; z \leq z ; y^* + w$
have $x ; ((z + x^* ; w) ; y^*) \leq x ; z ; y^* + x^* ; w ; y^*$
by (*smt add-right-isotone mult-associative mult-left-dist-add mult-right-dist-add mult-right-subdist-add-left star.circ-loop-fixpoint*)
also have $\dots \leq (z + w + x^* ; w) ; y^*$ **using** l
by (*smt add-left-divisibility add-left-isotone mult-associative mult-right-dist-add star.circ-transitive-equal*)
also have $\dots = (z + x^* ; w) ; y^*$
by (*metis add-associative add-right-upper-bound less-eq-def star.circ-loop-fixpoint*)
finally show $x^* ; z \leq (z + x^* ; w) ; y^*$
by (*metis add-least-upper-bound mult-left-subdist-add-left mult-right-one star.circ-right-unfold star-left-induct*)

qed

end

sublocale *kleene-algebra* < *star!*: *itering-3* **where** *circ* = *star*

apply *unfold-locales*

apply (*smt add-associative add-commutative add-left-upper-bound mult-associative mult-left-dist-add order-trans star.circ-loop-fixpoint star-circ-simulate-right-plus*)

apply (*smt add-commutative add-right-isotone mult-right-isotone order-trans star.circ-increasing star-circ-simulate-left-plus*)

apply (*rule star-circ-simulate-right-plus*)

by (*rule star-circ-simulate-left-plus*)

class *kleene-itering* = *kleene-algebra* + *itering-3*

begin

lemma *star-below-circ*: $x^* \leq x^\circ$

by (*metis circ-right-unfold mult-left-one order-refl star-right-induct*)

lemma *star-zero-below-circ-mult*: $x^* ; 0 \leq x^\circ ; y$

by (*metis mult-isotone star-below-circ zero-least*)

lemma *star-mult-circ*: $x^* ; x^\circ = x^\circ$

by (*metis add-right-divisibility antisym circ-left-unfold star-left-induct-mult star.circ-loop-fixpoint*)

lemma *circ-mult-star*: $x^\circ ; x^* = x^\circ$
by (*metis circ-slide mult-left-one mult-right-one star-mult-circ star-simulation-right-equal*)

lemma *circ-star*: $x^{\circ*} = x^\circ$
by (*metis circ-reflexive circ-transitive-equal less-eq-def mult-left-one star-one star-simulation-right-equal star-square*)

lemma *star-circ*: $x^{*\circ} = x^{\circ\circ}$
by (*metis antisym circ-circ-add circ-sub-dist less-eq-def star.circ-rtc-2 star-below-circ*)

lemma *circ-add-3*: $(x^\circ ; y^\circ)^* \leq (x + y)^\circ$
by (*metis circ-add-1 circ-isotone circ-left-unfold circ-star mult-left-subdist-add-left mult-right-isotone mult-right-one star.circ-isotone*)

lemma *circ-isolate*: $x^\circ = x^\circ ; 0 + x^*$
by (*metis add-commutative antisym circ-add-upper-bound circ-mult-star circ-simulate-absorb star.left-plus-below-circ star-below-circ zero-right-mult-decreasing*)

lemma *circ-isolate-mult*: $x^\circ ; y = x^\circ ; 0 + x^* ; y$
by (*metis circ-isolate mult-associative mult-left-zero mult-right-dist-add*)

lemma *circ-isolate-mult-sub*: $x^\circ ; y \leq x^\circ + x^* ; y$
by (*metis add-left-isotone circ-isolate-mult zero-right-mult-decreasing*)

lemma *circ-sub-decompose*: $(x^\circ ; y)^\circ \leq (x^* ; y)^\circ ; x^\circ$
by (*smt add-commutative add-least-upper-bound add-right-upper-bound circ-back-loop-fixpoint circ-isolate-mult mult-zero-add-circ-2 zero-right-mult-decreasing*)

lemma *circ-add-4*: $(x + y)^\circ = (x^* ; y)^\circ ; x^\circ$
by (*smt antisym circ-add circ-isotone circ-sub-decompose circ-transitive-equal mult-associative mult-left-isotone star-below-circ*)

lemma *circ-add-5*: $(x^\circ ; y)^\circ ; x^\circ = (x^* ; y)^\circ ; x^\circ$
by (*metis circ-add circ-add-4*)

lemma *plus-circ*: $(x^* ; x)^\circ = x^\circ$
by (*smt add-idempotent circ-add-4 circ-decompose-7 circ-star star.circ-decompose-5 star.right-plus-circ*)

end

class *kleene-algebra-T* = *semiring-T* + *kleene-algebra*

sublocale *kleene-algebra-T* < *star!*: *itering-T* **where** *circ* = *star* ..

class *omega-algebra-T* = *kleene-algebra-T* +
fixes *omega* :: 'a \Rightarrow 'a ($-\omega$ [100] 100)
assumes *omega-unfold*: $y^\omega = y ; y^\omega$
assumes *omega-induct*: $x \leq z + y ; x \rightarrow x \leq y^\omega + y^* ; z$

begin

Most lemmas in this class are taken from Georg Struth's theories.

lemma *star-zero-below-omega*: $x^* ; 0 \leq x^\omega$

by (*metis add-left-zero omega-unfold star-left-induct-equal*)

lemma *star-zero-below-omega-zero*: $x^* ; 0 \leq x^\omega ; 0$

by (*metis add-left-zero mult-associative omega-unfold star-left-induct-equal*)

lemma *omega-induct-mult*: $y \leq x ; y \rightarrow y \leq x^\omega$

by (*metis add-commutative add-left-zero less-eq-def omega-induct star-zero-below-omega*)

lemma *omega-sub-dist*: $x^\omega \leq (x+y)^\omega$

by (*metis mult-right-subdist-add-left omega-induct-mult omega-unfold*)

lemma *omega-isotone*: $x \leq y \rightarrow x^\omega \leq y^\omega$

by (*metis less-eq-def omega-sub-dist*)

lemma *omega-induct-equal*: $y = z + x ; y \rightarrow y \leq x^\omega + x^* ; z$

by (*metis omega-induct order-refl*)

lemma *omega-zero*: $0^\omega = 0$

by (*metis mult-left-zero omega-unfold*)

lemma *omega-one-greatest*: $x \leq 1^\omega$

by (*metis mult-left-one omega-induct-mult order-refl*)

lemma *omega-one*: $1^\omega = T$

by (*metis add-left-top less-eq-def omega-one-greatest*)

lemma *star-mult-omega*: $x^* ; x^\omega = x^\omega$

by (*metis antisym-conv mult-isotone omega-unfold star.circ-increasing star-left-induct-mult-equal star-left-induct-mult-iff*)

lemma *star-omega-top*: $x^{*\omega} = T$

by (*metis add-left-top less-eq-def omega-one omega-sub-dist star.circ-plus-one*)

lemma *omega-sub-vector*: $x^\omega ; y \leq x^\omega$

by (*metis mult-associative omega-induct-mult omega-unfold order-refl*)

lemma *omega-vector*: $x^\omega ; T = x^\omega$

by (*metis add-commutative less-eq-def omega-sub-vector top-right-mult-increasing*)

lemma *omega-simulation*: $z ; x \leq y ; z \rightarrow z ; x^\omega \leq y^\omega$

by (*smt less-eq-def mult-associative mult-right-subdist-add-left omega-induct-mult omega-unfold*)

lemma *omega-omega*: $x^{\omega\omega} \leq x^\omega$

by (*metis omega-sub-vector omega-unfold*)

lemma *left-plus-omega*: $(x ; x^*)^\omega = x^\omega$

by (*metis antisym mult-associative mult-right-isotone mult-right-one omega-induct-mult omega-unfold star-mult-omega star-one star.circ-reflexive star.circ-right-unfold star.circ-plus-same star-sum star-sup-one*)

lemma *omega-slide*: $x ; (y ; x)^\omega = (x ; y)^\omega$

by (*metis antisym mult-associative mult-right-isotone omega-simulation omega-unfold order-refl*)

lemma *omega-simulation-2*: $y ; x \leq x ; y \rightarrow (x ; y)^\omega \leq x^\omega$

by (*metis less-eq-def mult-right-isotone omega-induct-mult omega-slide omega-sub-dist*)

lemma *wagner*: $(x + y)^\omega = x ; (x + y)^\omega + z \rightarrow (x + y)^\omega = x^\omega + x^* ; z$

by (*metis add-commutative add-least-upper-bound eq-iff omega-induct omega-sub-dist star-left-induct*)

lemma *right-plus-omega*: $(x^* ; x)^\omega = x^\omega$

by (*metis left-plus-omega star.circ-plus-same*)

lemma *omega-sub-dist-1*: $(x ; y^*)^\omega \leq (x + y)^\omega$

by (*metis add-least-upper-bound left-plus-omega mult-isotone mult-left-one mult-right-dist-add omega-isotone order-refl star-decompose-1 star.circ-increasing star.circ-plus-one*)

lemma *omega-sub-dist-2*: $(x^* ; y)^\omega \leq (x + y)^\omega$

by (*metis add-commutative mult-isotone omega-slide omega-sub-dist-1 star-mult-omega star.circ-sub-dist*)

lemma *omega-star*: $(x^\omega)^* = 1 + x^\omega$

by (*metis mult-associative omega-unfold omega-vector star.circ-plus-same star-left-unfold-equal star-omega-top*)

lemma *omega-mult-omega-star*: $x^\omega ; x^{\omega*} = x^\omega$

by (*metis mult-associative omega-unfold omega-vector star.circ-plus-same star-omega-top*)

lemma *omega-sum-unfold-1*: $(x + y)^\omega = x^\omega + x^* ; y ; (x + y)^\omega$

by (*metis mult-associative mult-right-dist-add omega-unfold wagner*)

lemma *omega-sum-unfold-2*: $(x + y)^\omega \leq (x^* ; y)^\omega + (x^* ; y)^* ; x^\omega$

by (*metis omega-induct-equal omega-sum-unfold-1*)

lemma *omega-sum-unfold-3*: $(x^* ; y)^* ; x^\omega \leq (x + y)^\omega$

by (*metis omega-sum-unfold-1 star-left-induct-equal*)

lemma *omega-decompose*: $(x + y)^\omega = (x^* ; y)^\omega + (x^* ; y)^* ; x^\omega$

by (*metis add-least-upper-bound antisym omega-sub-dist-2 omega-sum-unfold-2 omega-sum-unfold-3*)

lemma *omega-loop-fixpoint*: $y ; (y^\omega + y^* ; z) + z = y^\omega + y^* ; z$

by (*metis add-associative mult-left-dist-add omega-unfold star.circ-loop-fixpoint*)

lemma *omega-loop-greatest-fixpoint*: $y ; x + z = x \rightarrow x \leq y^\omega + y^* ; z$

by (*metis add-commutative omega-induct-equal*)

lemma *omega-square*: $x^\omega = (x ; x)^\omega$

by (smt antisym mult-associative omega-induct-mult omega-slide omega-sub-vector omega-unfold omega-vector top-mult-top)

lemma *mult-top-omega*: $(x ; T)^\omega \leq x ; T$

by (metis mult-right-isotone omega-slide top-greatest)

lemma *mult-zero-omega*: $(x ; 0)^\omega = x ; 0$

by (metis mult-left-zero omega-slide)

lemma *mult-zero-add-omega*: $(x + y ; 0)^\omega = x^\omega + x^* ; y ; 0$

by (smt add-associative add-commutative add-idempotent mult-associative mult-left-one mult-left-zero mult-right-dist-add mult-zero-omega star.mult-zero-circ omega-decompose)

lemma *omega-mult-star*: $x^\omega ; x^* = x^\omega$

by (metis mult-associative omega-vector star.circ-left-top)

lemma *omega-loop-is-greatest-fixpoint*: *is-greatest-fixpoint* $(\lambda x . y ; x + z) (y^\omega + y^* ; z)$

by (smt is-greatest-fixpoint-def omega-loop-fixpoint omega-loop-greatest-fixpoint)

lemma *omega-loop-nu*: $\nu (\lambda x . y ; x + z) = y^\omega + y^* ; z$

by (metis greatest-fixpoint-same omega-loop-is-greatest-fixpoint)

lemma *omega-loop-zero-is-greatest-fixpoint*: *is-greatest-fixpoint* $(\lambda x . y ; x) (y^\omega)$

by (metis is-greatest-fixpoint-def omega-induct-mult omega-unfold order-refl)

lemma *omega-loop-zero-nu*: $\nu (\lambda x . y ; x) = y^\omega$

by (metis greatest-fixpoint-same omega-loop-zero-is-greatest-fixpoint)

lemma *omega-separate-unfold*: $(x^* ; y)^\omega = y^\omega + y^* ; x ; (x^* ; y)^\omega$

by (metis add-commutative mult-associative omega-slide omega-sum-unfold-1 star.circ-loop-fixpoint)

lemma *omega-circ-simulate-right-plus*: $z ; x \leq y ; (y^\omega ; 0 + y^*) ; z + w \rightarrow z ; (x^\omega ; 0 + x^*) \leq (y^\omega ; 0 + y^*) ; (z + w ; (x^\omega ; 0 + x^*))$

proof

assume $z ; x \leq y ; (y^\omega ; 0 + y^*) ; z + w$

hence 1: $z ; x \leq y^\omega ; 0 + y ; y^* ; z + w$

by (metis mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add omega-unfold)

hence $(y^\omega ; 0 + y^* ; z + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*) ; x \leq y^\omega ; 0 + y^* ; (y^\omega ; 0 + y ; y^* ; z + w) + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*$

by (smt add-associative add-left-upper-bound add-right-upper-bound less-eq-def mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add star.circ-back-loop-fixpoint)

also have $\dots = y^\omega ; 0 + y^* ; y ; y^* ; z + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*$

by (smt add-associative add-right-upper-bound less-eq-def mult-associative mult-left-dist-add star.circ-back-loop-fixpoint star-mult-omega)

finally have $z + (y^\omega ; 0 + y^* ; z + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*) ; x \leq y^\omega ; 0 + y^* ; z + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*$

by (smt add-least-upper-bound add-left-isotone add-left-upper-bound mult-left-isotone order-trans star.circ-loop-fixpoint star.circ-plus-same star.circ-transitive-equal star.left-plus-below-circ)

hence 2: $z ; x^* \leq y^\omega ; 0 + y^* ; z + y^* ; w ; x^\omega ; 0 + y^* ; w ; x^*$

by (metis star-right-induct)

have $z ; x^\omega ; 0 \leq (y^\omega ; 0 + y ; y^* ; z + w) ; x^\omega ; 0$ **using** 1

by (smt add-left-divisibility mult-associative mult-right-subdist-add-left omega-unfold)

hence $z ; x^\omega ; 0 \leq y^\omega + y^* ; (y^\omega ; 0 + w ; x^\omega ; 0)$

by (smt add-associative add-commutative left-plus-omega mult-associative mult-left-zero mult-right-dist-add omega-induct star.left-plus-circ)

thus $z ; (x^\omega ; 0 + x^*) \leq (y^\omega ; 0 + y^*) ; (z + w ; (x^\omega ; 0 + x^*))$ **using** 2

by (*smt add-associative add-commutative less-eq-def mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add omega-unfold omega-zero star-mult-omega zero-right-mult-decreasing*)
qed

lemma *omega-circ-simulate-left-plus*: $x ; z \leq z ; (y^\omega ; 0 + y^*) + w \rightarrow (x^\omega ; 0 + x^*) ; z \leq (z + (x^\omega ; 0 + x^*) ; w) ; (y^\omega ; 0 + y^*)$

proof

assume 1: $x ; z \leq z ; (y^\omega ; 0 + y^*) + w$

have $x ; (z ; y^\omega ; 0 + z ; y^* + x^\omega ; 0 + x^* ; w ; y^\omega ; 0 + x^* ; w ; y^*) \leq x ; z ; y^\omega ; 0 + x ; z ; y^* + x^\omega ; 0 + x^* ; w ; y^\omega ; 0 + x^* ; w ; y^*$

by (*smt add-associative add-right-upper-bound less-eq-def mult-associative mult-left-dist-add omega-unfold star.circ-loop-fixpoint*)

also have $\dots \leq (z ; y^\omega ; 0 + z ; y^* + w) ; y^\omega ; 0 + (z ; y^\omega ; 0 + z ; y^* + w) ; y^* + x^\omega ; 0 + x^* ; w ; y^\omega ; 0 + x^* ; w ; y^*$ **using** 1

by (*metis add-left-isotone mult-associative mult-left-dist-add mult-left-isotone*)

also have $\dots = z ; y^\omega ; 0 + z ; y^* + x^\omega ; 0 + x^* ; w ; y^\omega ; 0 + x^* ; w ; y^*$

by (*smt add-associative add-commutative add-idempotent mult-associative mult-left-zero mult-right-dist-add star.circ-loop-fixpoint star.circ-transitive-equal star-mult-omega*)

finally have $(x^\omega ; 0 + x^*) ; z \leq z ; y^\omega ; 0 + z ; y^* + x^\omega ; 0 + x^* ; w ; y^\omega ; 0 + x^* ; w ; y^*$

by (*smt add-least-upper-bound add-left-upper-bound mult-associative mult-left-zero mult-right-dist-add star.circ-back-loop-fixpoint star-left-induct*)

thus $(x^\omega ; 0 + x^*) ; z \leq (z + (x^\omega ; 0 + x^*) ; w) ; (y^\omega ; 0 + y^*)$

by (*smt add-associative mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add*)

qed

end

sublocale *omega-algebra-T* < *comb0!*: *itering-T* **where** *circ* = $(\lambda x . x^\omega ; 0 + x^*)$

apply *unfold-locales*

apply (*smt add-associative add-commutative mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add omega-slide star.circ-mult*)

apply (*smt add-associative add-commutative star.circ-add mult-associative mult-left-dist-add mult-left-zero mult-right-dist-add star.mult-zero-add-circ-2 mult-zero-add-omega omega-decompose*)

apply (*smt add-isotone star.circ-simulate mult-associative mult-left-dist-add mult-left-isotone mult-left-zero mult-right-dist-add omega-simulation*)

apply (*smt add-commutative add-left-isotone mult-left-isotone mult-left-subdist-add-left mult-right-one omega-circ-simulate-right-plus order-trans star.circ-plus-one*)

apply (*smt add-commutative add-right-isotone mult-left-subdist-add-left mult-right-isotone omega-circ-simulate-left-plus order-trans star.circ-increasing*)

apply (*metis omega-circ-simulate-right-plus*)

by (*metis omega-circ-simulate-left-plus*)

class *omega-itering* = *omega-algebra-T* + *itering-T*

begin

subclass *kleene-itering* ..

lemma *circ-below-omega-star*: $x^\circ \leq x^\omega + x^*$

by (*metis circ-left-unfold mult-right-one omega-induct order-refl*)

lemma *omega-mult-circ*: $x^\omega ; x^\circ = x^\omega$

by (*smt add-commutative circ-reflexive less-eq-def mult-left-dist-add mult-right-one omega-sub-vector*)

lemma *circ-mult-omega*: $x^\circ ; x^\omega = x^\omega$

by (*metis antisym circ-left-unfold circ-slide le-less mult-left-one mult-right-one mult-right-subdist-add-left omega-simulation*)

lemma *circ-omega*: $x^{\circ\omega} = T$

by (*metis circ-star star-omega-top*)

lemma *omega-circ*: $x^{\omega^\circ} = 1 + x^\omega$

by (*metis circ-left-unfold circ-star mult-associative omega-vector star.circ-left-top*)

end

class *Omega* =

fixes *Omega* :: 'a \Rightarrow 'a ($^{-\Omega}$ [100] 100)

class *dra* = *kleene-algebra-T* + *Omega* +

assumes *Omega-unfold* : $y^\Omega = 1 + y$; y^Ω

assumes *Omega-isolate* : $y^\Omega = y^\Omega$; $0 + y^*$

assumes *Omega-induct* : $x \leq z + y$; $x \rightarrow x \leq y^\Omega$; z

begin

lemma *Omega-mult*: $(x ; y)^\Omega = 1 + x$; $(y ; x)^\Omega$; y

by (*smt Omega-induct Omega-unfold antisym mult-associative mult-left-dist-add mult-left-one mult-right-dist-add mult-right-isotone mult-right-one order-refl*)

lemma *Omega-add-1*: $(x + y)^\Omega = x^\Omega$; $(y ; x^\Omega)^\Omega$

by (*smt Omega-induct Omega-unfold add-associative add-commutative add-right-isotone antisym mult-associative mult-left-one mult-right-dist-add mult-right-isotone mult-right-one order-refl*)

lemma *Omega-add*: $(x + y)^\Omega = (x^\Omega ; y)^\Omega$; x^Ω

by (*smt Omega-add-1 Omega-mult mult-associative mult-left-dist-add mult-left-one mult-right-dist-add mult-right-one*)

lemma *Omega-simulate*: $z ; x \leq y$; $z \rightarrow z$; $x^\Omega \leq y^\Omega$; z

by (*smt Omega-induct Omega-unfold add-right-isotone mult-associative mult-left-dist-add mult-left-isotone mult-right-one*)

end

sublocale *dra* < *Omega*!: *itering-1* **where** *circ* = *Omega*

apply *unfold-locales*

apply (*metis Omega-mult*)

apply (*metis Omega-add*)

by (*metis Omega-simulate*)

context *dra*

begin

lemma *star-below-Omega*: $x^* \leq x^\Omega$

by (*metis Omega-isolate add-commutative add-left-divisibility*)

lemma *Omega-sum-unfold-1*: $(x + y)^\Omega = y^\Omega + y^*$; x ; $(x + y)^\Omega$

by (*smt Omega.circ-add Omega.circ-loop-fixpoint Omega-isolate add-associative add-commutative mult-associative mult-left-zero mult-right-dist-add*)

lemma *Omega-add-3*: $(x + y)^\Omega = (x^* ; y)^\Omega ; x^\Omega$

by (*smt Omega.circ-add Omega.circ-isotone Omega-induct Omega-sum-unfold-1 add-commutative antisym mult-left-isotone order-refl star-below-Omega*)

lemma *Omega-one*: $1^\Omega = T$

by (*metis Omega.circ-transitive-equal Omega-induct add-left-top add-right-upper-bound less-eq-def mult-left-one*)

lemma *top-left-zero*: $T ; x = T$

by (*metis Omega-induct Omega-one add-left-top add-right-upper-bound less-eq-def mult-left-one*)

lemma *star-mult-Omega*: $x^\Omega = x^* ; x^\Omega$

by (*metis Omega.circ-plus-one Omega.circ-transitive-equal Omega-isolate add-commutative less-eq-def order-refl star.circ-add-mult-zero star.circ-increasing star.circ-plus-same star-involutive star-left-induct star-square*)

lemma *Omega-separate-2*: $y ; x \leq x ; (x + y) \rightarrow (x + y)^\Omega = x^\Omega ; y^\Omega$

by (*smt Omega.circ-sub-dist-3 Omega-induct Omega-sum-unfold-1 add-right-isotone antisym mult-associative mult-left-isotone star-mult-Omega star-simulation-left*)

lemma *Omega-circ-simulate-right-plus*: $z ; x \leq y ; y^\Omega ; z + w \rightarrow z ; x^\Omega \leq y^\Omega ; (z + w ; x^\Omega)$

by (*smt Omega.circ-back-loop-fixpoint Omega.circ-plus-same Omega.left-plus-circ Omega-induct add-associative add-commutative add-right-isotone less-eq-def mult-associative mult-right-dist-add*)

lemma *Omega-circ-simulate-left-plus*: $x ; z \leq z ; y^\Omega + w \rightarrow x^\Omega ; z \leq (z + x^\Omega ; w) ; y^\Omega$

proof

assume 1: $x ; z \leq z ; y^\Omega + w$

hence $x ; ((z + x^\Omega ; w) ; y^\Omega) \leq z ; y^\Omega ; y^\Omega + w ; y^\Omega + x^\Omega ; w ; y^\Omega$

by (*smt Omega.left-plus-below-circ add-left-isotone add-right-isotone mult-associative mult-left-dist-add mult-left-isotone mult-right-dist-add order-trans*)

hence $x ; ((z + x^\Omega ; w) ; y^\Omega) \leq (z + x^\Omega ; w) ; y^\Omega$

by (*metis Omega.circ-transitive-equal mult-associative Omega.circ-reflexive add-associative less-eq-def mult-left-one mult-right-dist-add*)

thus $x^\Omega ; z \leq (z + x^\Omega ; w) ; y^\Omega$

by (*smt Omega.circ-back-loop-fixpoint Omega-isolate add-least-upper-bound mult-associative mult-left-zero mult-right-dist-add mult-right-subdist-add-left mult-right-subdist-add-right star-left-induct*)

qed

end

sublocale *dra* < *Omega!*: *itering-T* **where** *circ* = *Omega*

apply *unfold-locales*

apply (*smt Omega.circ-loop-fixpoint Omega-circ-simulate-right-plus add-associative add-commutative add-left-upper-bound mult-associative mult-left-dist-add order-trans*)

apply (*smt Omega.circ-increasing Omega-circ-simulate-left-plus add-commutative add-right-isotone mult-right-isotone order-trans*)

apply (*metis Omega-circ-simulate-right-plus*)

by (*metis Omega-circ-simulate-left-plus*)

class *dra-omega* = *omega-algebra-T* + *Omega* +

assumes *top-left-zero*: $T ; x = T$

assumes *Omega-def*: $x^\Omega = x^\omega + x^*$

begin

lemma *omega-left-zero*: $x^\omega ; y = x^\omega$

by (*metis mult-associative omega-vector top-left-zero*)

lemma *Omega-mult*: $(x ; y)^\Omega = 1 + x ; (y ; x)^\Omega ; y$

by (*smt add-associative add-commutative Omega-def mult-left-dist-add mult-right-dist-add omega-left-zero omega-slide star-unfold-slide*)

lemma *Omega-add*: $(x + y)^\Omega = (x^\Omega ; y)^\Omega ; x^\Omega$

proof –

have $(x^\Omega ; y)^\Omega ; x^\Omega = (x^* ; y)^* ; x^\omega + (x^* ; y)^\omega + (x^* ; y)^* ; x^{\omega*} ; x^\Omega$

by (*smt add-commutative Omega-def mult-associative mult-right-dist-add mult-zero-add-omega omega-left-zero star.circ-add-1*)

thus *?thesis*

by (*smt add-associative add-commutative Omega-def mult-associative mult-left-dist-add omega-decompose omega-left-zero star.circ-add-1 star.circ-loop-fixpoint star.circ-slide*)

qed

lemma *Omega-simulate*: $z ; x \leq y ; z \rightarrow z ; x^\Omega \leq y^\Omega ; z$

by (*smt add-isotone Omega-def mult-left-dist-add mult-right-dist-add omega-left-zero omega-simulation star-simulation-right*)

subclass *dra*

apply *unfold-locales*

apply (*metis Omega-def comb0.circ-left-unfold omega-left-zero*)

apply (*smt Omega-def add-associative less-eq-def mult-right-dist-add omega-left-zero zero-right-mult-decreasing*)

by (*metis Omega-def mult-right-dist-add omega-induct omega-left-zero*)

end

class *dom* =

fixes $d :: 'a \Rightarrow 'a$

class *domain-semiring* = *semiring* + *dom* +

assumes *d-restrict*: $x + d(x) ; x = d(x) ; x$

assumes *d-mult-d* : $d(x ; y) = d(x ; d(y))$

assumes *d-plus-one*: $d(x) + 1 = 1$

assumes *d-zero* : $d(0) = 0$

assumes *d-dist-add*: $d(x + y) = d(x) + d(y)$

begin

Most lemmas in this class are taken from Georg Struth's theories.

lemma *d-restrict-equals*: $x = d(x) ; x$

by (*metis add-commutative d-plus-one d-restrict mult-left-one mult-right-dist-add*)

lemma *d-involutive*: $d(d(x)) = d(x)$

by (*metis d-mult-d mult-left-one*)

lemma *d-fixpoint*: $(\exists y . x = d(y)) \leftrightarrow x = d(x)$

by (*metis d-involutive*)

lemma *d-type*: $\forall P . (\forall x . x = d(x) \rightarrow P(x)) \leftrightarrow (\forall x . P(d(x)))$

by (*metis d-involutive*)

lemma *d-mult-sub*: $d(x ; y) \leq d(x)$

by (*metis d-dist-add d-mult-d d-plus-one less-eq-def mult-left-subdist-add-left mult-right-one*)

lemma *d-sub-one*: $x \leq 1 \rightarrow x \leq d(x)$

by (*metis d-restrict-equals mult-right-isotone mult-right-one*)

lemma *d-strict*: $d(x) = 0 \leftrightarrow x = 0$

by (*metis d-restrict-equals d-zero mult-left-zero*)

lemma *d-one*: $d(1) = 1$

by (*metis d-restrict-equals mult-right-one*)

lemma *d-below-one*: $d(x) \leq 1$

by (*metis d-plus-one less-eq-def*)

lemma *d-isotone*: $x \leq y \rightarrow d(x) \leq d(y)$

by (*metis d-dist-add less-eq-def*)

lemma *d-plus-left-upper-bound*: $d(x) \leq d(x + y)$

by (*metis add-left-upper-bound d-isotone*)

lemma *d-export*: $d(d(x) ; y) = d(x) ; d(y)$

by (*smt add-commutative antisym d-mult-d d-mult-sub d-plus-left-upper-bound d-plus-one d-restrict d-sub-one mult-isotone mult-left-one mult-left-subdist-add-right mult-right-dist-add mult-right-one*)

lemma *d-idempotent*: $d(x) ; d(x) = d(x)$

by (*metis d-export d-restrict-equals*)

lemma *d-commutative*: $d(x) ; d(y) = d(y) ; d(x)$

by (*smt antisym d-export d-mult-d d-mult-sub d-one d-restrict-equals mult-isotone mult-left-one*)

lemma *d-least-left-preserved*: $x \leq d(y) ; x \leftrightarrow d(x) \leq d(y)$

by (*metis d-below-one d-involutive d-mult-sub d-restrict-equals eq-iff mult-left-isotone mult-left-one*)

lemma *d-weak-locality*: $x ; y = 0 \leftrightarrow x ; d(y) = 0$

by (*metis d-mult-d d-strict*)

lemma *d-add-closed*: $d(d(x) + d(y)) = d(x) + d(y)$

by (*metis d-dist-add d-involutive*)

lemma *d-mult-closed*: $d(d(x) ; d(y)) = d(x) ; d(y)$

by (*metis d-export d-mult-d*)

lemma *d-mult-left-lower-bound*: $d(x) ; d(y) \leq d(x)$
by (*metis d-export d-involutive d-mult-sub*)

lemma *d-mult-greatest-lower-bound*: $d(x) \leq d(y) ; d(z) \leftrightarrow d(x) \leq d(y) \wedge d(x) \leq d(z)$
by (*metis d-commutative d-idempotent d-mult-left-lower-bound mult-isotone order-trans*)

lemma *d-mult-left-absorb-add*: $d(x) ; (d(x) + d(y)) = d(x)$
by (*metis add-commutative d-idempotent d-plus-one mult-left-dist-add mult-right-one*)

lemma *d-add-left-absorb-mult*: $d(x) + d(x) ; d(y) = d(x)$
by (*metis add-commutative d-mult-left-lower-bound less-eq-def*)

lemma *d-add-left-dist-mult*: $d(x) + d(y) ; d(z) = (d(x) + d(y)) ; (d(x) + d(z))$
by (*smt add-associative d-commutative d-idempotent d-mult-left-absorb-add mult-left-dist-add mult-right-dist-add*)

lemma *d-order*: $d(x) \leq d(y) \leftrightarrow d(x) = d(x) ; d(y)$
by (*metis d-mult-greatest-lower-bound d-mult-left-absorb-add less-eq-def order-refl*)

lemma *d-mult-below*: $d(x) ; y \leq y$
by (*metis add-left-divisibility d-plus-one mult-left-one mult-right-dist-add*)

lemma *d-preserves-equation*: $d(y) ; x \leq x ; d(y) \leftrightarrow d(y) ; x = d(y) ; x ; d(y)$
by (*smt antisym d-below-one d-idempotent mult-associative mult-left-isotone mult-left-one mult-right-isotone mult-right-one*)

end

class *semiring-dL-below* = *domain-semiring* +
fixes $L :: 'a$
assumes *L-left-zero-below*: $L ; x \leq L$
assumes *mult-L-split*: $x ; L = x ; 0 + d(x) ; L$

begin

lemma *d-zero-mult-L*: $d(x ; 0) ; L \leq x$
by (*metis add-least-upper-bound mult-L-split mult-associative mult-left-zero zero-right-mult-decreasing*)

lemma *mult-L*: $x ; L \leq x ; 0 + L$
by (*metis add-right-isotone d-mult-below mult-L-split*)

lemma *d-mult-L*: $d(x) ; L \leq x ; L$
by (*metis add-right-divisibility mult-L-split*)

lemma *d-L-split*: $x ; d(y) ; L = x ; 0 + d(x ; y) ; L$
by (*metis d-commutative d-mult-d d-zero mult-L-split mult-associative mult-left-zero*)

lemma *d-mult-mult-L*: $d(x ; y) ; L \leq x ; d(y) ; L$
by (*metis add-right-divisibility d-L-split*)

lemma *L-L*: $L ; L = L$
by (*metis d-restrict-equals less-eq-def mult-L-split zero-right-mult-decreasing*)

end

class *ed-below* = *omega-algebra-T* + *semiring-dL-below* + *Omega* +
assumes *Omega-def*: $x^\Omega = d(x^\omega) ; L + x^*$

begin

lemma *Omega-isotone*: $x \leq y \rightarrow x^\Omega \leq y^\Omega$
by (*metis Omega-def add-isotone d-isotone mult-left-isotone omega-isotone star.circ-isotone*)

lemma *star-below-Omega*: $x^* \leq x^\Omega$
by (*metis Omega-def add-right-upper-bound*)

lemma *one-below-Omega*: $1 \leq x^\Omega$
by (*metis add-least-upper-bound star.circ-plus-one star-below-Omega*)

lemma *L-left-zero-star*: $L ; x^* = L$
by (*metis L-left-zero-below add-right-upper-bound antisym star.circ-back-loop-fixpoint*)

lemma *L-left-zero-Omega*: $L ; x^\Omega = L$
by (*metis L-left-zero-below L-left-zero-star Omega-def less-eq-def mult-left-dist-add*)

lemma *mult-L-star*: $(x ; L)^* = 1 + x ; L$
by (*metis L-left-zero-star mult-associative star.circ-left-unfold*)

lemma *mult-L-omega-below*: $(x ; L)^\omega \leq x ; L$
by (*metis L-left-zero-below mult-right-isotone omega-slide*)

lemma *mult-L-add-star*: $(x ; L + y)^* = y^* + y^* ; x ; L$
by (*metis L-left-zero-star add-commutative mult-associative star.circ-unfold-sum*)

lemma *mult-L-add-omega-below*: $(x ; L + y)^\omega \leq y^\omega + y^* ; x ; L$
proof –
have $(x ; L + y)^\omega = (y^* ; x ; L)^\omega + (y^* ; x ; L)^* ; y^\omega$
by (*metis add-commutative mult-associative omega-decompose*)
also have $\dots \leq y^* ; x ; L + (y^* ; x ; L)^* ; y^\omega$
by (*metis add-left-isotone mult-L-omega-below*)
also have $\dots = y^* ; x ; L + y^* ; x ; L ; y^\omega + y^\omega$

by (smt L-left-zero-star add-associative add-commutative mult-associative star.circ-loop-fixpoint)
also have ... $\leq y^\omega + y^* ; x ; L$
by (metis L-left-zero-star add-commutative eq-refl mult-associative star.circ-back-loop-fixpoint)
finally show ?thesis

qed

lemma *mult-L-add-circ*: $(x ; L + y)^\Omega = d(y^\omega) ; L + y^* + y^* ; x ; L$

proof –

have $(x ; L + y)^\Omega = d((x ; L + y)^\omega) ; L + (x ; L + y)^*$
by (metis Omega-def)
also have ... $\leq d(y^\omega + y^* ; x ; L) ; L + (x ; L + y)^*$
by (metis add-left-isotone d-isotone mult-L-add-omega-below mult-left-isotone)
also have ... $= d(y^\omega) ; L + d(y^* ; x ; L) ; L + (x ; L + y)^*$
by (metis d-dist-add mult-right-dist-add)
also have ... $\leq d(y^\omega) ; L + y^* ; x ; L ; L + (x ; L + y)^*$
by (metis add-left-isotone add-right-isotone d-mult-L)
also have ... $= d(y^\omega) ; L + y^* + y^* ; x ; L$
by (smt L-L add-associative add-commutative less-eq-def mult-L-add-star mult-associative order-refl)
finally have 1: $(x ; L + y)^\Omega \leq d(y^\omega) ; L + y^* + y^* ; x ; L$
.
have 2: $d(y^\omega) ; L \leq (x ; L + y)^\Omega$
by (metis Omega-def add-left-upper-bound add-right-upper-bound d-isotone mult-left-isotone omega-isotone order-trans)
have $y^* + y^* ; x ; L \leq (x ; L + y)^\Omega$
by (metis Omega-def add-right-upper-bound mult-L-add-star)
hence $d(y^\omega) ; L + y^* + y^* ; x ; L \leq (x ; L + y)^\Omega$ **using** 2
by (metis Omega-def add-least-upper-bound add-right-upper-bound mult-L-add-star)
thus ?thesis **using** 1
by (metis antisym)

qed

lemma *circ-add-d*: $(x^\Omega ; y)^\Omega ; x^\Omega = d((x^* ; y)^\omega) ; L + ((x^* ; y)^* ; x^* + (x^* ; y)^* ; d(x^\omega) ; L)$

proof –

have $(x^\Omega ; y)^\Omega ; x^\Omega = ((d(x^\omega) ; L + x^*) ; y)^\Omega ; x^\Omega$
by (metis Omega-def)
also have ... $= (d(x^\omega) ; L ; y + x^* ; y)^\Omega ; x^\Omega$
by (metis mult-right-dist-add)
also have ... $\leq (d(x^\omega) ; L + x^* ; y)^\Omega ; x^\Omega$
by (metis L-left-zero-below Omega-isotone add-left-isotone mult-associative mult-left-isotone mult-right-isotone)
also have ... $= (d((x^* ; y)^\omega) ; L + (x^* ; y)^* + (x^* ; y)^* ; d(x^\omega) ; L) ; x^\Omega$
by (metis mult-L-add-circ)
also have ... $= d((x^* ; y)^\omega) ; L ; x^\Omega + (x^* ; y)^* ; x^\Omega + (x^* ; y)^* ; d(x^\omega) ; L ; x^\Omega$
by (metis mult-right-dist-add)
also have ... $= d((x^* ; y)^\omega) ; L + (x^* ; y)^* ; x^\Omega + (x^* ; y)^* ; d(x^\omega) ; L$
by (smt L-left-zero-Omega mult-associative)
also have ... $= (d((x^* ; y)^\omega) ; L + ((x^* ; y)^* ; x^* + (x^* ; y)^* ; d(x^\omega) ; L)$
by (smt Omega-def add-associative add-commutative add-idempotent mult-associative mult-left-dist-add)

finally have 1: $(x^\Omega ; y)^\Omega ; x^\Omega \leq d((x^* ; y)^\omega) ; L + ((x^* ; y)^* ; x^* + (x^* ; y)^* ; d(x^\omega) ; L)$

•
have $d((x^* ; y)^\omega) ; L \leq (x^\Omega ; y)^\Omega$

by (*metis Omega-def Omega-isotone add-commutative add-right-upper-bound mult-left-isotone order-trans*)

also have $\dots \leq (x^\Omega ; y)^\Omega ; x^\Omega$

by (*metis mult-right-isotone mult-right-one one-below-Omega*)

finally have 2: $d((x^* ; y)^\omega) ; L \leq (x^\Omega ; y)^\Omega ; x^\Omega$

•
have 3: $(x^* ; y)^* ; x^* \leq (x^\Omega ; y)^\Omega ; x^\Omega$

by (*smt Omega-isotone mult-left-isotone mult-right-isotone order-trans star-below-Omega*)

have $(x^* ; y)^* ; d(x^\omega) ; L \leq (x^* ; y)^* ; x^\Omega$

by (*metis Omega-def add-commutative mult-associative mult-left-subdist-add-right*)

also have $\dots \leq (x^\Omega ; y)^\Omega ; x^\Omega$

by (*metis Omega-isotone mult-left-isotone order-trans star-below-Omega*)

finally have $d((x^* ; y)^\omega) ; L + ((x^* ; y)^* ; x^* + (x^* ; y)^* ; d(x^\omega) ; L) \leq (x^\Omega ; y)^\Omega ; x^\Omega$ using 2 3

by (*smt add-associative less-eq-def*)

thus *?thesis* using 1

by (*metis antisym*)

qed

end

class *ed* = *ed-below* +

assumes *L-left-zero*: $L ; x = L$

begin

lemma *mult-L-omega*: $(x ; L)^\omega = x ; L$

by (*metis L-left-zero omega-slide*)

lemma *mult-L-add-omega*: $(x ; L + y)^\omega = y^\omega + y^* ; x ; L$

by (*smt L-left-zero add-commutative add-left-upper-bound less-eq-def mult-L-omega mult-L-star mult-associative mult-left-one mult-right-dist-add omega-decompose*)

lemma *d-Omega-circ-simulate-right-plus*: $z ; x \leq y ; y^\Omega ; z + w \rightarrow z ; x^\Omega \leq y^\Omega ; (z + w ; x^\Omega)$

proof

assume $z ; x \leq y ; y^\Omega ; z + w$

hence $z ; x \leq y ; d(y^\omega) ; L ; z + y ; y^* ; z + w$

by (*metis Omega-def mult-associative mult-left-dist-add mult-right-dist-add*)

also have $\dots \leq y ; d(y^\omega) ; L + y ; y^* ; z + w$

by (*metis L-left-zero-below add-commutative add-right-isotone mult-associative mult-right-isotone*)

also have $\dots = y ; 0 + d(y ; y^\omega) ; L + y ; y^* ; z + w$

by (*metis d-L-split*)

also have $\dots = d(y^\omega) ; L + y ; y^* ; z + w$

by (*smt add-associative add-commutative add-left-zero mult-associative mult-left-dist-add omega-unfold*)

finally have 1: $z ; x \leq d(y^\omega) ; L + y ; y^* ; z + w$

have $(d(y^\omega); L + y^*; z + y^*; w; d(x^\omega); L + y^*; w; x^*); x = d(y^\omega); L; x + y^*; z; x + y^*; w; d(x^\omega); L; x + y^*; w; x^*; x$
by (*metis mult-right-dist-add*)
also have $\dots \leq d(y^\omega); L + y^*; z; x + y^*; w; d(x^\omega); L; x + y^*; w; x^*; x$
by (*metis L-left-zero-below add-left-isotone mult-associative mult-right-isotone*)
also have $\dots \leq d(y^\omega); L + y^*; z; x + y^*; w; d(x^\omega); L + y^*; w; x^*; x$
by (*metis L-left-zero-below add-commutative add-left-isotone mult-associative mult-right-isotone*)
also have $\dots \leq d(y^\omega); L + y^*; z; x + y^*; w; d(x^\omega); L + y^*; w; x^*$
by (*metis add-left-upper-bound add-right-isotone star.circ-back-loop-fixpoint*)
also have $\dots \leq d(y^\omega); L + y^*; (d(y^\omega); L + y; y^*; z + w) + y^*; w; d(x^\omega); L + y^*; w; x^*$ **using** 1
by (*smt add-left-isotone add-right-isotone less-eq-def mult-associative mult-left-dist-add*)
also have $\dots = d(y^\omega); L + y^*; y; y^*; z + y^*; w; d(x^\omega); L + y^*; w; x^*$
by (*smt add-associative add-commutative add-idempotent mult-associative mult-left-dist-add d-L-split star.circ-back-loop-fixpoint star-mult-omega*)
also have $\dots \leq d(y^\omega); L + y^*; z + y^*; w; d(x^\omega); L + y^*; w; x^*$
by (*metis add-left-isotone add-right-isotone mult-left-isotone star.circ-plus-same star.circ-transitive-equal star.left-plus-below-circ*)
finally have 2: $z; x^* \leq d(y^\omega); L + y^*; z + y^*; w; d(x^\omega); L + y^*; w; x^*$
by (*smt add-least-upper-bound add-left-upper-bound star.circ-loop-fixpoint star-right-induct*)
have $z; x; x^\omega \leq y; y^*; z; x^\omega + d(y^\omega); L; x^\omega + w; x^\omega$ **using** 1
by (*metis add-commutative mult-left-isotone mult-right-dist-add*)
also have $\dots \leq y; y^*; z; x^\omega + d(y^\omega); L + w; x^\omega$
by (*metis L-left-zero-below add-commutative add-right-isotone mult-associative mult-right-isotone*)
finally have $z; x^\omega \leq y^\omega + y^*; d(y^\omega); L + y^*; w; x^\omega$
by (*smt add-associative add-commutative left-plus-omega mult-associative mult-left-dist-add omega-induct omega-unfold star.left-plus-circ*)
hence $z; x^\omega \leq y^\omega + y^*; d(y^\omega); L + y^*; w; x^\omega$
by (*smt add-associative add-commutative left-plus-omega mult-associative mult-left-dist-add omega-induct omega-unfold star.left-plus-circ*)
hence $z; x^\omega \leq y^\omega + y^*; w; x^\omega$
by (*metis add-commutative d-mult-L less-eq-def mult-associative mult-right-isotone omega-sub-vector order-trans star-mult-omega*)
hence $d(z; x^\omega); L \leq d(y^\omega); L + y^*; w; d(x^\omega); L$
by (*smt add-associative add-commutative d-L-split d-dist-add less-eq-def mult-right-dist-add*)
hence $z; d(x^\omega); L \leq z; 0 + d(y^\omega); L + y^*; w; d(x^\omega); L$
by (*metis add-associative add-right-isotone d-L-split*)
hence $z; d(x^\omega); L \leq d(y^\omega); L + y^*; z + y^*; w; d(x^\omega); L + y^*; w; x^*$
by (*smt add-commutative add-left-isotone add-left-upper-bound order-trans star.circ-loop-fixpoint zero-right-mult-decreasing*)
thus $z; x^\Omega \leq y^\Omega; (z + w; x^\Omega)$ **using** 2
by (*smt L-left-zero Omega-def add-associative less-eq-def mult-associative mult-left-dist-add mult-right-dist-add*)

qed

lemma *d-Omega-circ-simulate-left-plus*: $x; z \leq z; y^\Omega + w \rightarrow x^\Omega; z \leq (z + x^\Omega; w); y^\Omega$

proof

assume 1: $x; z \leq z; y^\Omega + w$
have $x; (z; d(y^\omega); L + z; y^* + d(x^\omega); L + x^*; w; d(y^\omega); L + x^*; w; y^*) = x; z; d(y^\omega); L + x; z; y^* + d(x^\omega); L + x; x^*; w; d(y^\omega); L + x; x^*; w; y^*$
by (*smt add-associative add-commutative mult-associative mult-left-dist-add d-L-split omega-unfold*)
also have $\dots \leq (z; d(y^\omega); L + z; y^* + w); d(y^\omega); L + (z; d(y^\omega); L + z; y^* + w); y^* + d(x^\omega); L + x^*; w; d(y^\omega); L + x^*; w; y^*$ **using** 1
by (*smt Omega-def add-associative add-right-upper-bound less-eq-def mult-associative mult-left-dist-add mult-right-dist-add star.circ-loop-fixpoint*)
also have $\dots = z; d(y^\omega); L + z; y^*; d(y^\omega); L + w; d(y^\omega); L + z; y^* + w; y^* + d(x^\omega); L + x^*; w; d(y^\omega); L + x^*; w; y^*$
by (*smt L-left-zero add-associative add-commutative add-idempotent mult-associative mult-right-dist-add star.circ-transitive-equal*)
also have $\dots = z; d(y^\omega); L + w; d(y^\omega); L + z; y^* + w; y^* + d(x^\omega); L + x^*; w; d(y^\omega); L + x^*; w; y^*$

by (smt add-associative add-commutative add-idempotent less-eq-def mult-associative d-L-split star-mult-omega zero-right-mult-decreasing)
 finally have $x ; (z ; d(y^\omega) ; L + z ; y^* + d(x^\omega) ; L + x^* ; w ; d(y^\omega) ; L + x^* ; w ; y^*) \leq z ; d(y^\omega) ; L + z ; y^* + d(x^\omega) ; L + x^* ; w ; d(y^\omega) ; L + x^* ; w ; y^*$
 by (smt add-associative add-commutative add-idempotent mult-associative star.circ-loop-fixpoint)
 thus $x^\Omega ; z \leq (z + x^\Omega ; w) ; y^\Omega$
 by (smt L-left-zero Omega-def add-associative add-least-upper-bound add-left-upper-bound mult-associative mult-left-dist-add mult-right-dist-add star.circ-back-loop-fixpoint star-left-induct)
 qed

end

sublocale $ed < ed\text{-}\omega!$: ω -itering where $circ = \Omega$

apply unfold-locales

apply (smt L-left-zero add-associative add-commutative add-left-zero Omega-def mult-associative mult-left-dist-add mult-right-dist-add d-L-split omega-slide star-unfold-slide)

apply (smt add-associative add-commutative add-left-zero circ-add-d Omega-def mult-left-dist-add mult-right-dist-add d-L-split d-dist-add omega-decompose star.circ-add-1 star.circ-slide)

apply (smt L-left-zero add-associative add-commutative add-isotone add-left-zero Omega-def mult-associative mult-left-dist-add mult-left-isotone mult-right-dist-add d-L-split d-isotone omega-simulation star-simulation-right)

apply (smt d-Omega-circ-simulate-right-plus Omega-def add-left-isotone add-right-upper-bound mult-left-isotone mult-right-isotone order-trans star.circ-back-loop-fixpoint)

apply (smt Omega-def add-commutative add-least-upper-bound add-right-isotone mult-right-isotone d-Omega-circ-simulate-left-plus order-trans star.circ-increasing)

apply (metis d-Omega-circ-simulate-right-plus)

by (metis d-Omega-circ-simulate-left-plus)

sublocale $ed < ed\text{-}\star!$: ω -itering where $circ = \star ..$

class meet =

fixes meet :: 'a \Rightarrow 'a \Rightarrow 'a (infixl \frown 65)

class semiring-lattice = semiring-T + meet +

assumes meet-associative : $(x \frown y) \frown z = x \frown (y \frown z)$

assumes meet-commutative : $x \frown y = y \frown x$

assumes meet-idempotent : $x \frown x = x$

assumes meet-left-zero : $0 \frown x = 0$

assumes meet-left-top : $T \frown x = x$

assumes meet-left-dist-add : $x \frown (y + z) = (x \frown y) + (x \frown z)$

assumes add-left-dist-meet : $x + (y \frown z) = (x + y) \frown (x + z)$

assumes meet-absorb : $x \frown (x + y) = x$

assumes add-absorb : $x + (x \frown y) = x$

begin

lemma less-eq-meet: $x \leq y \leftrightarrow x \frown y = x$

by (metis add-absorb add-right-upper-bound less-eq-def meet-absorb meet-commutative)

lemma meet-left-isotone: $x \leq y \rightarrow x \frown z \leq y \frown z$

by (smt less-eq-def meet-commutative meet-left-dist-add)

lemma meet-right-isotone: $x \leq y \rightarrow z \frown x \leq z \frown y$

by (*metis meet-commutative meet-left-isotone*)

lemma *meet-isotone*: $w \leq y \wedge x \leq z \rightarrow w \frown x \leq y \frown z$
by (*smt less-eq-meet meet-associative meet-commutative*)

lemma *meet-left-upper-bound*: $x \frown y \leq x$
by (*metis add-absorb add-right-divisibility*)

lemma *meet-right-upper-bound*: $x \frown y \leq y$
by (*metis meet-commutative meet-left-upper-bound*)

lemma *meet-least-upper-bound*: $z \leq x \wedge z \leq y \leftrightarrow z \leq x \frown y$
by (*metis less-eq-meet meet-associative meet-left-upper-bound*)

lemma *meet-left-divisibility*: $y \leq x \leftrightarrow (\exists z . x \frown z = y)$
by (*metis less-eq-meet meet-commutative meet-left-upper-bound*)

lemma *meet-right-divisibility*: $y \leq x \leftrightarrow (\exists z . z \frown x = y)$
by (*metis meet-commutative meet-left-divisibility*)

lemma *meet-right-zero*: $x \frown 0 = 0$
by (*metis meet-commutative meet-left-zero*)

lemma *mult-left-subdist-meet-left*: $x ; (y \frown z) \leq x ; y$
by (*metis meet-left-upper-bound mult-right-isotone*)

lemma *mult-left-subdist-meet-right*: $x ; (y \frown z) \leq x ; z$
by (*metis meet-commutative mult-left-subdist-meet-left*)

lemma *mult-right-subdist-meet-left*: $(x \frown y) ; z \leq x ; z$
by (*metis meet-left-upper-bound mult-left-isotone*)

lemma *mult-right-subdist-meet-right*: $(x \frown y) ; z \leq y ; z$
by (*metis meet-right-upper-bound mult-left-isotone*)

lemma *mult-same-context*: $x \leq y \frown z \wedge y \leq x \frown z \rightarrow x \frown z = y \frown z$
by (*metis eq-iff meet-least-upper-bound*)

lemma *mult-right-top*: $x \frown T = x$
by (*metis meet-commutative meet-left-top*)

lemma *vector-mult-closed*: $\text{vector } x \wedge \text{vector } y \rightarrow \text{vector } (x \frown y)$
by (*metis antisym meet-least-upper-bound mult-right-subdist-meet-left mult-right-subdist-meet-right top-right-mult-increasing vector-def*)

lemma *relative-equality*: $x + z = y + z \wedge x \frown z = y \frown z \rightarrow x = y$
by (*metis add-absorb add-commutative add-left-dist-meet*)

end

class *domain-semiring-lattice* = *domain-semiring* + *semiring-lattice*

begin

lemma *d-top*: $d(T) = 1$

by (*metis add-left-top d-dist-add d-one d-plus-one*)

lemma *mult-domain-top*: $x ; d(y) ; T \leq d(x ; y) ; T$

by (*smt d-mult-d d-restrict-equals mult-associative mult-right-isotone top-greatest*)

lemma *meet-domain*: $x \frown d(y) ; z = d(y) ; (x \frown z)$

apply (*rule antisym*)

apply (*smt add-commutative add-right-divisibility d-export d-isotone d-mult-left-lower-bound d-plus-one d-restrict-equals meet-right-isotone meet-right-upper-bound mult-left-isotone mult-left-one mult-right-isotone order-trans*)

by (*metis d-plus-one meet-least-upper-bound mult-left-one mult-left-subdist-meet-right mult-right-subdist-add-left*)

lemma *meet-intro-domain*: $x \frown y = d(y) ; x \frown y$

by (*metis d-restrict-equals meet-commutative meet-domain*)

lemma *meet-domain-top*: $x \frown d(y) ; T = d(y) ; x$

by (*metis meet-domain mult-right-top*)

lemma *d-galois*: $d(x) \leq d(y) \leftrightarrow x \leq d(y) ; T$

by (*metis d-isotone d-least-left-presenter meet-domain-top meet-least-upper-bound*)

lemma *vector-meet*: $x ; T \frown y \leq d(x) ; y$

by (*metis d-galois d-mult-sub meet-commutative meet-domain-top meet-right-isotone*)

end

class *domain-semiring-lattice-L* = *domain-semiring-lattice* +

fixes *L* :: 'a

assumes *l1*: $x ; L = x ; 0 + d(x) ; L$

assumes *l2*: $d(L) ; x \leq x ; d(L)$

assumes *l3*: $d(L) ; T \leq L + d(L ; 0) ; T$

assumes *l4*: $L ; T \leq L$

assumes *l5*: $x ; 0 \frown L \leq (x \frown L) ; 0$

begin

lemma *l8*: $(x \frown L) ; 0 \leq x ; 0 \frown L$

by (*metis meet-commutative meet-least-upper-bound mult-right-subdist-meet-right zero-right-mult-decreasing*)

lemma l9: $x ; 0 \frown L \leq d(x ; 0) ; L$

by (*metis d-restrict-equals meet-commutative meet-domain meet-right-upper-bound*)

lemma l10: $L ; L = L$

by (*metis d-restrict-equals l1 less-eq-def zero-right-mult-decreasing*)

lemma l11: $d(x) ; L \leq x ; L$

by (*metis add-right-upper-bound l1*)

lemma l12: $d(x ; 0) ; L \leq x ; 0$

by (*metis add-right-divisibility l1 mult-associative mult-left-zero*)

lemma l13: $d(x ; 0) ; L \leq x$

by (*metis l12 order-trans zero-right-mult-decreasing*)

lemma l14: $x ; L \leq x ; 0 + L$

by (*metis add-right-isotone l1 meet-domain-top meet-left-upper-bound*)

lemma l15: $x ; d(y) ; L = x ; 0 + d(x ; y) ; L$

by (*metis d-commutative d-mult-d d-zero l1 mult-associative mult-left-zero*)

lemma l16: $x ; T \frown L \leq x ; L$

by (*metis l11 order-trans vector-meet*)

lemma l17: $d(x) ; L \leq d(x ; L) ; L$

by (*smt d-restrict-equals l11 meet-domain-top meet-least-upper-bound meet-left-divisibility order-trans*)

lemma l18: $d(x) ; L = d(x ; L) ; L$

by (*metis antisym d-mult-sub l17 mult-left-isotone*)

lemma l19: $d(x ; T ; 0) ; L \leq d(x ; L) ; L$

by (*metis d-mult-sub l18 mult-associative mult-left-isotone*)

lemma l20: $x \leq y \leftrightarrow x \leq y + L \wedge x \leq y + d(y ; 0) ; T$

apply rule

apply (*metis add-absorb less-eq-meet meet-left-dist-add*)

by (*smt add-commutative add-left-dist-meet l13 less-eq-def meet-domain-top*)

lemma l21: $d(x ; 0) ; L \leq x ; 0 \frown L$

by (*metis l12 meet-domain-top meet-least-upper-bound meet-left-upper-bound*)

lemma l22: $x ; 0 \frown L = d(x ; 0) ; L$

by (*metis antisym l9 l21*)

lemma l23: $x ; T \frown L = d(x) ; L$

apply (*rule antisym*)
apply (*metis vector-meet*)
by (*metis add-least-upper-bound d-mult-below l1 meet-least-upper-bound mult-right-isotone top-greatest*)

lemma l29: $L ; d(L) = L$
by (*metis d-preserves-equation d-restrict-equals l2*)

lemma l30: $d(L) ; x \leq (x \frown L) + d(L ; 0) ; x$
by (*metis l3 less-eq-def meet-domain-top meet-left-dist-add*)

lemma l31: $d(L) ; x = (x \frown L) + d(L ; 0) ; x$
by (*smt add-least-upper-bound antisym d-mult-sub d-restrict-equals l30 meet-domain mult-left-isotone mult-left-subdist-meet-left*)

lemma l40: $L ; x \leq L$
by (*metis add-right-top l4 mult-left-subdist-add-left order-trans*)

lemma l41: $L ; T = L$
by (*metis antisym l4 top-right-mult-increasing*)

lemma l50: $x ; 0 \frown L = (x \frown L) ; 0$
by (*metis eq-iff l5 l8*)

lemma l51: $d(x ; 0) ; L = (x \frown L) ; 0$
by (*metis l22 l50*)

lemma l90: $L ; T ; L = L$
by (*metis add-commutative d-restrict-equals d-top l1 l51 l31 meet-absorb meet-commutative mult-associative mult-left-dist-add mult-left-zero mult-right-one*)

lemma l91: $x = x ; T \rightarrow d(L ; 0) ; x \leq d(x ; 0) ; T$
proof –

have $d(L ; 0) ; x \leq d(d(L ; 0) ; x) ; T$
by (*metis d-galois order-refl*)

also have $\dots = d(d(L ; 0) ; d(x)) ; T$
by (*metis d-mult-d*)

also have $\dots = d(d(x) ; L ; 0) ; T$
by (*metis d-commutative d-mult-d mult-associative*)

also have $\dots \leq d(x ; L ; 0) ; T$
by (*metis d-isotone l11 mult-left-isotone*)

also have $\dots \leq d(x ; T ; 0) ; T$
by (*metis d-isotone mult-left-isotone mult-right-isotone top-greatest*)

finally show *?thesis*
by *metis*

qed

lemma l92: $x = x ; T \rightarrow d(L ; 0) ; x \leq d((x \frown L) ; 0) ; T$
proof

assume *1:* $x = x ; T$

have $d(L ; 0) ; x = d(L) ; d(L ; 0) ; x$
by (*metis d-commutative d-mult-sub d-order*)
also have $\dots \leq d(L) ; d(x ; 0) ; T$ **using** 1
by (*metis eq-iff l91 mult-associative mult-isotone*)
also have $\dots = d(d(x ; 0) ; L) ; T$
by (*metis d-commutative d-export*)
also have $\dots \leq d((x \frown L) ; 0) ; T$
by (*metis l51 order-refl*)
finally show $d(L ; 0) ; x \leq d((x \frown L) ; 0) ; T$
by *metis*

qed

end

class *domain-itering-lattice-L = itering-T + domain-semiring-lattice-L*

begin

lemma *mult-L-circ*: $(x ; L)^\circ = 1 + x ; L$
by (*metis circ-back-loop-fixpoint circ-mult l40 less-eq-def mult-associative*)

lemma *mult-L-circ-mult-below*: $(x ; L)^\circ ; y \leq y + x ; L$
by (*smt add-right-isotone l40 mult-L-circ mult-associative mult-left-one mult-right-dist-add mult-right-isotone*)

lemma *circ-L*: $L^\circ = L + 1$
by (*metis add-commutative l10 mult-L-circ*)

lemma *circ-d0-L*: $x^\circ ; d(x ; 0) ; L = x^\circ ; 0$
by (*metis add-right-zero circ-loop-fixpoint circ-plus-same d-zero l15 mult-associative mult-left-zero*)

lemma *d0-circ-left-unfold*: $d(x^\circ ; 0) = d(x ; x^\circ ; 0)$
by (*metis add-commutative add-left-zero circ-loop-fixpoint mult-associative*)

lemma *d-circ-import*: $d(y) ; x \leq x ; d(y) \rightarrow d(y) ; x^\circ = d(y) ; (d(y) ; x)^\circ$
by (*rule, rule antisym,*
metis circ-simulate circ-slide mult-associative d-idempotent d-preserves-equation,
metis circ-isotone mult-left-isotone mult-left-one mult-right-isotone d-below-one)

end

class *domain-omega-algebra-lattice-L = omega-algebra-T + domain-semiring-lattice-L*

begin

lemma *mult-L-star*: $(x ; L)^* = 1 + x ; L$

by (metis l40 less-eq-def mult-associative star.circ-back-loop-fixpoint star.circ-mult)

lemma *mult-L-omega*: $(x ; L)^\omega \leq x ; L$

by (smt l41 less-eq-def mult-associative mult-left-dist-add omega-unfold top-greatest)

lemma *mult-L-add-star*: $(x ; L + y)^* = y^* + y^* ; x ; L$

proof (rule antisym)

have $(x ; L + y) ; (y^* + y^* ; x ; L) = x ; L ; (y^* + y^* ; x ; L) + y ; (y^* + y^* ; x ; L)$

by (metis mult-associative mult-right-dist-add)

also have $\dots \leq x ; L + y ; (y^* + y^* ; x ; L)$

by (metis add-left-isotone l40 mult-associative mult-right-isotone)

also have $\dots \leq x ; L + y ; y^* + y^* ; x ; L$

by (smt add-associative add-commutative add-right-upper-bound mult-associative mult-left-dist-add star.circ-loop-fixpoint)

also have $\dots \leq x ; L + y^* + y^* ; x ; L$

by (metis add-left-isotone add-right-isotone star.left-plus-below-circ)

also have $\dots = y^* + y^* ; x ; L$

by (metis add-associative add-commutative mult-associative star.circ-loop-fixpoint star.circ-reflexive star.circ-sup-one-right-unfold star-involutive)

finally have $1 + (x ; L + y) ; (y^* + y^* ; x ; L) \leq y^* + y^* ; x ; L$

by (metis add-commutative add-least-upper-bound add-right-divisibility star.circ-left-unfold)

thus $(x ; L + y)^* \leq y^* + y^* ; x ; L$

by (metis mult-right-one star-left-induct)

next

show $y^* + y^* ; x ; L \leq (x ; L + y)^*$

by (metis add-commutative add-least-upper-bound mult-associative star.circ-increasing star.circ-mult-upper-bound star.circ-sub-dist)

qed

lemma *mult-L-add-omega*: $(x ; L + y)^\omega \leq y^\omega + y^* ; x ; L$

proof –

have 1: $(y^* ; x ; L)^\omega \leq y^\omega + y^* ; x ; L$

by (metis add-least-upper-bound add-right-isotone mult-L-omega)

have $(y^* ; x ; L)^* ; y^\omega \leq y^\omega + y^* ; x ; L$

by (metis add-right-isotone l40 mult-associative mult-right-isotone star-left-induct)

thus *?thesis* using 1

by (smt add-associative add-commutative less-eq-def mult-associative omega-decompose)

qed

end

sublocale *domain-omega-algebra-lattice-L* < *dL-star!*: *omega-itering* **where** *circ* = *star* ..

sublocale *domain-omega-algebra-lattice-L* < *dL-star!*: *domain-itering-lattice-L* **where** *circ* = *star* ..

context *domain-omega-algebra-lattice-L*

begin

lemma *d0-star-below-d0-omega*: $d(x^* ; 0) \leq d(x^\omega ; 0)$

by (*metis d-isotone star-zero-below-omega-zero*)

lemma *d0-below-d0-omega*: $d(x ; 0) \leq d(x^\omega ; 0)$

by (*metis d0-star-below-d0-omega d-isotone mult-left-isotone order-trans star.circ-increasing*)

lemma *star-d0-L*: $x^* ; d(x ; 0) ; L = x^* ; 0$

by (*metis dL-star.circ-d0-L*)

lemma *star-L-split*: $y \leq z \wedge x ; z ; L \leq x ; 0 + z ; L \rightarrow x^* ; y ; L \leq x^* ; 0 + z ; L$

by (*smt add-least-upper-bound add-left-isotone star.left-plus-below-circ mult-associative mult-left-dist-add mult-left-isotone order-trans star.circ-increasing star-left-induct*)

lemma *star-L-split-same*: $x ; y ; L \leq x ; 0 + y ; L \rightarrow x^* ; y ; L = x^* ; 0 + y ; L$

by (*smt add-associative add-left-zero antisym less-eq-def mult-associative mult-left-dist-add mult-left-one mult-right-subdist-add-left order-refl star-L-split star.circ-right-unfold*)

lemma *star-d-L-split-equal*: $d(x ; y) \leq d(y) \rightarrow x^* ; d(y) ; L = x^* ; 0 + d(y) ; L$

by (*metis add-right-isotone l15 less-eq-def mult-right-subdist-add-left star-L-split-same*)

lemma *d0-omega-mult*: $d(x^\omega ; y ; 0) = d(x^\omega ; 0)$

by (*smt antisym mult-associative mult-left-isotone mult-right-isotone omega-sub-vector zero-least*)

lemma *d-star-left-unfold*: $d(x^* ; 0) = d(x ; x^* ; 0)$

by (*metis dL-star.d0-circ-left-unfold*)

lemma *d-star-import*: $d(y) ; x \leq x ; d(y) \rightarrow d(y) ; x^* = d(y) ; (d(y) ; x)^*$

by (*metis dL-star.d-circ-import*)

lemma *d-omega-export*: $d(y) ; x \leq x ; d(y) \rightarrow d(y) ; x^\omega = (d(y) ; x)^\omega$

apply (*rule impI, rule antisym*)

apply (*metis d-preserves-equation omega-simulation order-refl*)

by (*smt less-eq-def mult-left-dist-add omega-simulation-2 omega-slide*)

lemma *d-omega-import*: $d(y) ; x \leq x ; d(y) \rightarrow d(y) ; x^\omega = d(y) ; (d(y) ; x)^\omega$

by (*metis d-idempotent d-omega-export mult-associative omega-slide*)

lemma *star-d-omega-top*: $x^* ; d(x^\omega) ; T = x^* ; 0 + d(x^\omega) ; T$

apply (*rule antisym*)

apply (*metis add-right-upper-bound dL-star.circ-mult-omega mult-domain-top order-trans*)

by (*metis add-commutative add-least-upper-bound add-left-divisibility dL-star.star-zero-below-circ-mult mult-associative star.circ-loop-fixpoint*)

lemma *omega-meet-L*: $x^\omega \frown L = d(x^\omega) ; L$

by (*metis l23 omega-vector*)

end

class *domain-semiring-lattice-apx* = *domain-semiring-lattice-L* +
fixes *apx* :: 'a ⇒ 'a ⇒ bool (**infix** \sqsubseteq 50)
assumes *apx-def*: $x \sqsubseteq y \leftrightarrow x \leq y + L \wedge d(L) ; y \leq x + d(x ; 0) ; T$

begin

definition *apx-isotone* :: ('a ⇒ 'a) ⇒ bool
where *apx-isotone* $f \leftrightarrow (\forall x y . x \sqsubseteq y \rightarrow f(x) \sqsubseteq f(y))$

lemma *apx-reflexive*: $x \sqsubseteq x$
by (*metis add-least-upper-bound add-left-upper-bound apx-def d-plus-one mult-left-one mult-right-dist-add*)

lemma *apx-L-least*: $L \sqsubseteq x$
by (*metis add-right-upper-bound apx-def l3 mult-right-isotone order-trans top-greatest*)

lemma *apx-transitive*: $x \sqsubseteq y \wedge y \sqsubseteq z \rightarrow x \sqsubseteq z$

proof

assume 1: $x \sqsubseteq y \wedge y \sqsubseteq z$

hence 2: $x \leq z + L$

by (*smt add-associative add-commutative apx-def less-eq-def*)

have $d(d(L) ; y ; 0) ; T \leq d((x + d(x ; 0) ; T) ; 0) ; T$ **using** 1

by (*metis apx-def d-isotone mult-left-isotone*)

also have $\dots \leq d(x ; 0) ; T$

by (*metis add-least-upper-bound d-galois mult-left-isotone mult-right-dist-add order-refl zero-right-mult-decreasing*)

finally have 3: $d(d(L) ; y ; 0) ; T \leq d(x ; 0) ; T$

by *metis*

have $d(L) ; z = d(L) ; (d(L) ; z)$

by (*metis d-idempotent mult-associative*)

also have $\dots \leq d(L) ; y + d(d(L) ; y ; 0) ; T$ **using** 1

by (*metis apx-def calculation d-export less-eq-meet meet-domain mult-associative mult-left-dist-add*)

also have $\dots \leq x + d(x ; 0) ; T$ **using** 1 3

by (*smt add-least-upper-bound add-right-upper-bound apx-def order-trans*)

finally show $x \sqsubseteq z$ **using** 2

by (*metis apx-def*)

qed

lemma *apx-meet-L*: $y \sqsubseteq x \rightarrow x \frown L \leq y \frown L$

proof

assume 1: $y \sqsubseteq x$

have $x \frown L = d(L) ; x \frown L$

by (*metis d-restrict-equals meet-commutative meet-domain*)

also have $\dots \leq (y + d(y ; 0) ; T) \frown L$ **using** 1

by (*metis apx-def meet-left-isotone*)

also have $\dots \leq y$

by (*metis add-least-upper-bound l13 meet-commutative meet-domain meet-left-dist-add meet-right-upper-bound mult-right-top*)
finally show $x \frown L \leq y \frown L$
 by (*metis meet-associative meet-idempotent meet-left-isotone*)

qed

lemma *apx-antisymmetric*: $x \sqsubseteq y \wedge y \sqsubseteq x \rightarrow x = y$
 by (*metis add-same-context antisym apx-def apx-meet-L relative-equality*)

lemma *add-apx-left-isotone*: $x \sqsubseteq y \rightarrow x + z \sqsubseteq y + z$

proof

assume 1: $x \sqsubseteq y$
hence 2: $x + z \leq y + z + L$
 by (*smt add-associative add-commutative add-left-isotone apx-def*)
have $d(L) ; (y + z) = d(L) ; y + d(L) ; z$
 by (*metis mult-left-dist-add*)
also have $\dots \leq d(L) ; y + z$
 by (*metis add-commutative add-least-upper-bound add-right-upper-bound d-below-one mult-left-dist-add mult-left-isotone mult-left-one*)
also have $\dots \leq x + d(x ; 0) ; T + z$ **using** 1
 by (*metis add-left-isotone apx-def*)
also have $\dots \leq x + z + d((x + z) ; 0) ; T$
 by (*smt add-associative add-commutative add-right-isotone d-isotone mult-left-isotone mult-right-subdist-add-left*)
finally show $x + z \sqsubseteq y + z$ **using** 2
 by (*metis apx-def*)

qed

lemma *add-apx-right-isotone*: $x \sqsubseteq y \rightarrow z + x \sqsubseteq z + y$
 by (*metis add-commutative add-apx-left-isotone*)

lemma *add-apx-isotone*: $w \sqsubseteq y \wedge x \sqsubseteq z \rightarrow w + x \sqsubseteq y + z$
 by (*metis add-apx-left-isotone add-apx-right-isotone apx-transitive*)

lemma *mult-apx-left-isotone*: $x \sqsubseteq y \rightarrow x ; z \sqsubseteq y ; z$

proof

assume 1: $x \sqsubseteq y$
hence $x ; z \leq y ; z + L ; z$
 by (*metis apx-def mult-left-isotone mult-right-dist-add*)
hence 2: $x ; z \leq y ; z + L$
 by (*metis add-commutative add-left-isotone l40 order-trans*)
have $d(L) ; y ; z \leq x ; z + d(x ; 0) ; T ; z$ **using** 1
 by (*metis apx-def mult-left-isotone mult-right-dist-add*)
also have $\dots \leq x ; z + d(x ; z ; 0) ; T$
 by (*metis add-right-isotone d-isotone mult-associative mult-isotone mult-right-isotone top-greatest zero-least*)
finally show $x ; z \sqsubseteq y ; z$ **using** 2
 by (*metis apx-def mult-associative*)

qed

lemma *mult-apx-right-isotone*: $x \sqsubseteq y \rightarrow z ; x \sqsubseteq z ; y$

proof

assume $I: x \sqsubseteq y$
hence $z ; x \leq z ; y + z ; L$
by (*metis apx-def mult-left-dist-add mult-right-isotone*)
also have $\dots \leq z ; y + z ; 0 + L$
by (*metis add-associative add-right-isotone l14*)
finally have $2: z ; x \leq z ; y + L$
by (*metis add-right-zero mult-left-dist-add*)
have $d(L) ; z ; y \leq z ; d(L) ; y$
by (*metis l2 mult-left-isotone*)
also have $\dots \leq z ; (x + d(x ; 0)) ; T$ **using** 1
by (*metis apx-def mult-associative mult-right-isotone*)
also have $\dots = z ; x + z ; d(x ; 0) ; T$
by (*metis mult-associative mult-left-dist-add*)
also have $\dots \leq z ; x + d(z ; x ; 0) ; T$
by (*metis add-right-isotone mult-associative mult-domain-top*)
finally show $z ; x \sqsubseteq z ; y$ **using** 2
by (*metis apx-def mult-associative*)

qed

lemma *mult-apx-isotone*: $w \sqsubseteq y \wedge x \sqsubseteq z \rightarrow w ; x \sqsubseteq y ; z$

by (*metis apx-transitive mult-apx-left-isotone mult-apx-right-isotone*)

lemma *meet-L-apx-isotone*: $x \sqsubseteq y \rightarrow x \frown L \sqsubseteq y \frown L$

by (*smt add-absorb add-commutative apx-def apx-meet-L d-restrict-equals l20 meet-commutative meet-domain meet-left-upper-bound*)

definition *is-apx-prefixpoint* :: $('a \Rightarrow 'a) \Rightarrow 'a \Rightarrow \text{bool}$ **where** *is-apx-prefixpoint* $f x \leftrightarrow f(x) \sqsubseteq x$

definition *is-apx-least-fixpoint* :: $('a \Rightarrow 'a) \Rightarrow 'a \Rightarrow \text{bool}$ **where** *is-apx-least-fixpoint* $f x \leftrightarrow f(x) = x \wedge (\forall y . f(y) = y \rightarrow x \sqsubseteq y)$

definition *is-apx-least-prefixpoint* :: $('a \Rightarrow 'a) \Rightarrow 'a \Rightarrow \text{bool}$ **where** *is-apx-least-prefixpoint* $f x \leftrightarrow f(x) \sqsubseteq x \wedge (\forall y . f(y) \sqsubseteq y \rightarrow x \sqsubseteq y)$

definition *has-apx-prefixpoint* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$ **where** *has-apx-prefixpoint* $f \leftrightarrow (\exists x . \text{is-apx-prefixpoint } f x)$

definition *has-apx-least-fixpoint* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$ **where** *has-apx-least-fixpoint* $f \leftrightarrow (\exists x . \text{is-apx-least-fixpoint } f x)$

definition *has-apx-least-prefixpoint* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$ **where** *has-apx-least-prefixpoint* $f \leftrightarrow (\exists x . \text{is-apx-least-prefixpoint } f x)$

definition *the-apx-least-fixpoint* :: $('a \Rightarrow 'a) \Rightarrow 'a$ (ξ - [201] 200) **where** $\xi f = (\text{THE } x . \text{is-apx-least-fixpoint } f x)$

definition *the-apx-least-prefixpoint* :: $('a \Rightarrow 'a) \Rightarrow 'a$ ($p\xi$ - [201] 200) **where** $p\xi f = (\text{THE } x . \text{is-apx-least-prefixpoint } f x)$

lemma *apx-least-fixpoint-unique*: $\text{has-apx-least-fixpoint } f \rightarrow (\exists! x . \text{is-apx-least-fixpoint } f x)$

by (*smt apx-antisymmetric has-apx-least-fixpoint-def is-apx-least-fixpoint-def*)

lemma *apx-least-prefixpoint-unique*: $\text{has-apx-least-prefixpoint } f \rightarrow (\exists! x . \text{is-apx-least-prefixpoint } f x)$

by (*smt apx-antisymmetric has-apx-least-prefixpoint-def is-apx-least-prefixpoint-def*)

lemma *apx-least-fixpoint*: $\text{has-apx-least-fixpoint } f \rightarrow \text{is-apx-least-fixpoint } f$ (ξf)

proof

assume *has-apx-least-fixpoint* f

hence *is-apx-least-fixpoint* f ($\text{THE } x . \text{is-apx-least-fixpoint } f x$)

by (*smt apx-least-fixpoint-unique theI'*)

thus *is-apx-least-fixpoint* f (ξf)

by (simp add: is-apx-least-fixpoint-def the-apx-least-fixpoint-def)
qed

lemma *apx-least-prefixpoint*: $has-apx-least-prefixpoint f \rightarrow is-apx-least-prefixpoint f (p\xi f)$

proof

assume *has-apx-least-prefixpoint* f

hence *is-apx-least-prefixpoint* f (THE $x . is-apx-least-prefixpoint f x$)

by (smt *apx-least-prefixpoint-unique theI'*)

thus *is-apx-least-prefixpoint* f ($p\xi f$)

by (simp add: is-apx-least-prefixpoint-def the-apx-least-prefixpoint-def)

qed

lemma *apx-least-fixpoint-same*: $is-apx-least-fixpoint f x \rightarrow x = \xi f$

by (metis *apx-least-fixpoint apx-least-fixpoint-unique has-apx-least-fixpoint-def*)

lemma *apx-least-prefixpoint-same*: $is-apx-least-prefixpoint f x \rightarrow x = p\xi f$

by (metis *apx-least-prefixpoint apx-least-prefixpoint-unique has-apx-least-prefixpoint-def*)

lemma *apx-least-fixpoint-char*: $is-apx-least-fixpoint f x \leftrightarrow has-apx-least-fixpoint f \wedge x = \xi f$

by (metis *apx-least-fixpoint-same has-apx-least-fixpoint-def*)

lemma *apx-least-prefixpoint-char*: $is-apx-least-prefixpoint f x \leftrightarrow has-apx-least-prefixpoint f \wedge x = p\xi f$

by (metis *apx-least-prefixpoint-same has-apx-least-prefixpoint-def*)

lemma *apx-least-prefixpoint-fixpoint*: $has-apx-least-prefixpoint f \wedge apx-isotone f \rightarrow is-apx-least-fixpoint f (p\xi f)$

by (smt *apx-antisymmetric apx-isotone-def apx-reflexive is-apx-least-fixpoint-def is-apx-least-prefixpoint-def apx-least-prefixpoint*)

lemma *pxi-xi*: $has-apx-least-prefixpoint f \wedge apx-isotone f \rightarrow p\xi f = \xi f$

by (smt *has-apx-least-fixpoint-def is-apx-least-fixpoint-def apx-least-fixpoint-unique apx-least-prefixpoint-fixpoint apx-least-fixpoint*)

definition *lifted-apx-less-eq* :: $('a \Rightarrow 'a) \Rightarrow ('a \Rightarrow 'a) \Rightarrow bool$ ($(- \sqsubseteq\sqsubseteq -)$ [51, 51] 50)

where $f \sqsubseteq\sqsubseteq g \leftrightarrow (\forall x . f(x) \sqsubseteq g(x))$

lemma *lifted-reflexive*: $f = g \rightarrow f \sqsubseteq\sqsubseteq g$

by (metis *lifted-apx-less-eq-def apx-reflexive*)

lemma *lifted-transitive*: $f \sqsubseteq\sqsubseteq g \wedge g \sqsubseteq\sqsubseteq h \rightarrow f \sqsubseteq\sqsubseteq h$

by (smt *lifted-apx-less-eq-def apx-transitive*)

lemma *lifted-antisymmetric*: $f \sqsubseteq\sqsubseteq g \wedge g \sqsubseteq\sqsubseteq f \rightarrow f = g$

by (metis *apx-antisymmetric ext lifted-apx-less-eq-def*)

lemma *pxi-isotone*: $has-apx-least-prefixpoint f \wedge has-apx-least-prefixpoint g \wedge f \sqsubseteq\sqsubseteq g \rightarrow p\xi f \sqsubseteq p\xi g$

by (metis *is-apx-least-prefixpoint-def apx-transitive apx-least-prefixpoint lifted-apx-less-eq-def*)

lemma *xi-isotone*: $has-apx-least-prefixpoint f \wedge has-apx-least-prefixpoint g \wedge apx-isotone f \wedge apx-isotone g \wedge f \sqsubseteq\sqsubseteq g \rightarrow \xi f \sqsubseteq \xi g$

by (metis *pxi-isotone pxi-xi*)

lemma *xi-square*: $\text{apx-isotone } f \wedge \text{has-apx-least-fixpoint } f \wedge \text{has-apx-least-fixpoint } (f \circ f) \rightarrow \xi f = \xi (f \circ f)$
by (*metis apx-antisymmetric is-apx-least-fixpoint-def apx-isotone-def apx-least-fixpoint-char apx-least-fixpoint-unique o-apply*)

lemma *mu-below-xi*: $\text{has-least-fixpoint } f \wedge \text{has-apx-least-fixpoint } f \rightarrow \mu f \leq \xi f$
by (*metis apx-least-fixpoint is-apx-least-fixpoint-def is-least-fixpoint-def least-fixpoint*)

lemma *xi-below-nu*: $\text{has-greatest-fixpoint } f \wedge \text{has-apx-least-fixpoint } f \rightarrow \xi f \leq \nu f$
by (*metis apx-least-fixpoint greatest-fixpoint is-apx-least-fixpoint-def is-greatest-fixpoint-def*)

lemma *xi-apx-below-mu*: $\text{has-least-fixpoint } f \wedge \text{has-apx-least-fixpoint } f \rightarrow \xi f \sqsubseteq \mu f$
by (*metis apx-least-fixpoint is-apx-least-fixpoint-def is-least-fixpoint-def least-fixpoint*)

lemma *xi-apx-below-nu*: $\text{has-greatest-fixpoint } f \wedge \text{has-apx-least-fixpoint } f \rightarrow \xi f \sqsubseteq \nu f$
by (*metis apx-least-fixpoint greatest-fixpoint is-apx-least-fixpoint-def is-greatest-fixpoint-def*)

definition *is-apx-meet* :: $'a \Rightarrow 'a \Rightarrow 'a \Rightarrow \text{bool}$ **where** $\text{is-apx-meet } x \ y \ z \leftrightarrow z \sqsubseteq x \wedge z \sqsubseteq y \wedge (\forall w . w \sqsubseteq x \wedge w \sqsubseteq y \rightarrow w \sqsubseteq z)$
definition *has-apx-meet* :: $'a \Rightarrow 'a \Rightarrow \text{bool}$ **where** $\text{has-apx-meet } x \ y \leftrightarrow (\exists z . \text{is-apx-meet } x \ y \ z)$
definition *the-apx-meet* :: $'a \Rightarrow 'a \Rightarrow 'a$ (**infixl** \square 66) **where** $x \square y = (\text{THE } z . \text{is-apx-meet } x \ y \ z)$

lemma *apx-meet-unique*: $\text{has-apx-meet } x \ y \rightarrow (\exists! z . \text{is-apx-meet } x \ y \ z)$
by (*smt apx-antisymmetric has-apx-meet-def is-apx-meet-def*)

lemma *apx-meet*: $\text{has-apx-meet } x \ y \rightarrow \text{is-apx-meet } x \ y \ (x \square y)$
proof

assume $\text{has-apx-meet } x \ y$
hence $\text{is-apx-meet } x \ y \ (\text{THE } z . \text{is-apx-meet } x \ y \ z)$
by (*smt apx-meet-unique theI'*)
thus $\text{is-apx-meet } x \ y \ (x \square y)$
by (*simp add: is-apx-meet-def the-apx-meet-def*)

qed

lemma *apx-greatest-lower-bound*: $\text{has-apx-meet } x \ y \rightarrow (w \sqsubseteq x \wedge w \sqsubseteq y \leftrightarrow w \sqsubseteq x \square y)$
by (*smt apx-meet apx-transitive is-apx-meet-def*)

lemma *apx-meet-same*: $\text{is-apx-meet } x \ y \ z \rightarrow z = x \square y$
by (*metis apx-meet apx-meet-unique has-apx-meet-def*)

lemma *apx-meet-char*: $\text{is-apx-meet } x \ y \ z \leftrightarrow \text{has-apx-meet } x \ y \wedge z = x \square y$
by (*metis apx-meet-same has-apx-meet-def*)

definition *xi-apx-meet* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where $\text{xi-apx-meet } f \leftrightarrow \text{has-apx-least-fixpoint } f \wedge \text{has-apx-meet } (\mu f) (\nu f) \wedge \xi f = \mu f \square \nu f$

definition *xi-mu-nu* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where $\text{xi-mu-nu } f \leftrightarrow \text{has-apx-least-fixpoint } f \wedge \xi f = \mu f + (\nu f \frown L)$

definition *nu-below-mu-nu* :: ('a ⇒ 'a) ⇒ bool
where *nu-below-mu-nu* f ↔ d(L) ; ν f ≤ μ f + (ν f ∩ L) + d(ν f ; 0) ; T

definition *nu-below-mu-nu-2* :: ('a ⇒ 'a) ⇒ bool
where *nu-below-mu-nu-2* f ↔ d(L) ; ν f ≤ μ f + (ν f ∩ L) + d((μ f + (ν f ∩ L)) ; 0) ; T

definition *mu-nu-apx-nu* :: ('a ⇒ 'a) ⇒ bool
where *mu-nu-apx-nu* f ↔ μ f + (ν f ∩ L) ⊆ ν f

definition *mu-nu-apx-meet* :: ('a ⇒ 'a) ⇒ bool
where *mu-nu-apx-meet* f ↔ *has-apx-meet* (μ f) (ν f) ∧ μ f ⊓ ν f = μ f + (ν f ∩ L)

definition *apx-meet-below-nu* :: ('a ⇒ 'a) ⇒ bool
where *apx-meet-below-nu* f ↔ *has-apx-meet* (μ f) (ν f) ∧ μ f ⊓ ν f ≤ ν f

lemma *mu-below-l*: μ f ≤ μ f + (ν f ∩ L)
by (*metis add-left-upper-bound*)

lemma *l-below-nu*: *has-least-fixpoint* f ∧ *has-greatest-fixpoint* f → μ f + (ν f ∩ L) ≤ ν f
by (*metis add-least-upper-bound meet-left-upper-bound mu-below-nu*)

lemma *n-l-nu*: *has-least-fixpoint* f ∧ *has-greatest-fixpoint* f → (μ f + (ν f ∩ L)) ∩ L = ν f ∩ L
by (*smt add-commutative add-left-dist-meet less-eq-def meet-absorb meet-associative meet-commutative mu-below-nu*)

lemma *l-apx-mu*: μ f + (ν f ∩ L) ⊆ μ f
by (*metis add-right-isotone apx-def meet-absorb meet-domain-top meet-least-upper-bound meet-left-upper-bound meet-right-upper-bound*)

lemma *nu-below-mu-nu-nu-below-mu-nu-2*: *nu-below-mu-nu* f → *nu-below-mu-nu-2* f

proof

assume 1: *nu-below-mu-nu* f

have d(L) ; ν f = d(L) ; (d(L) ; ν f)

by (*metis d-idempotent mult-associative*)

also have ... ≤ d(L) ; (μ f + (ν f ∩ L) + d(ν f ; 0)) ; T **using** 1

by (*metis mult-right-isotone nu-below-mu-nu-def*)

also have ... = d(L) ; (μ f + (ν f ∩ L)) + d(L) ; d(ν f ; 0) ; T

by (*metis mult-associative mult-left-dist-add*)

also have ... ≤ μ f + (ν f ∩ L) + d(L) ; d(ν f ; 0) ; T

by (*metis add-left-isotone meet-domain-top meet-left-upper-bound*)

also have ... = μ f + (ν f ∩ L) + d(d(ν f ; 0) ; L) ; T

by (*smt d-commutative d-export*)

also have ... = μ f + (ν f ∩ L) + d((ν f ∩ L) ; 0) ; T

by (*metis l51*)

also have ... ≤ μ f + (ν f ∩ L) + d((μ f + (ν f ∩ L)) ; 0) ; T

by (*metis add-right-isotone add-right-upper-bound d-dist-add mult-right-dist-add*)

finally show *nu-below-mu-nu-2* f

by (*metis nu-below-mu-nu-2-def*)

qed

lemma *nu-below-mu-nu-2-nu-below-mu-nu*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f \rightarrow nu\text{-}below\text{-}mu\text{-}nu\ f$
proof

assume 1: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f$

hence $d(L) ; \nu f \leq \mu f + (\nu f \frown L) + d((\mu f + (\nu f \frown L)) ; 0) ; T$

by (*metis nu-below-mu-nu-2-def*)

also have $\dots \leq \mu f + (\nu f \frown L) + d(\nu f ; 0) ; T$ **using** 1

by (*smt add-absorb add-associative add-commutative d-dist-add l-below-nu less-eq-def meet-absorb mult-right-dist-add*)

finally show $nu\text{-}below\text{-}mu\text{-}nu\ f$

by (*metis nu-below-mu-nu-def*)

qed

lemma *nu-below-mu-nu-equivalent*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \rightarrow (nu\text{-}below\text{-}mu\text{-}nu\ f \leftrightarrow nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f)$

by (*metis nu-below-mu-nu-2-nu-below-mu-nu nu-below-mu-nu-nu-below-mu-nu-2*)

lemma *nu-below-mu-nu-2-mu-nu-apx-nu*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f \rightarrow mu\text{-}nu\text{-}apx\text{-}nu\ f$

proof

assume 1: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f$

hence $\mu f + (\nu f \frown L) \leq \nu f + L$

by (*metis add-commutative add-right-upper-bound l-below-nu order-trans*)

thus $mu\text{-}nu\text{-}apx\text{-}nu\ f$ **using** 1

by (*metis apx-def mu-nu-apx-nu-def nu-below-mu-nu-2-def*)

qed

lemma *mu-nu-apx-nu-mu-nu-apx-meet*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge mu\text{-}nu\text{-}apx\text{-}nu\ f \rightarrow mu\text{-}nu\text{-}apx\text{-}meet\ f$

proof

let $?l = \mu f + (\nu f \frown L)$

assume $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge mu\text{-}nu\text{-}apx\text{-}nu\ f$

hence $is\text{-}apx\text{-}meet\ (\mu f)\ (\nu f)\ ?l$

by (*smt add-apx-left-isotone add-commutative apx-meet-L is-apx-meet-def l-apx-mu less-eq-def meet-least-upper-bound mu-nu-apx-nu-def*)

thus $mu\text{-}nu\text{-}apx\text{-}meet\ f$

by (*smt apx-meet-char mu-nu-apx-meet-def*)

qed

lemma *mu-nu-apx-meet-apx-meet-below-nu*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge mu\text{-}nu\text{-}apx\text{-}meet\ f \rightarrow apx\text{-}meet\text{-}below\text{-}nu\ f$

by (*metis apx-meet-below-nu-def l-below-nu mu-nu-apx-meet-def*)

lemma *apx-meet-below-nu-nu-below-mu-nu-2*: $apx\text{-}meet\text{-}below\text{-}nu\ f \rightarrow nu\text{-}below\text{-}mu\text{-}nu\text{-}2\ f$

proof –

let $?l = \mu f + (\nu f \frown L)$

have $\forall m . m \sqsubseteq \mu f \wedge m \sqsubseteq \nu f \wedge m \leq \nu f \rightarrow d(L) ; \nu f \leq ?l + d(?l ; 0) ; T$

proof

fix m

show $m \sqsubseteq \mu f \wedge m \sqsubseteq \nu f \wedge m \leq \nu f \rightarrow d(L) ; \nu f \leq ?l + d(?l ; 0) ; T$

proof

assume 1: $m \sqsubseteq \mu f \wedge m \sqsubseteq \nu f \wedge m \leq \nu f$

hence $m \leq ?l$

by (smt add-commutative add-left-dist-meet add-left-upper-bound apx-def less-eq-meet meet-least-upper-bound)
 hence $m + d(m ; 0) ; T \leq ?l + d(?l ; 0) ; T$
 by (metis add-isotone d-dist-add less-eq-def mult-right-dist-add)
 thus $d(L) ; \nu f \leq ?l + d(?l ; 0) ; T$ using 1
 by (smt apx-def order-trans)
 qed
 qed
 thus ?thesis
 by (smt apx-meet-below-nu-def apx-meet-same apx-meet-unique is-apx-meet-def nu-below-mu-nu-2-def)
 qed

lemma *has-apx-least-fixpoint-xi-apx-meet*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge has\text{-}apx\text{-}least\text{-}fixpoint\ f \rightarrow xi\text{-}apx\text{-}meet\ f$
proof
 assume 1: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge has\text{-}apx\text{-}least\text{-}fixpoint\ f$
 hence $\forall w . w \sqsubseteq \mu f \wedge w \sqsubseteq \nu f \rightarrow w \sqsubseteq \xi f$
 by (smt add-left-isotone apx-def mu-below-xi mult-right-isotone order-trans xi-below-nu)
 thus $xi\text{-}apx\text{-}meet\ f$ using 1
 by (smt apx-meet-char is-apx-meet-def xi-apx-below-mu xi-apx-below-nu xi-apx-meet-def)
 qed

lemma *xi-apx-meet-apx-meet-below-nu*: $has\text{-}greatest\text{-}fixpoint\ f \wedge xi\text{-}apx\text{-}meet\ f \rightarrow apx\text{-}meet\text{-}below\text{-}nu\ f$
 by (metis apx-meet-below-nu-def xi-apx-meet-def xi-below-nu)

lemma *apx-meet-below-nu-xi-mu-nu*: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge isotone\ f \wedge apx\text{-}isotone\ f \wedge apx\text{-}meet\text{-}below\text{-}nu\ f \rightarrow xi\text{-}mu\text{-}nu\ f$
proof
 let $?l = \mu f + (\nu f \frown L)$
 let $?m = \mu f \sqcap \nu f$
 assume 1: $has\text{-}least\text{-}fixpoint\ f \wedge has\text{-}greatest\text{-}fixpoint\ f \wedge isotone\ f \wedge apx\text{-}isotone\ f \wedge apx\text{-}meet\text{-}below\text{-}nu\ f$
 hence 2: $?m = ?l$
 by (metis apx-meet-below-nu-nu-below-mu-nu-2 mu-nu-apx-meet-def mu-nu-apx-nu-mu-nu-apx-meet nu-below-mu-nu-2-mu-nu-apx-nu)
 have 3: $?l \leq f(?l) + L$
proof –
 have $?l \leq \mu f + L$
 by (metis add-right-isotone meet-right-upper-bound)
 also have $\dots = f(\mu f) + L$ using 1
 by (metis is-least-fixpoint-def least-fixpoint)
 also have $\dots \leq f(?l) + L$ using 1
 by (metis add-left-isotone add-left-upper-bound isotone-def)
 finally show $?l \leq f(?l) + L$
 by metis
 qed
 have $d(L) ; f(?l) \leq ?l + d(?l ; 0) ; T$
proof –
 have $d(L) ; f(?l) \leq d(L) ; f(\nu f)$ using 1 2
 by (metis apx-meet-below-nu-def isotone-def mult-right-isotone)
 also have $\dots = d(L) ; \nu f$ using 1
 by (metis greatest-fixpoint is-greatest-fixpoint-def)

also have $\dots \leq ?l + d(?l ; 0) ; T$ **using** 1
by (*metis apx-meet-below-nu-nu-below-mu-nu-2 nu-below-mu-nu-2-def*)
finally show $d(L) ; f(?l) \leq ?l + d(?l ; 0) ; T$
by *metis*
qed
hence 4: $?l \sqsubseteq f(?l)$ **using** 3
by (*metis apx-def*)
have 5: $f(?l) \sqsubseteq \mu f$
proof –
have $?l \sqsubseteq \mu f$
by (*metis l-apx-mu*)
thus $f(?l) \sqsubseteq \mu f$ **using** 1
by (*metis apx-isotone-def is-least-fixpoint-def least-fixpoint*)
qed
have 6: $f(?l) \sqsubseteq \nu f$
proof –
have $?l \sqsubseteq \nu f$ **using** 1 2
by (*metis apx-greatest-lower-bound apx-meet-below-nu-def apx-reflexive*)
thus $f(?l) \sqsubseteq \nu f$ **using** 1
by (*metis apx-isotone-def greatest-fixpoint is-greatest-fixpoint-def*)
qed
hence $f(?l) \sqsubseteq ?l$ **using** 1 2 5
by (*metis apx-greatest-lower-bound apx-meet-below-nu-def*)
hence 7: $f(?l) = ?l$ **using** 4
by (*metis apx-antisymmetric*)
have $\forall y . f(y) = y \rightarrow ?l \sqsubseteq y$
proof
fix y
show $f(y) = y \rightarrow ?l \sqsubseteq y$
proof
assume 8: $f(y) = y$
hence 9: $?l \leq y + L$ **using** 1
by (*metis add-isotone is-least-fixpoint-def least-fixpoint meet-right-upper-bound*)
have $y \leq \nu f$ **using** 1 8
by (*metis greatest-fixpoint is-greatest-fixpoint-def*)
hence $d(L) ; y \leq ?l + d(?l ; 0) ; T$ **using** 4 6
by (*smt apx-def apx-transitive mult-right-isotone order-trans*)
thus $?l \sqsubseteq y$ **using** 9
by (*metis apx-def*)
qed
qed
thus *xi-mu-nu f* **using** 1 2 7
by (*smt apx-least-fixpoint-same has-apx-least-fixpoint-def is-apx-least-fixpoint-def xi-mu-nu-def*)
qed

lemma *xi-mu-nu-has-apx-least-fixpoint*: $xi-mu-nu f \rightarrow has-apx-least-fixpoint f$
by (*metis xi-mu-nu-def*)

lemma *nu-below-mu-nu-xi-mu-nu*: $\text{has-least-fixpoint } f \wedge \text{has-greatest-fixpoint } f \wedge \text{isotone } f \wedge \text{apx-isotone } f \wedge \text{nu-below-mu-nu } f \rightarrow \text{xi-mu-nu } f$
by (*metis apx-meet-below-nu-xi-mu-nu mu-nu-apx-meet-apx-meet-below-nu mu-nu-apx-nu-mu-nu-apx-meet nu-below-mu-nu-nu-below-mu-nu-2 nu-below-mu-nu-2-mu-nu-apx-nu*)

lemma *xi-mu-nu-nu-below-mu-nu*: $\text{has-least-fixpoint } f \wedge \text{has-greatest-fixpoint } f \wedge \text{xi-mu-nu } f \rightarrow \text{nu-below-mu-nu } f$
by (*metis apx-meet-below-nu-nu-below-mu-nu-2 has-apx-least-fixpoint-xi-apx-meet nu-below-mu-nu-2-nu-below-mu-nu xi-apx-meet-apx-meet-below-nu xi-mu-nu-has-apx-least-fixpoint*)

definition *xi-mu-nu-L* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where *xi-mu-nu-L* $f \leftrightarrow \text{has-apx-least-fixpoint } f \wedge \xi f = \mu f + d(\nu f ; 0) ; L$

definition *nu-below-mu-nu-L* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where *nu-below-mu-nu-L* $f \leftrightarrow d(L) ; \nu f \leq \mu f + d(\nu f ; 0) ; T$

definition *mu-nu-apx-nu-L* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where *mu-nu-apx-nu-L* $f \leftrightarrow \mu f + d(\nu f ; 0) ; L \sqsubseteq \nu f$

definition *mu-nu-apx-meet-L* :: $('a \Rightarrow 'a) \Rightarrow \text{bool}$
where *mu-nu-apx-meet-L* $f \leftrightarrow \text{has-apx-meet } (\mu f) (\nu f) \wedge \mu f \sqcap \nu f = \mu f + d(\nu f ; 0) ; L$

lemma *n-below-l*: $x + d(y ; 0) ; L \leq x + (y \frown L)$
by (*metis add-right-isotone d-mult-below l13 meet-least-upper-bound*)

lemma *n-equal-l*: $\text{nu-below-mu-nu-L } f \rightarrow \mu f + d(\nu f ; 0) ; L = \mu f + (\nu f \frown L)$

proof

assume *nu-below-mu-nu-L* f
hence $\nu f \frown L \leq (\mu f + d(\nu f ; 0) ; T) \frown L$
by (*smt meet-associative meet-intro-domain meet-right-divisibility nu-below-mu-nu-L-def*)
also have $\dots \leq \mu f + d(\nu f ; 0) ; L$
by (*smt add-left-dist-meet add-right-divisibility meet-commutative meet-domain-top meet-left-isotone*)
finally have $\mu f + (\nu f \frown L) \leq \mu f + d(\nu f ; 0) ; L$
by (*metis add-least-upper-bound add-left-upper-bound*)
thus $\mu f + d(\nu f ; 0) ; L = \mu f + (\nu f \frown L)$
by (*metis antisym n-below-l*)

qed

lemma *nu-below-mu-nu-L-nu-below-mu-nu*: $\text{nu-below-mu-nu-L } f \rightarrow \text{nu-below-mu-nu } f$
by (*metis add-associative add-right-top mult-left-dist-add n-equal-l nu-below-mu-nu-L-def nu-below-mu-nu-def*)

lemma *nu-below-mu-nu-L-xi-mu-nu-L*: $\text{has-least-fixpoint } f \wedge \text{has-greatest-fixpoint } f \wedge \text{isotone } f \wedge \text{apx-isotone } f \wedge \text{nu-below-mu-nu-L } f \rightarrow \text{xi-mu-nu-L } f$
by (*metis n-equal-l nu-below-mu-nu-L-nu-below-mu-nu nu-below-mu-nu-xi-mu-nu xi-mu-nu-L-def xi-mu-nu-def*)

lemma *nu-below-mu-nu-L-mu-nu-apx-nu-L*: $\text{has-least-fixpoint } f \wedge \text{has-greatest-fixpoint } f \wedge \text{nu-below-mu-nu-L } f \rightarrow \text{mu-nu-apx-nu-L } f$
by (*metis mu-nu-apx-nu-L-def mu-nu-apx-nu-def n-equal-l nu-below-mu-nu-2-mu-nu-apx-nu nu-below-mu-nu-L-nu-below-mu-nu nu-below-mu-nu-nu-below-mu-nu-2*)

lemma *nu-below-mu-nu-L-mu-nu-apx-meet-L*: $\text{has-least-fixpoint } f \wedge \text{has-greatest-fixpoint } f \wedge \text{nu-below-mu-nu-L } f \rightarrow \text{mu-nu-apx-meet-L } f$
by (*metis mu-nu-apx-meet-L-def mu-nu-apx-meet-def mu-nu-apx-nu-mu-nu-apx-meet n-equal-l nu-below-mu-nu-2-mu-nu-apx-nu nu-below-mu-nu-L-nu-below-mu-nu nu-below-mu-nu-nu-below-mu-nu-2*)

lemma *mu-nu-apx-nu-L-nu-below-mu-nu-L*: *has-least-fixpoint f* \wedge *has-greatest-fixpoint f* \wedge *mu-nu-apx-nu-L f* \rightarrow *nu-below-mu-nu-L f*

proof

let $?n = \mu f + d(\nu f ; 0)$; L

let $?l = \mu f + (\nu f \frown L)$

assume 1: *has-least-fixpoint f* \wedge *has-greatest-fixpoint f* \wedge *mu-nu-apx-nu-L f*

hence $d(L) ; \nu f \leq ?n + d(?n ; 0)$; T

by (*metis apx-def mu-nu-apx-nu-L-def*)

also have $\dots \leq ?n + d(?l ; 0)$; T

by (*metis add-right-isotone d-isotone mult-left-isotone n-below-l*)

also have $\dots \leq ?n + d(\nu f ; 0)$; T **using** 1

by (*metis add-right-isotone d-isotone l-below-nu mult-left-isotone*)

finally show *nu-below-mu-nu-L f*

by (*metis add-associative add-right-top mult-left-dist-add nu-below-mu-nu-L-def*)

qed

lemma *xi-mu-nu-L-mu-nu-apx-nu-L*: *has-greatest-fixpoint f* \wedge *xi-mu-nu-L f* \rightarrow *mu-nu-apx-nu-L f*

by (*metis mu-nu-apx-nu-L-def xi-apx-below-nu xi-mu-nu-L-def*)

lemma *mu-nu-apx-meet-L-mu-nu-apx-nu-L*: *mu-nu-apx-meet-L f* \rightarrow *mu-nu-apx-nu-L f*

by (*metis apx-meet-same has-apx-meet-def is-apx-meet-def mu-nu-apx-meet-L-def mu-nu-apx-nu-L-def*)

lemma *xi-mu-nu-L-nu-below-mu-nu-L*: *has-least-fixpoint f* \wedge *has-greatest-fixpoint f* \wedge *xi-mu-nu-L f* \rightarrow *nu-below-mu-nu-L f*

by (*metis mu-nu-apx-nu-L-nu-below-mu-nu-L xi-mu-nu-L-mu-nu-apx-nu-L*)

end

class *itering-apx* = *domain-itering-lattice-L* + *domain-semiring-lattice-apx*

begin

lemma *circ-apx-isotone*: $x \sqsubseteq y \rightarrow x^\circ \sqsubseteq y^\circ$

proof

assume $x \sqsubseteq y$

hence 1: $x \leq y + L \wedge d(L) ; y \leq x + d(x ; 0)$; T

by (*metis apx-def*)

have $d(L) ; y^\circ \leq (d(L) ; y)^\circ$

by (*smt circ-reflexive circ-transitive-equal d-below-one d-circ-import l2 mult-left-isotone order-trans*)

also have $\dots \leq x^\circ ; (d(x ; 0) ; T ; x^\circ)^\circ$ **using** 1

by (*metis circ-add-1 circ-isotone*)

also have $\dots = x^\circ + x^\circ ; d(x ; 0) ; T$

by (*metis circ-left-top mult-associative mult-left-dist-add mult-right-one mult-top-circ*)

also have $\dots \leq x^\circ + d(x^\circ ; x ; 0) ; T$

by (*metis add-right-isotone mult-associative mult-domain-top*)

finally have 2: $d(L) ; y^\circ \leq x^\circ + d(x^\circ ; 0) ; T$

by (*metis circ-plus-same d0-circ-left-unfold*)

have $x^\circ \leq y^\circ ; L^\circ$ **using** 1

by (*metis circ-add-1 circ-back-loop-fixpoint circ-isotone l40 less-eq-def mult-associative*)

also have $\dots = y^\circ + y^\circ ; L$
by (*metis add-commutative circ-L mult-left-dist-add mult-right-one*)
also have $\dots \leq y^\circ + y^\circ ; 0 + L$
by (*metis add-associative add-right-isotone l14*)
finally have $x^\circ \leq y^\circ + L$
by (*metis add-commutative less-eq-def zero-right-mult-decreasing*)
thus $x^\circ \sqsubseteq y^\circ$ **using** 2
by (*metis apx-def*)

qed

end

class *omega-algebra-apx* = *domain-omega-algebra-lattice-L* + *domain-semiring-lattice-apx*

sublocale *omega-algebra-apx* < *star!*: *itering-apx* **where** *circ* = *star* ..

context *omega-algebra-apx*

begin

lemma *omega-apx-isotone*: $x \sqsubseteq y \rightarrow x^\omega \sqsubseteq y^\omega$

proof

assume $x \sqsubseteq y$

hence 1: $x \leq y + L \wedge d(L) ; y \leq x + d(x ; 0) ; T$

by (*metis apx-def*)

have $d(L) ; y^\omega = (d(L) ; y)^\omega$

by (*metis d-omega-export l2*)

also have $\dots \leq (x + d(x ; 0) ; T)^\omega$ **using** 1

by (*metis omega-isotone*)

also have $\dots = (x^* ; d(x ; 0) ; T)^\omega + (x^* ; d(x ; 0) ; T)^* ; x^\omega$

by (*metis mult-associative omega-decompose*)

also have $\dots \leq x^* ; d(x ; 0) ; T + (x^* ; d(x ; 0) ; T)^* ; x^\omega$

by (*metis add-left-isotone mult-top-omega*)

also have $\dots = x^* ; d(x ; 0) ; T + (1 + x^* ; d(x ; 0) ; T ; (x^* ; d(x ; 0) ; T)^*) ; x^\omega$

by (*metis mult-associative star.circ-left-top star.mult-top-circ*)

also have $\dots \leq x^\omega + x^* ; d(x ; 0) ; T$

by (*smt add-isotone add-least-upper-bound mult-associative mult-left-one mult-right-dist-add mult-right-isotone order-refl top-greatest*)

also have $\dots \leq x^\omega + d(x^* ; x ; 0) ; T$

by (*metis add-right-isotone mult-associative mult-domain-top*)

also have $\dots \leq x^\omega + d(x^* ; 0) ; T$

by (*metis dL-star.d0-circ-left-unfold eq-refl star.circ-plus-same*)

finally have 2: $d(L) ; y^\omega \leq x^\omega + d(x^\omega ; 0) ; T$

by (*smt add-right-isotone d0-star-below-d0-omega mult-left-isotone order-trans*)

have $x^\omega \leq (y + L)^\omega$ **using** 1

by (*metis omega-isotone*)

also have $\dots = (y^* ; L)^\omega + (y^* ; L)^* ; y^\omega$

by (*metis omega-decompose*)

also have $\dots = y^* ; L ; (y^* ; L)^\omega + (y^* ; L)^* ; y^\omega$
by (*metis omega-unfold*)
also have $\dots \leq y^* ; L + (y^* ; L)^* ; y^\omega$
by (*metis add-left-isotone l40 mult-associative mult-right-isotone*)
also have $\dots = y^* ; L + (1 + y^* ; L ; (y^* ; L)^*) ; y^\omega$
by (*metis star.circ-left-unfold*)
also have $\dots \leq y^* ; L + y^\omega$
by (*metis add-commutative add-least-upper-bound add-right-upper-bound dL-star.mult-L-circ-mult-below mult-associative star.circ-mult star.circ-slide*)
also have $\dots \leq y^* ; 0 + L + y^\omega$
by (*metis add-left-isotone l14*)
finally have $x^\omega \leq y^\omega + L$
by (*metis add-associative add-commutative less-eq-def star-zero-below-omega*)
thus $x^\omega \sqsubseteq y^\omega$ **using** 2
by (*metis apx-def*)
qed

lemma *combined-apx-isotone*: $x \sqsubseteq y \rightarrow (x^\omega \frown L) + x^* ; z \sqsubseteq (y^\omega \frown L) + y^* ; z$
by (*metis add-apx-isotone mult-apx-left-isotone omega-apx-isotone star.circ-apx-isotone meet-L-apx-isotone*)

lemma *d-split-nu-mu*: $d(L) ; (y^\omega + y^* ; z) \leq y^* ; z + ((y^\omega + y^* ; z) \frown L) + d((y^\omega + y^* ; z) ; 0) ; T$

proof –

have $d(L) ; y^\omega \leq (y^\omega \frown L) + d(y^\omega ; 0) ; T$
by (*metis add-right-isotone l31 l91 omega-vector*)
hence $d(L) ; (y^\omega + y^* ; z) \leq y^* ; z + (y^\omega \frown L) + d(y^\omega ; 0) ; T$
by (*smt add-associative add-commutative add-isotone d-mult-below mult-left-dist-add*)
thus *?thesis*
by (*smt add-commutative add-isotone add-right-isotone add-right-upper-bound d-isotone meet-commutative meet-right-isotone mult-left-isotone order-trans*)
qed

lemma *loop-exists*: $d(L) ; \nu (\lambda x . y ; x + z) \leq \mu (\lambda x . y ; x + z) + (\nu (\lambda x . y ; x + z) \frown L) + d(\nu (\lambda x . y ; x + z) ; 0) ; T$
by (*metis d-split-nu-mu omega-loop-nu star-loop-mu*)

lemma *loop-isotone*: *isotone* $(\lambda x . y ; x + z)$
by (*smt add-commutative add-right-isotone isotone-def mult-right-isotone*)

lemma *loop-apx-isotone*: *apx-isotone* $(\lambda x . y ; x + z)$
by (*smt add-apx-left-isotone apx-isotone-def mult-apx-right-isotone*)

lemma *loop-has-least-fixpoint*: *has-least-fixpoint* $(\lambda x . y ; x + z)$
by (*metis has-least-fixpoint-def star-loop-is-least-fixpoint*)

lemma *loop-has-greatest-fixpoint*: *has-greatest-fixpoint* $(\lambda x . y ; x + z)$
by (*metis has-greatest-fixpoint-def omega-loop-is-greatest-fixpoint*)

lemma *loop-apx-least-fixpoint*: *is-apx-least-fixpoint* $(\lambda x . y ; x + z)$ $(\mu (\lambda x . y ; x + z) + (\nu (\lambda x . y ; x + z) \frown L))$
by (*metis apx-least-fixpoint-char loop-apx-isotone loop-exists loop-has-greatest-fixpoint loop-has-least-fixpoint loop-isotone nu-below-mu-nu-def nu-below-mu-nu-xi-mu-nu xi-mu-nu-def*)

lemma *loop-has-apx-least-fixpoint*: $\text{has-apx-least-fixpoint } (\lambda x . y ; x + z)$
by (*metis has-apx-least-fixpoint-def loop-apx-least-fixpoint*)

lemma *loop-semantics*: $\xi (\lambda x . y ; x + z) = \mu (\lambda x . y ; x + z) + (\nu (\lambda x . y ; x + z) \frown L)$
by (*metis apx-least-fixpoint-char loop-apx-least-fixpoint*)

lemma *loop-semantics-xi-mu-nu*: $\xi (\lambda x . y ; x + z) = (y^\omega \frown L) + y^* ; z$

proof –

have $\xi (\lambda x . y ; x + z) = y^* ; z + ((y^\omega + y^* ; z) \frown L)$

by (*metis loop-semantics omega-loop-nu star-loop-mu*)

thus *?thesis*

by (*smt add-absorb add-associative add-commutative add-left-dist-meet*)

qed

lemma *loop-semantics-xi-mu-nu-domain*: $\xi (\lambda x . y ; x + z) = d(y^\omega) ; L + y^* ; z$
by (*metis loop-semantics-xi-mu-nu omega-meet-L*)

lemma *loop-semantics-apx-isotone*: $w \sqsubseteq y \rightarrow \xi (\lambda x . w ; x + z) \sqsubseteq \xi (\lambda x . y ; x + z)$
by (*metis loop-semantics-xi-mu-nu combined-apx-isotone*)

end

class *while* =

fixes *while* :: 'a \Rightarrow 'a \Rightarrow 'a (**infixr** * 60)

class *binary-itering* = *semiring* + *while* +

assumes *while-productstar*: $(x ; y) * z = z + x ; ((y ; x) * (y ; z))$

assumes *while-sumstar*: $(x + y) * z = (x * y) * (x * z)$

assumes *while-left-dist-add*: $x * (y + z) = (x * y) + (x * z)$

assumes *while-sub-associative*: $(x * y) ; z \leq x * (y ; z)$

assumes *while-simulate-left-plus*: $x ; z \leq z ; (y * 1) + w \rightarrow x * (z ; v) \leq z ; (y * v) + (x * (w ; (y * v)))$

assumes *while-simulate-right-plus*: $z ; x \leq y ; (y * z) + w \rightarrow z ; (x * v) \leq y * (z ; v + w ; (x * v))$

begin

lemma *while-zero*: $0 * x = x$

by (*metis add-right-zero mult-left-zero while-productstar*)

lemma *while-mult-increasing*: $x ; y \leq x * y$

by (*metis add-least-upper-bound mult-left-one order-refl while-productstar*)

lemma *while-one-increasing*: $x \leq x * 1$

by (*metis mult-right-one while-mult-increasing*)

lemma *while-increasing*: $y \leq x * y$

by (*metis add-left-divisibility mult-left-one while-productstar*)

lemma *while-right-isotone*: $y \leq z \rightarrow x * y \leq x * z$

by (*metis less-eq-def while-left-dist-add*)

lemma *while-left-isotone*: $x \leq y \rightarrow x * z \leq y * z$

by (*metis less-eq-def while-increasing while-sumstar*)

lemma *while-isotone*: $w \leq x \wedge y \leq z \rightarrow w * y \leq x * z$

by (*smt order-trans while-left-isotone while-right-isotone*)

lemma *while-left-unfold*: $x * y = y + x ; (x * y)$

by (*metis mult-left-one mult-right-one while-productstar*)

lemma *while-simulate-left-plus-1*: $x ; z \leq z ; (y * 1) \rightarrow x * (z ; w) \leq z ; (y * w) + (x * 0)$

by (*metis add-right-zero mult-left-zero while-simulate-left-plus*)

lemma *while-simulate-absorb*: $y ; x \leq x \rightarrow y * x \leq x + (y * 0)$

by (*metis while-simulate-left-plus-1 while-zero mult-right-one*)

lemma *while-transitive*: $x * (x * y) = x * y$

by (*metis add-right-upper-bound add-right-zero antisym while-increasing while-left-dist-add while-left-unfold while-simulate-absorb*)

lemma *while-slide*: $(x ; y) * (x ; z) = x ; ((y ; x) * z)$

by (*metis mult-associative mult-left-dist-add while-left-unfold while-productstar*)

lemma *while-zero-2*: $(x ; 0) * y = x ; 0 + y$

by (*metis add-commutative mult-associative mult-left-zero while-left-unfold*)

lemma *while-mult-star-exchange*: $x ; (x * y) = x * (x ; y)$

by (*metis mult-left-one while-slide*)

lemma *while-right-unfold*: $x * y = y + (x * (x ; y))$

by (*metis while-left-unfold while-mult-star-exchange*)

lemma *while-one-mult-below*: $(x * 1) ; y \leq x * y$

by (*metis mult-left-one while-sub-associative*)

lemma *while-plus-one*: $x * y = y + (x * y)$

by (*metis less-eq-def while-increasing*)

lemma *while-rtc-2*: $y + x ; y + (x * (x * y)) = x * y$

by (*metis add-associative less-eq-def while-mult-increasing while-plus-one while-transitive*)

lemma *while-left-plus-below*: $x ; (x * y) \leq x * y$

by (*metis add-right-divisibility while-left-unfold*)

lemma *while-right-plus-below*: $x * (x ; y) \leq x * y$
by (*metis while-left-plus-below while-mult-star-exchange*)

lemma *while-right-plus-below-2*: $(x * x) ; y \leq x * y$
by (*smt order-trans while-right-plus-below while-sub-associative*)

lemma *while-mult-transitive*: $x \leq z * y \wedge y \leq z * w \rightarrow x \leq z * w$
by (*smt order-trans while-right-isotone while-transitive*)

lemma *while-mult-upper-bound*: $x \leq z * 1 \wedge y \leq z * w \rightarrow x ; y \leq z * w$
by (*metis less-eq-def mult-right-subdist-add-left order-trans while-mult-transitive while-one-mult-below*)

lemma *while-sub-dist*: $x * z \leq (x + y) * z$
by (*metis add-left-upper-bound while-left-isotone*)

lemma *while-sub-dist-1*: $x ; z \leq (x + y) * z$
by (*metis order-trans while-mult-increasing while-sub-dist*)

lemma *while-sub-dist-2*: $x ; y ; z \leq (x + y) * z$
by (*smt mult-associative mult-right-subdist-add-right order-trans while-mult-increasing while-mult-transitive while-sub-dist-1*)

lemma *while-sub-dist-3*: $x * (y * z) \leq (x + y) * z$
by (*metis add-right-upper-bound while-left-isotone while-mult-transitive while-sub-dist*)

lemma *while-absorb-2*: $x \leq y \rightarrow y * (x * z) = y * z$
by (*metis add-commutative less-eq-def while-left-dist-add while-plus-one while-sub-dist-3*)

lemma *while-simulate-right-plus-1*: $z ; x \leq y ; (y * z) \rightarrow z ; (x * w) \leq y * (z ; w)$
by (*metis add-right-zero mult-left-zero while-simulate-right-plus*)

lemma *while-sumstar-1-below*: $x * ((y ; (x * 1)) * z) \leq ((x * 1) ; y) * (x * z)$

proof –

have 1: $x ; (((x * 1) ; y) * (x * z)) \leq ((x * 1) ; y) * (x * z)$

by (*smt add-isotone add-right-upper-bound mult-associative mult-left-dist-add mult-right-subdist-add-right while-left-unfold*)

have $x * ((y ; (x * 1)) * z) \leq (x * z) + (x * (y ; (((x * 1) ; y) * ((x * 1) ; z))))$

by (*metis eq-refl while-left-dist-add while-productstar*)

also have $\dots \leq (x * z) + (x * ((x * 1) ; y ; (((x * 1) ; y) * ((x * 1) ; z))))$

by (*metis add-right-isotone mult-associative mult-left-one mult-right-subdist-add-left while-left-unfold while-right-isotone*)

also have $\dots \leq (x * z) + (x * (((x * 1) ; y) * ((x * 1) ; z)))$

by (*metis add-right-isotone add-right-upper-bound while-left-unfold while-right-isotone*)

also have $\dots \leq x * (((x * 1) ; y) * (x * z))$

by (*smt add-associative add-left-upper-bound less-eq-def mult-left-one while-left-dist-add while-left-unfold while-sub-associative*)

also have $\dots \leq (((x * 1) ; y) * (x * z)) + (x * 0)$ **using** 1

by (*metis while-simulate-absorb*)

also have $\dots = ((x * 1) ; y) * (x * z)$

by (*smt add-associative add-commutative add-left-zero while-left-dist-add while-left-unfold*)

finally show *?thesis*

·
qed

lemma *while-sumstar-2-below*: $((x * 1) ; y) * (x * z) \leq (x * y) * (x * z)$
by (*metis mult-left-one while-left-isotone while-sub-associative*)

lemma *while-add-1-below*: $x * ((y ; (x * 1)) * z) \leq (x + y) * z$

proof –

have $((x * 1) ; y) * ((x * 1) ; z) \leq (x + y) * z$

by (*metis while-isotone while-one-mult-below while-sumstar*)

hence $(y ; (x * 1)) * z \leq z + y ; ((x + y) * z)$

by (*metis add-right-isotone mult-right-isotone while-productstar*)

also have $\dots \leq (x + y) * z$

by (*metis add-right-isotone add-right-upper-bound mult-left-isotone while-left-unfold*)

finally show *?thesis*

by (*metis add-commutative add-right-upper-bound while-isotone while-transitive*)

qed

lemma *while-while-while*: $((x * 1) * 1) * y = (x * 1) * y$

by (*smt add-commutative less-eq-def mult-right-one while-left-plus-below while-mult-star-exchange while-plus-one while-sumstar while-transitive*)

lemma *while-one*: $(1 * 1) * y = 1 * y$

by (*metis while-while-while while-zero*)

lemma *while-add-below*: $x + y \leq x * (y * 1)$

by (*smt add-commutative add-isotone case-split-right order-trans while-increasing while-left-plus-below while-mult-increasing while-plus-one*)

lemma *while-add-2*: $(x + y) * z \leq (x * (y * 1)) * z$

by (*metis while-add-below while-left-isotone*)

lemma *while-sup-one-left-unfold*: $1 \leq x \rightarrow x ; (x * y) = x * y$

by (*metis less-eq-def mult-left-one mult-right-dist-add while-mult-star-exchange while-right-unfold while-transitive*)

lemma *while-sup-one-right-unfold*: $1 \leq x \rightarrow x * (x ; y) = x * y$

by (*metis while-mult-star-exchange while-sup-one-left-unfold*)

lemma *while-decompose-7*: $(x + y) * z = x * (y * ((x + y) * z))$

by (*metis eq-iff order-trans while-increasing while-sub-dist-3 while-transitive*)

lemma *while-decompose-8*: $(x + y) * z = (x + y) * (x * (y * z))$

by (*metis add-commutative while-sumstar while-transitive*)

lemma *while-decompose-9*: $(x * (y * 1)) * z = x * (y * ((x * (y * 1)) * z))$

by (*smt add-commutative less-eq-def order-trans while-add-below while-increasing while-sub-dist-3*)

lemma *while-decompose-10*: $(x * (y * 1)) * z = (x * (y * 1)) * (x * (y * z))$

proof –

have $1: (x * (y * 1)) * z \leq (x * (y * 1)) * (x * (y * z))$
by (*metis add-associative less-eq-def while-left-dist-add while-plus-one*)
have $x + (y * 1) \leq x * (y * 1)$
by (*metis add-least-upper-bound while-add-below while-increasing*)
hence $(x * (y * 1)) * (x * (y * z)) \leq (x * (y * 1)) * z$
by (*smt antisym while-decompose-9 while-increasing while-mult-transitive while-right-isotone*)
thus *?thesis* **using** 1
by (*metis antisym*)
qed

lemma *while-back-loop-fixpoint*: $z ; (y * (y ; x)) + z ; x = z ; (y * x)$
by (*metis add-commutative mult-left-dist-add while-right-unfold*)

lemma *while-back-loop-prefixpoint*: $z ; (y * 1) ; y + z \leq z ; (y * 1)$
by (*metis add-least-upper-bound mult-associative mult-right-isotone mult-right-one order-refl while-increasing while-mult-upper-bound while-one-increasing*)

lemma *while-loop-is-fixpoint*: *is-fixpoint* $(\lambda x . y ; x + z) (y * z)$
by (*smt add-commutative is-fixpoint-def while-left-unfold*)

lemma *while-back-loop-is-prefixpoint*: *is-prefixpoint* $(\lambda x . x ; y + z) (z ; (y * 1))$
by (*metis is-prefixpoint-def while-back-loop-prefixpoint*)

lemma *while-while-add*: $(1 + x) * y = (x * 1) * y$
by (*metis add-commutative while-decompose-10 while-sumstar while-zero*)

lemma *while-while-mult-sub*: $x * (1 * y) \leq (x * 1) * y$
by (*metis add-commutative while-sub-dist-3 while-while-add*)

lemma *while-right-plus*: $(x * x) * y = x * y$
by (*metis add-idempotent while-plus-one while-sumstar while-transitive*)

lemma *while-left-plus*: $(x ; (x * 1)) * y = x * y$
by (*metis mult-right-one while-mult-star-exchange while-right-plus*)

lemma *while-below-while-one*: $x * x \leq x * 1$
by (*metis while-one-increasing while-right-plus*)

lemma *while-below-while-one-mult*: $x ; (x * x) \leq x ; (x * 1)$
by (*metis mult-right-isotone while-below-while-one*)

lemma *while-add-sub-add-one*: $x * (x + y) \leq x * (1 + y)$
by (*metis add-left-isotone while-below-while-one while-left-dist-add*)

lemma *while-add-sub-add-one-mult*: $x ; (x * (x + y)) \leq x ; (x * (1 + y))$
by (*metis mult-right-isotone while-add-sub-add-one*)

lemma *while-elimination*: $x ; y = 0 \rightarrow x ; (y * z) = x ; z$

by (*metis add-right-zero mult-associative mult-left-dist-add mult-left-zero while-left-unfold*)

lemma *while-square*: $(x ; x) * y \leq x * y$

by (*metis while-left-isotone while-mult-increasing while-right-plus*)

lemma *while-mult-sub-add*: $(x ; y) * z \leq (x + y) * z$

by (*metis while-increasing while-isotone while-mult-increasing while-sumstar*)

lemma *while-absorb-1*: $x \leq y \rightarrow x * (y * z) = y * z$

by (*metis antisym less-eq-def while-increasing while-sub-dist-3*)

lemma *while-absorb-3*: $x \leq y \rightarrow x * (y * z) = y * (x * z)$

by (*metis while-absorb-1 while-absorb-2*)

lemma *while-square-2*: $(x ; x) * ((x + 1) ; y) \leq x * y$

by (*smt add-least-upper-bound while-increasing while-mult-transitive while-mult-upper-bound while-one-increasing while-square*)

lemma *while-separate-unfold-below*: $(y ; (x * 1)) * z \leq (y * z) + (y * (y ; x ; (x * ((y ; (x * 1)) * z))))$

proof –

have $(y ; (x * 1)) * z = (y * (y ; x ; (x * 1))) * (y * z)$

by (*metis mult-associative mult-left-dist-add mult-right-one while-left-unfold while-sumstar*)

hence $(y ; (x * 1)) * z = (y * z) + (y * (y ; x ; (x * 1))) ; ((y ; (x * 1)) * z)$

by (*metis while-left-unfold*)

also have $\dots \leq (y * z) + (y * (y ; x ; (x * 1))) ; ((y ; (x * 1)) * z)$

by (*metis add-right-isotone while-sub-associative*)

also have $\dots \leq (y * z) + (y * (y ; x ; (x * ((y ; (x * 1)) * z))))$

by (*smt add-right-isotone mult-associative mult-right-isotone while-one-mult-below while-right-isotone*)

finally show *?thesis*

qed

lemma *while-mult-zero-add*: $(x + y ; 0) * z = x * ((y ; 0) * z)$

proof –

have $(x + y ; 0) * z = (x * (y ; 0)) * (x * z)$

by (*metis while-sumstar*)

also have $\dots = (x * z) + (x * (y ; 0)) ; ((x * (y ; 0)) * (x * z))$

by (*metis while-left-unfold*)

also have $\dots \leq (x * z) + (x * (y ; 0))$

by (*metis add-right-isotone mult-associative mult-left-zero while-sub-associative*)

also have $\dots = x * ((y ; 0) * z)$

by (*metis add-commutative while-left-dist-add while-zero-2*)

finally show *?thesis*

by (*metis le-neq-trans less-def while-sub-dist-3*)

qed

lemma *while-add-mult-zero*: $(x + y ; 0) * y = x * y$

by (*metis less-eq-def while-mult-zero-add while-zero-2 zero-right-mult-decreasing*)

lemma *while-mult-zero-add-2*: $(x + y ; 0) * z = (x * z) + (x * (y ; 0))$
by (*metis add-commutative while-left-dist-add while-mult-zero-add while-zero-2*)

lemma *while-add-zero-star*: $(x + y ; 0) * z = x * (y ; 0 + z)$
by (*metis while-mult-zero-add while-zero-2*)

lemma *while-unfold-sum*: $(x + y) * z = (x * z) + (x * (y ; ((x + y) * z)))$
apply (*rule antisym*)
apply (*smt add-associative less-eq-def while-absorb-1 while-increasing while-mult-star-exchange while-right-unfold while-sub-associative while-sumstar*)
by (*metis add-least-upper-bound while-decompose-7 while-mult-increasing while-right-isotone while-sub-dist*)

lemma *while-simulate-left*: $x ; z \leq z ; y + w \rightarrow x * (z ; v) \leq z ; (y * v) + (x * (w ; (y * v)))$
by (*smt add-commutative add-right-isotone mult-right-isotone order-trans while-one-increasing while-simulate-left-plus*)

lemma *while-simulate-right*: $z ; x \leq y ; z + w \rightarrow z ; (x * v) \leq y * (z ; v + w ; (x * v))$
proof –
have $y ; z + w \leq y ; (y * z) + w$
by (*metis add-left-isotone mult-right-isotone while-increasing*)
thus *?thesis*
by (*smt order-trans while-simulate-right-plus*)
qed

lemma *while-simulate*: $z ; x \leq y ; z \rightarrow z ; (x * v) \leq y * (z ; v)$
by (*metis add-right-zero mult-left-zero while-simulate-right*)

lemma *while-while-mult*: $1 * (x * y) = (x * 1) * y$
proof –
have $(x * 1) * y \leq (x * 1) ; ((x * 1) * y)$
by (*metis order-refl while-increasing while-sup-one-left-unfold*)
also have $\dots \leq 1 * ((x * 1) ; y)$
by (*metis mult-left-one order-refl while-mult-upper-bound while-simulate*)
also have $\dots \leq 1 * (x * y)$
by (*metis while-one-mult-below while-right-isotone*)
finally show *?thesis*
by (*metis antisym while-sub-dist-3 while-while-add*)
qed

lemma *while-simulate-left-1*: $x ; z \leq z ; y \rightarrow x * (z ; v) \leq z ; (y * v) + (x * 0)$
by (*metis add-right-zero mult-left-zero while-simulate-left*)

lemma *while-associative-1*: $1 \leq z \rightarrow x * (y ; z) = (x * y) ; z$
proof
assume *1*: $1 \leq z$
have $x * (y ; z) \leq x * ((x * y) ; z)$
by (*metis less-eq-def mult-right-dist-add while-plus-one while-right-isotone*)
also have $\dots \leq (x * y) ; (0 * z) + (x * 0)$

by (*metis mult-associative mult-right-subdist-add-right while-left-unfold while-simulate-absorb while-zero*)
 also have $\dots \leq (x * y) ; z + (x * 0) ; z$ using 1
 by (*metis add-least-upper-bound add-left-upper-bound add-right-upper-bound case-split-right while-plus-one while-zero*)
 also have $\dots = (x * y) ; z$
 by (*metis add-right-zero mult-right-dist-add while-left-dist-add*)
 finally show $x * (y ; z) = (x * y) ; z$
 by (*metis antisym while-sub-associative*)
 qed

lemma *while-associative-while-1*: $x * (y ; (z * 1)) = (x * y) ; (z * 1)$
 by (*metis while-associative-1 while-increasing*)

lemma *while-one-while*: $(x * 1) ; (y * 1) = x * (y * 1)$
 by (*metis mult-left-one while-associative-while-1*)

lemma *while-decompose-5-below*: $(x * (y * 1)) * z \leq (y * (x * 1)) * z$
 by (*smt add-commutative mult-left-dist-add mult-right-one while-increasing while-left-unfold while-mult-star-exchange while-one-while while-plus-one while-sumstar*)

lemma *while-decompose-5*: $(x * (y * 1)) * z = (y * (x * 1)) * z$
 by (*metis antisym while-decompose-5-below*)

lemma *while-decompose-4*: $(x * (y * 1)) * z = x * ((y * (x * 1)) * z)$
 by (*metis while-decompose-5 while-decompose-9 while-transitive*)

lemma *while-simulate-2*: $y ; (x * 1) \leq x * (y * 1) \leftrightarrow y * (x * 1) \leq x * (y * 1)$
proof (*rule iffI*)

assume $y ; (x * 1) \leq x * (y * 1)$
 hence $y ; (x * 1) \leq (x * 1) ; (y * 1)$
 by (*metis while-one-while*)
 hence $y * ((x * 1) ; 1) \leq (x * 1) ; (y * 1) + (y * 0)$
 by (*metis while-simulate-left-plus-1*)
 hence $y * (x * 1) \leq (x * (y * 1)) + (y * 0)$
 by (*metis mult-right-one while-one-while*)
 also have $\dots = x * (y * 1)$
 by (*metis add-commutative less-eq-def order-trans while-increasing while-right-isotone zero-least*)
 finally show $y * (x * 1) \leq x * (y * 1)$

next
 assume $y * (x * 1) \leq x * (y * 1)$
 thus $y ; (x * 1) \leq x * (y * 1)$
 by (*metis order-trans while-mult-increasing*)

qed

lemma *while-simulate-1*: $y ; x \leq x ; y \rightarrow y * (x * 1) \leq x * (y * 1)$
 by (*smt mult-right-one order-trans while-one-increasing while-right-isotone while-simulate while-simulate-2*)

lemma *while-simulate-3*: $y ; (x * 1) \leq x * 1 \rightarrow y * (x * 1) \leq x * (y * 1)$

by (*metis add-idempotent case-split-right while-increasing while-mult-upper-bound while-simulate-2*)

lemma *while-extra-while*: $(y ; (x * 1)) * z = (y ; (y * (x * 1))) * z$

proof –

have $y ; (y * (x * 1)) \leq y ; (x * 1) ; (y ; (x * 1) * 1)$

by (*smt add-commutative add-left-upper-bound mult-right-one order-trans while-back-loop-prefixpoint while-left-isotone while-mult-star-exchange*)

hence 1: $(y ; (y * (x * 1))) * z \leq (y ; (x * 1)) * z$

by (*metis while-simulate-right-plus-1 mult-left-one*)

have $(y ; (x * 1)) * z \leq (y ; (y * (x * 1))) * z$

by (*metis while-increasing while-left-isotone while-mult-star-exchange*)

thus *?thesis* **using** 1

by (*metis antisym*)

qed

lemma *while-separate-4*: $y ; x \leq x ; (x * (1 + y)) \rightarrow (x + y) * z = x * (y * z)$

proof

assume 1: $y ; x \leq x ; (x * (1 + y))$

hence $(1 + y) ; x \leq x ; (x * (1 + y))$

by (*smt add-associative add-least-upper-bound mult-left-one mult-left-subdist-add-left mult-right-dist-add mult-right-one while-left-unfold*)

hence 2: $(1 + y) ; (x * 1) \leq x * (1 + y)$

by (*metis mult-right-one while-simulate-right-plus-1*)

have $y ; x ; (x * 1) \leq x ; (x * ((1 + y) ; (x * 1)))$ **using** 1

by (*smt less-eq-def mult-associative mult-right-dist-add while-associative-1 while-increasing*)

also have $\dots \leq x ; (x * (1 + y))$ **using** 2

by (*smt less-eq-def mult-left-dist-add order-refl while-absorb-2 while-left-dist-add*)

also have $\dots \leq x ; (x * 1) ; (y * 1)$

by (*metis add-least-upper-bound mult-associative mult-right-isotone while-increasing while-one-increasing while-one-while while-right-isotone*)

finally have $y * (x ; (x * 1)) \leq x ; (x * 1) ; (y * 1) + (y * 0)$

by (*metis mult-associative mult-right-one while-simulate-left-plus-1*)

hence $(y * 1) ; (y * x) \leq x ; (x * y * 1) + (y * 0)$

by (*smt less-eq-def mult-associative mult-right-one order-refl order-trans while-absorb-2 while-left-dist-add while-mult-star-exchange while-one-mult-below while-one-while while-plus-one*)

hence $(y * 1) ; ((y * x) * (y * z)) \leq x * ((y * 1) ; (y * z) + (y * 0) ; ((y * x) * (y * z)))$

by (*metis while-simulate-right-plus*)

also have $\dots \leq x * ((y * z) + (y * 0))$

by (*metis add-isotone mult-left-zero order-refl while-absorb-2 while-one-mult-below while-right-isotone while-sub-associative*)

also have $\dots = x * y * z$

by (*metis add-right-zero while-left-dist-add*)

finally show $(x + y) * z = x * (y * z)$

by (*smt add-commutative less-eq-def mult-left-one mult-right-dist-add while-plus-one while-sub-associative while-sumstar*)

qed

lemma *while-separate-5*: $y ; x \leq x ; (x * (x + y)) \rightarrow (x + y) * z = x * (y * z)$

by (*smt order-trans while-add-sub-add-one-mult while-separate-4*)

lemma *while-separate-6*: $y ; x \leq x ; (x + y) \rightarrow (x + y) * z = x * (y * z)$

by (*smt order-trans while-increasing while-mult-star-exchange while-separate-5*)

lemma *while-separate-1*: $y ; x \leq x ; y \rightarrow (x + y) * z = x * (y * z)$
by (*metis add-least-upper-bound less-eq-def mult-left-subdist-add-right while-separate-6*)

lemma *while-separate-mult-1*: $y ; x \leq x ; y \rightarrow (x ; y) * z \leq x * (y * z)$
by (*metis while-mult-sub-add while-separate-1*)

lemma *separation*: $y ; x \leq x ; (y * 1) \rightarrow (x + y) * z = x * (y * z)$

proof

assume $y ; x \leq x ; (y * 1)$

hence $y * x \leq x ; (y * 1) + (y * 0)$

by (*metis mult-right-one while-simulate-left-plus-1*)

also have $\dots \leq x ; (x * y * 1) + (y * 0)$

by (*metis add-left-isotone while-increasing while-mult-star-exchange*)

finally have $(y * 1) ; (y * x) \leq x ; (x * y * 1) + (y * 0)$

by (*metis order-refl order-trans while-absorb-2 while-one-mult-below*)

hence $(y * 1) ; ((y * x) * (y * z)) \leq x * ((y * 1) ; (y * z) + (y * 0)) ; ((y * x) * (y * z))$

by (*metis while-simulate-right-plus*)

also have $\dots \leq x * ((y * z) + (y * 0))$

by (*metis add-isotone mult-left-zero order-refl while-absorb-2 while-one-mult-below while-right-isotone while-sub-associative*)

also have $\dots = x * y * z$

by (*metis add-right-zero while-left-dist-add*)

finally show $(x + y) * z = x * (y * z)$

by (*smt add-commutative less-eq-def mult-left-one mult-right-dist-add while-plus-one while-sub-associative while-sumstar*)

qed

lemma *while-separate-left*: $y ; x \leq x ; (y * 1) \rightarrow y * (x * z) \leq x * (y * z)$
by (*metis add-commutative separation while-sub-dist-3*)

lemma *while-simulate-4*: $y ; x \leq x ; (x * (1 + y)) \rightarrow y * (x * z) \leq x * (y * z)$
by (*metis add-commutative while-separate-4 while-sub-dist-3*)

lemma *while-simulate-5*: $y ; x \leq x ; (x * (x + y)) \rightarrow y * (x * z) \leq x * (y * z)$
by (*smt order-trans while-add-sub-add-one-mult while-simulate-4*)

lemma *while-simulate-6*: $y ; x \leq x ; (x + y) \rightarrow y * (x * z) \leq x * (y * z)$
by (*smt order-trans while-increasing while-mult-star-exchange while-simulate-5*)

lemma *while-simulate-7*: $y ; x \leq x ; y \rightarrow y * (x * z) \leq x * (y * z)$
by (*metis add-commutative mult-left-subdist-add-left order-trans while-simulate-6*)

lemma *while-while-mult-1*: $x * (1 * y) = 1 * (x * y)$
by (*metis add-commutative mult-left-one mult-right-one order-refl while-separate-1*)

lemma *while-while-mult-2*: $x * (1 * y) = (x * 1) * y$
by (*metis while-while-mult while-while-mult-1*)

lemma *while-import*: $p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p \rightarrow p ; (x * y) = p ; ((p ; x) * y)$

proof

assume $1: p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p$
hence $p ; (x * y) \leq (p ; x) * (p ; y)$
by (*smt add-commutative less-eq-def mult-associative mult-left-dist-add mult-right-one while-simulate*)
also have $\dots \leq (p ; x) * y$ **using** 1
by (*metis less-eq-def mult-left-one mult-right-dist-add while-right-isotone*)
finally have $2: p ; (x * y) \leq p ; ((p ; x) * y)$ **using** 1
by (*smt add-commutative less-eq-def mult-associative mult-left-dist-add mult-right-one*)
have $p ; ((p ; x) * y) \leq p ; (x * y)$ **using** 1
by (*metis mult-left-isotone mult-left-one mult-right-isotone while-left-isotone*)
thus $p ; (x * y) = p ; ((p ; x) * y)$ **using** 2
by (*metis antisym*)
qed

lemma *while-preserve*: $p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p \rightarrow p ; (x * y) = p ; (x * (p ; y))$
apply (*rule, rule antisym*)
apply (*smt mult-associative mult-left-isotone mult-right-isotone order-trans while-simulate*)
by (*metis mult-left-isotone mult-left-one mult-right-isotone while-right-isotone*)

lemma *while-plus-below-while*: $(x * 1) ; x \leq x * 1$
by (*metis order-refl while-mult-upper-bound while-one-increasing*)

lemma *while-01*: $(w ; (x * 1)) * (y ; z) \leq (x * w) * ((x * y) ; z)$

proof –

have $(w ; (x * 1)) * (y ; z) = y ; z + w ; (((x * 1) ; w) * ((x * 1) ; y ; z))$
by (*metis mult-associative while-productstar*)
also have $\dots \leq y ; z + w ; ((x * w) * ((x * y) ; z))$
by (*metis add-right-isotone mult-left-isotone mult-right-isotone while-isotone while-one-mult-below*)
also have $\dots \leq (x * y) ; z + (x * w) ; ((x * w) * ((x * y) ; z))$
by (*metis add-isotone mult-right-subdist-add-left while-left-unfold*)
finally show *?thesis*
by (*metis while-left-unfold*)

qed

lemma *while-while-sub-associative*: $x * (y * z) \leq ((x * y) * z) + (x * z)$

proof –

have $1: x ; (x * 1) \leq (x * 1) ; ((x * y) * 1)$
by (*metis add-least-upper-bound order-trans while-back-loop-prefixpoint while-left-plus-below*)
have $x * (y * z) \leq x * ((x * 1) ; (y * z))$
by (*metis mult-left-isotone mult-left-one while-increasing while-right-isotone*)
also have $\dots \leq (x * 1) ; ((x * y) * (y * z)) + (x * 0)$ **using** 1
by (*metis while-simulate-left-plus-1*)
also have $\dots \leq (x * 1) ; ((x * y) * z) + (x * z)$
by (*metis add-isotone order-refl while-absorb-2 while-increasing while-right-isotone zero-least*)
also have $\dots = (x * 1) ; z + (x * 1) ; (x * y) ; ((x * y) * z) + (x * z)$
by (*metis mult-associative mult-left-dist-add while-left-unfold*)
also have $\dots = (x * y) ; ((x * y) * z) + (x * z)$

```

    by (smt add-associative add-commutative less-eq-def mult-left-one mult-right-dist-add order-refl while-absorb-1 while-plus-one while-sub-associative)
  also have ... ≤ ((x * y) * z) + (x * z)
    by (metis add-left-isotone while-left-plus-below)
  finally show ?thesis
  .
qed

end

class binary-itering-T = semiring-T + binary-itering

begin

lemma while-right-top: x * T = T
  by (metis add-left-top while-left-unfold)

lemma while-left-top: T ; (x * 1) = T
  by (metis add-right-top antisym top-greatest while-back-loop-prefixpoint)

end

class binary-itering-1 = binary-itering +
  assumes while-denest-0: w ; (x * (y ; z)) ≤ (w ; (x * y)) * (w ; (x * y) ; z)

begin

lemma while-denest-1: w ; (x * (y ; z)) ≤ (w ; (x * y)) * z
  by (metis order-trans while-denest-0 while-right-plus-below)

lemma while-mult-sub-while-while: x * (y ; z) ≤ (x * y) * z
  by (metis mult-left-one while-denest-1)

lemma while-zero-zero: (x * 0) * 0 = x * 0
  by (smt less-eq-def mult-left-zero while-left-dist-add while-mult-star-exchange while-mult-sub-while-while while-mult-zero-add-2 while-plus-one while-sumstar)

lemma while-mult-zero-zero: (x ; (y * 0)) * 0 = x ; (y * 0)
  apply (rule antisym)
  apply (metis add-least-upper-bound add-right-zero mult-left-zero mult-right-isotone while-left-dist-add while-slide while-sub-associative)
  by (metis mult-left-zero while-denest-1)

lemma while-denest-2: w ; ((x * (y ; w)) * z) = w ; (((x * y) ; w) * z)
  apply (rule antisym)
  apply (metis mult-associative while-denest-0 while-simulate-right-plus-1 while-slide)

```

by (*metis mult-right-isotone while-left-isotone while-sub-associative*)

lemma *while-denest-3*: $(x * w) * (x * 0) = (x * w) * 0$

by (*metis while-absorb-2 while-right-isotone while-zero-zero zero-least*)

lemma *while-02*: $x * ((x * w) * ((x * y) ; z)) = (x * w) * ((x * y) ; z)$

proof –

have $x ; ((x * w) * ((x * y) ; z)) = x ; (x * y) ; z + x ; (x * w) ; ((x * w) * ((x * y) ; z))$

by (*metis mult-associative mult-left-dist-add while-left-unfold*)

also have $\dots \leq (x * w) * ((x * y) ; z)$

by (*metis add-isotone mult-right-subdist-add-right while-left-unfold*)

finally have $x * ((x * w) * ((x * y) ; z)) \leq ((x * w) * ((x * y) ; z)) + (x * 0)$

by (*metis while-simulate-absorb*)

hence $x * ((x * w) * ((x * y) ; z)) \leq (x * w) * ((x * y) ; z)$

by (*smt add-least-upper-bound order-refl order-trans while-mult-sub-while-while while-right-isotone zero-least*)

thus *?thesis*

by (*metis antisym while-increasing*)

qed

lemma *while-sumstar-3-below*: $(x * y) * (x * z) \leq (x * y) * ((x * 1) ; z)$

proof –

have $(x * y) * (x * z) = (x * z) + ((x * y) * ((x * y) ; (x * z)))$

by (*metis while-right-unfold*)

also have $\dots \leq (x * z) + ((x * y) * (x * (y ; (x * z))))$

by (*metis add-right-isotone while-right-isotone while-sub-associative*)

also have $\dots \leq (x * z) + ((x * y) * (x * ((x * y) * ((x * 1) ; z))))$

by (*smt add-right-isotone mult-left-isotone mult-left-one order-trans while-increasing while-mult-upper-bound while-one-increasing while-right-isotone while-sumstar while-transitive*)

also have $\dots = (x * z) + ((x * y) * ((x * 1) ; z))$

by (*metis while-02 while-transitive*)

also have $\dots = (x * y) * ((x * 1) ; z)$

by (*smt add-associative mult-left-one mult-right-dist-add while-02 while-left-dist-add while-plus-one*)

finally show *?thesis*

.

qed

lemma *while-sumstar-4-below*: $(x * y) * ((x * 1) ; z) \leq x * ((y ; (x * 1)) * z)$

proof –

have $(x * y) * ((x * 1) ; z) = (x * 1) ; z + (x * y) ; ((x * y) * ((x * 1) ; z))$

by (*metis while-left-unfold*)

also have $\dots \leq (x * z) + (x * (y ; ((x * y) * ((x * 1) ; z))))$

by (*metis add-isotone while-one-mult-below while-sub-associative*)

also have $\dots = (x * z) + (x * (y ; (((x * 1) ; y) * ((x * 1) ; z))))$

by (*metis mult-left-one while-denest-2*)

also have $\dots = x * ((y ; (x * 1)) * z)$

by (*metis while-left-dist-add while-productstar*)

finally show *?thesis*

.

qed

lemma *while-sumstar-1*: $(x + y) * z = (x * y) * ((x * 1) ; z)$

by (*smt eq-iff order-trans while-add-1-below while-sumstar while-sumstar-3-below while-sumstar-4-below*)

lemma *while-sumstar-2*: $(x + y) * z = x * ((y ; (x * 1)) * z)$

by (*metis eq-iff while-add-1-below while-sumstar-1 while-sumstar-4-below*)

lemma *while-sumstar-3*: $(x + y) * z = ((x * 1) ; y) * (x * z)$

by (*metis eq-iff while-sumstar while-sumstar-1-below while-sumstar-2 while-sumstar-2-below*)

lemma *while-decompose-6*: $x * ((y ; (x * 1)) * z) = y * ((x ; (y * 1)) * z)$

by (*metis add-commutative while-sumstar-2*)

lemma *while-denest-4*: $(x * w) * (x * (y ; z)) = (x * w) * ((x * y) ; z)$

proof –

have $(x * w) * (x * (y ; z)) = x * ((w ; (x * 1)) * (y ; z))$

by (*metis while-sumstar while-sumstar-2*)

also have $\dots \leq (x * w) * ((x * y) ; z)$

by (*smt antisym while-01 while-02 while-increasing while-right-isotone*)

finally show *?thesis*

by (*metis antisym while-right-isotone while-sub-associative*)

qed

lemma *while-denest-5*: $w ; ((x * (y ; w)) * (x * (y ; z))) = w ; (((x * y) ; w) * ((x * y) ; z))$

by (*metis while-denest-2 while-denest-4*)

lemma *while-denest-6*: $(w ; (x * y)) * z = z + w ; ((x + y ; w) * (y ; z))$

by (*metis while-denest-5 while-productstar while-sumstar*)

lemma *while-sum-below-one*: $y ; ((x + y) * z) \leq (y ; (x * 1)) * z$

by (*metis add-right-divisibility mult-left-one while-denest-6*)

lemma *while-separate-unfold*: $(y ; (x * 1)) * z = (y * z) + (y * (y ; x ; (x * ((y ; (x * 1)) * z))))$

proof –

have $y * (y ; x ; (x * ((y ; (x * 1)) * z))) \leq y * (y ; ((x + y) * z))$

by (*metis mult-associative mult-right-isotone while-sumstar-2 while-left-plus-below while-right-isotone*)

also have $\dots \leq (y ; (x * 1)) * z$

by (*metis add-commutative add-left-upper-bound while-absorb-1 while-mult-star-exchange while-sum-below-one*)

finally have $(y * z) + (y * (y ; x ; (x * ((y ; (x * 1)) * z)))) \leq (y ; (x * 1)) * z$

by (*metis add-least-upper-bound mult-left-subdist-add-left mult-right-one while-left-isotone while-left-unfold*)

thus *?thesis*

by (*metis antisym while-separate-unfold-below*)

qed

lemma *while-finite-associative*: $x * 0 = 0 \rightarrow (x * y) ; z = x * (y ; z)$

by (*metis while-denest-4 while-zero*)

lemma *atomicity-refinement*: $s = s ; q \wedge x = q ; x \wedge q ; b = 0 \wedge r ; b \leq b ; r \wedge r ; l \leq l ; r \wedge x ; l \leq l ; x \wedge b ; l \leq l ; b \wedge q ; l \leq l ; q \wedge r * q \leq q ; (r * 1) \wedge q \leq 1 \rightarrow s ; ((x + b + r + l) * (q ; z)) \leq s ; ((x ; (b * q) + r + l) * z)$

proof

assume 1: $s = s ; q \wedge x = q ; x \wedge q ; b = 0 \wedge r ; b \leq b ; r \wedge r ; l \leq l ; r \wedge x ; l \leq l ; x \wedge b ; l \leq l ; b \wedge q ; l \leq l ; q \wedge r * q \leq q ; (r * 1) \wedge q \leq 1$

hence 2: $(x + b + r) ; l \leq l ; (x + b + r)$

by (*smt add-commutative add-least-upper-bound mult-left-subdist-add-right mult-right-dist-add order-trans*)

have $q ; ((x ; (b * r * 1) ; q) * z) \leq (x ; (b * r * 1) ; q) * z$ **using** 1

by (*smt eq-refl order-trans while-increasing while-mult-upper-bound*)

also have $\dots \leq (x ; (b * ((r * 1) ; q))) * z$

by (*metis mult-associative mult-right-isotone while-left-isotone while-sub-associative*)

also have $\dots \leq (x ; (b * r * q)) * z$

by (*metis mult-right-isotone while-left-isotone while-one-mult-below while-right-isotone*)

also have $\dots \leq (x ; (b * (q ; (r * 1)))) * z$ **using** 1

by (*metis mult-right-isotone while-left-isotone while-right-isotone*)

finally have 3: $q ; ((x ; (b * r * 1) ; q) * z) \leq (x ; (b * q) ; (r * 1)) * z$

by (*metis mult-associative while-associative-while-1*)

have $s ; ((x + b + r + l) * (q ; z)) = s ; (l * (x + b + r) * (q ; z))$ **using** 2

by (*metis add-commutative while-separate-1*)

also have $\dots = s ; q ; (l * b * r * (q ; x ; (b * r * 1)) * (q ; z))$ **using** 1

by (*smt add-associative add-commutative while-sumstar-2 while-separate-1*)

also have $\dots = s ; q ; (l * b * r * (q ; ((x ; (b * r * 1) ; q) * z)))$

by (*smt mult-associative while-slide*)

also have $\dots \leq s ; q ; (l * b * r * (x ; (b * q) ; (r * 1)) * z)$ **using** 3

by (*metis mult-right-isotone while-right-isotone*)

also have $\dots \leq s ; (l * q ; (b * r * (x ; (b * q) ; (r * 1)) * z))$ **using** 1

by (*smt mult-associative mult-right-isotone while-simulate*)

also have $\dots = s ; (l * q ; (r * (x ; (b * q) ; (r * 1)) * z))$ **using** 1

by (*metis while-elimination*)

also have $\dots \leq s ; (l * r * (x ; (b * q) ; (r * 1)) * z)$ **using** 1

by (*smt add-left-divisibility mult-left-one mult-right-dist-add mult-right-isotone while-right-isotone*)

also have $\dots = s ; (l * (r + x ; (b * q)) * z)$

by (*metis while-sumstar-2*)

also have $\dots \leq s ; ((x ; (b * q) + r + l) * z)$

by (*metis add-commutative mult-right-isotone while-sub-dist-3*)

finally show $s ; ((x + b + r + l) * (q ; z)) \leq s ; ((x ; (b * q) + r + l) * z)$

qed

end

class *binary-itering-1-T* = *binary-itering-T* + *binary-itering-1*

begin

end

class *strict-itering* = *itering-3* + *while* +
 assumes *while-def*: $x * y = x^\circ ; y$

begin

subclass *binary-itering-1*
 apply *unfold-locales*
 apply (*metis add-commutative circ-loop-fixpoint circ-slide mult-associative while-def*)
 apply (*metis circ-add mult-associative while-def*)
 apply (*metis mult-left-dist-add while-def*)
 apply (*metis mult-associative order-refl while-def*)
 apply (*metis circ-simulate-left-plus mult-associative mult-left-isotone mult-right-dist-add mult-right-one while-def*)
 apply (*metis circ-simulate-right-plus mult-associative mult-left-isotone mult-right-dist-add while-def*)
 by (*metis add-right-divisibility circ-loop-fixpoint mult-associative while-def*)

lemma *while-associative*: $(x * y) ; z = x * (y ; z)$
 by (*metis mult-associative while-def*)

lemma *circ-import*: $p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p \rightarrow p ; x^\circ = p ; (p ; x)^\circ$
proof

assume 1: $p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p$
 hence $p ; x \leq p ; x ; p$
 by (*smt mult-associative mult-left-isotone mult-right-isotone order-trans*)
 hence $p ; x^\circ \leq (p ; x)^\circ$ using 1
 by (*smt circ-mult-upper-bound circ-reflexive circ-simulate order-refl order-trans*)
 hence 2: $p ; x^\circ \leq p ; (p ; x)^\circ$ using 1
 by (*smt mult-associative mult-left-isotone mult-right-isotone order-trans*)
 have $p ; (p ; x)^\circ \leq p ; x^\circ$ using 1
 by (*metis circ-isotone mult-left-isotone mult-left-one mult-right-isotone*)
 thus $p ; x^\circ = p ; (p ; x)^\circ$ using 2
 by (*metis antisym*)
qed

lemma *while-one-mult*: $(x * 1) ; x = x * x$
 by (*metis mult-right-one while-def*)

lemma *while-back-loop-is-fixpoint*: *is-fixpoint* $(\lambda x . x ; y + z)$ $(z ; (y * 1))$
 by (*metis circ-back-loop-is-fixpoint mult-right-one while-def*)

lemma $x ; z \leq z \wedge y \leq z \wedge x * 1 \leq z \rightarrow x * y \leq z$
 by (*metis add-commutative less-eq-def mult-left-one order-trans while-right-isotone while-simulate-absorb zero-right-mult-decreasing*)

lemma $(x + y) * z = ((x * 1) ; y) * ((x * 1) ; z)$

by (*metis mult-right-one while-def while-sumstar*)

lemma $(x * 1) ; y = x * y$

by (*metis mult-left-one while-associative*)

end

class *strict-itering-T* = *itering-T* + *strict-itering*

begin

subclass *binary-itering-1-T* ..

lemma *while-top-2*: $T * z = T ; z$

by (*metis circ-top while-def*)

lemma *while-mult-top-2*: $(x ; T) * z = z + x ; T ; z$

by (*metis circ-left-top mult-associative while-def while-left-unfold*)

end

class *nonstrict-itering* = *omega-algebra-T* + *while* +

assumes *while-def*: $x * y = x^\omega + x^* ; y$

begin

subclass *binary-itering-T*

proof (*unfold-locales*)

fix $x y z$

show $(x ; y) * z = z + x ; ((y ; x) * (y ; z))$

by (*metis add-commutative mult-associative mult-left-dist-add omega-loop-fixpoint omega-slide star.circ-slide while-def*)

next

fix $x y z$

show $(x + y) * z = (x * y) * (x * z)$

proof –

have 1: $(x + y) * z = (x^* ; y)^\omega + (x^* ; y)^* ; (x^\omega + x^* ; z)$

by (*smt add-associative mult-associative mult-left-dist-add omega-decompose star.circ-add while-def*)

hence 2: $(x + y) * z \leq (x * y) * (x * z)$

by (*smt add-isotone add-right-upper-bound less-eq-def mult-left-isotone omega-sub-dist star.circ-sub-dist while-def*)

let $?rhs = x^* ; y ; ((x^\omega + x^* ; y)^\omega + (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z)) + (x^\omega + x^* ; z)$

have $x^\omega ; (x^\omega + x^* ; y)^\omega \leq x^\omega$

```

  by (metis omega-sub-vector)
hence  $x^\omega ; (x^\omega + x^* ; y)^\omega + x^* ; y ; (x^\omega + x^* ; y)^\omega \leq ?rhs$ 
  by (smt add-commutative add-isotone add-left-upper-bound mult-left-dist-add order-trans)
hence 3:  $(x^\omega + x^* ; y)^\omega \leq ?rhs$ 
  by (metis mult-right-dist-add omega-unfold)
have  $x^\omega ; (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) \leq x^\omega$ 
  by (metis mult-associative omega-sub-vector)
hence  $x^\omega ; (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) + x^* ; y ; (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) \leq ?rhs$ 
  by (smt add-commutative add-isotone add-right-upper-bound mult-associative mult-left-dist-add order-trans)
hence  $(x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) \leq ?rhs$ 
  by (smt add-associative add-right-upper-bound less-eq-def mult-associative mult-right-dist-add star.circ-loop-fixpoint)
hence  $(x^\omega + x^* ; y)^\omega + (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) \leq ?rhs$  using 3
  by (metis add-least-upper-bound)
hence  $(x^\omega + x^* ; y)^\omega + (x^\omega + x^* ; y)^* ; (x^\omega + x^* ; z) \leq (x^* ; y)^\omega + (x^* ; y)^* ; (x^\omega + x^* ; z)$ 
  by (metis add-commutative omega-induct)
thus ?thesis using 1 2
  by (smt antisym while-def)
qed
next
fix x y z
show  $x * (y + z) = (x * y) + (x * z)$ 
  by (smt add-associative add-commutative add-left-upper-bound less-eq-def mult-left-dist-add while-def)
next
fix x y z
show  $(x * y) ; z \leq x * (y ; z)$ 
  by (metis mult-associative mult-right-dist-add omega-loop-fixpoint omega-loop-greatest-fixpoint while-def)
next
fix v w x y z
show  $x ; z \leq z ; (y * 1) + w \rightarrow x * (z ; v) \leq z ; (y * v) + (x * (w ; (y * v)))$ 
proof
  assume  $x ; z \leq z ; (y * 1) + w$ 
  hence 1:  $x ; z \leq z ; y^\omega + z ; y^* + w$ 
    by (metis mult-left-dist-add mult-right-one while-def)
  let ?rhs =  $z ; (y^\omega + y^* ; v) + x^\omega + x^* ; w ; (y^\omega + y^* ; v)$ 
  have 2:  $z ; v \leq ?rhs$ 
    by (metis add-least-upper-bound add-left-upper-bound mult-left-dist-add omega-loop-fixpoint)
  have  $x ; z ; (y^\omega + y^* ; v) \leq ?rhs$ 
proof -
  have  $x ; z ; (y^\omega + y^* ; v) \leq (z ; y^\omega + z ; y^* + w) ; (y^\omega + y^* ; v)$  using 1
    by (metis mult-left-isotone)
  also have  $\dots = z ; (y^\omega ; (y^\omega + y^* ; v) + y^* ; (y^\omega + y^* ; v)) + w ; (y^\omega + y^* ; v)$ 
    by (smt mult-associative mult-left-dist-add mult-right-dist-add)
  also have  $\dots = z ; (y^\omega ; (y^\omega + y^* ; v) + y^\omega + y^* ; v) + w ; (y^\omega + y^* ; v)$ 
    by (smt add-associative mult-associative mult-left-dist-add star.circ-transitive-equal star-mult-omega)
  also have  $\dots \leq z ; (y^\omega + y^* ; v) + x^* ; w ; (y^\omega + y^* ; v)$ 
    by (smt add-commutative add-isotone add-left-top mult-left-dist-add mult-left-one mult-right-dist-add mult-right-subdist-add-left omega-vector order-refl star.circ-plus-one)
  finally show ?thesis

```

```

    by (smt add-associative add-commutative less-eq-def)
qed
hence  $x ; ?rhs \leq ?rhs$ 
by (smt add-associative add-commutative add-left-upper-bound less-eq-def mult-associative mult-left-dist-add mult-right-dist-add omega-unfold star.circ-increasing star.circ-transitive-equal)
hence  $z ; v + x ; ?rhs \leq ?rhs$  using 2
    by (metis add-least-upper-bound)
hence  $x^* ; z ; v \leq ?rhs$ 
    by (metis mult-associative star-left-induct)
hence  $x^\omega + x^* ; z ; v \leq ?rhs$ 
    by (metis add-least-upper-bound add-left-upper-bound)
thus  $x * (z ; v) \leq z ; (y * v) + (x * (w ; (y * v)))$ 
    by (smt add-associative mult-associative mult-left-dist-add while-def)
qed
next
fix  $v w x y z$ 
show  $z ; x \leq y ; (y * z) + w \rightarrow z ; (x * v) \leq y * (z ; v + w ; (x * v))$ 
proof
    assume  $z ; x \leq y ; (y * z) + w$ 
    hence  $z ; x \leq y ; (y^\omega + y^* ; z) + w$ 
        by (metis while-def)
    hence 1:  $z ; x \leq y^\omega + y ; y^* ; z + w$ 
        by (metis mult-associative mult-left-dist-add omega-unfold)
    let  $?rhs = y^\omega + y^* ; z ; v + y^* ; w ; (x^\omega + x^* ; v)$ 
    have 2:  $z ; x^\omega \leq ?rhs$ 
proof -
    have  $z ; x^\omega \leq y ; y^* ; z ; x^\omega + y^\omega ; x^\omega + w ; x^\omega$  using 1
        by (smt add-commutative less-eq-def mult-associative mult-right-dist-add omega-unfold)
    also have  $\dots \leq y ; y^* ; z ; x^\omega + y^\omega + w ; x^\omega$ 
        by (metis add-left-isotone add-right-isotone omega-sub-vector)
    also have  $\dots = y ; y^* ; (z ; x^\omega) + (y^\omega + w ; x^\omega)$ 
        by (metis add-associative mult-associative)
    finally have  $z ; x^\omega \leq (y ; y^*)^\omega + (y ; y^*)^* ; (y^\omega + w ; x^\omega)$ 
        by (metis add-commutative omega-induct)
    also have  $\dots = y^\omega + y^* ; w ; x^\omega$ 
        by (metis left-plus-omega less-eq-def mult-associative mult-left-dist-add mult-left-subdist-add-left star.left-plus-circ star-mult-omega)
    also have  $\dots \leq ?rhs$ 
        by (metis add-isotone add-left-upper-bound mult-left-subdist-add-left)
    finally show  $?thesis$ 
        by metis
qed
let  $?rhs2 = y^\omega + y^* ; z + y^* ; w ; (x^\omega + x^*)$ 
have  $?rhs2 ; x \leq ?rhs2$ 
proof -
    have 3:  $y^\omega ; x \leq ?rhs2$ 
        by (metis add-associative less-eq-def omega-sub-vector)
    have  $y^* ; z ; x \leq y^* ; (y^\omega + y ; y^* ; z + w)$  using 1
        by (metis mult-associative mult-right-isotone)

```

also have $\dots = y^\omega + y^* ; y ; y^* ; z + y^* ; w$
by (*metis mult-associative mult-left-dist-add star-mult-omega*)
also have $\dots = y^\omega + y ; y^* ; z + y^* ; w$
by (*metis mult-associative star.circ-transitive-equal star-simulation-right-equal*)
also have $\dots \leq y^\omega + y^* ; z + y^* ; w$
by (*metis add-left-isotone add-right-isotone mult-left-isotone star.left-plus-below-circ*)
also have $\dots \leq y^\omega + y^* ; z + y^* ; w ; x^*$
by (*metis add-right-isotone add-right-upper-bound star.circ-back-loop-fixpoint*)
finally have 4: $y^* ; z ; x \leq ?rhs2$
by (*smt add-associative add-commutative less-eq-def mult-left-dist-add*)
have $(x^\omega + x^*) ; x \leq x^\omega + x^*$
by (*metis add-isotone mult-right-dist-add omega-sub-vector star.circ-plus-same star.left-plus-below-circ*)
hence $y^* ; w ; (x^\omega + x^*) ; x \leq ?rhs2$
by (*smt add-right-upper-bound mult-associative mult-right-isotone order-trans*)
thus *?thesis* **using** 3 4
by (*smt add-associative less-eq-def mult-right-dist-add*)
qed
hence $z + ?rhs2 ; x \leq ?rhs2$
by (*smt add-commutative add-least-upper-bound add-right-divisibility while-def omega-loop-fixpoint*)
hence 5: $z ; x^* \leq ?rhs2$
by (*metis star-right-induct*)
have $z ; x^* ; v \leq ?rhs$
proof –
have $z ; x^* ; v \leq ?rhs2 ; v$ **using** 5
by (*metis mult-left-isotone*)
also have $\dots = y^\omega ; v + y^* ; z ; v + y^* ; w ; (x^\omega ; v + x^* ; v)$
by (*metis mult-associative mult-right-dist-add*)
also have $\dots \leq y^\omega + y^* ; z ; v + y^* ; w ; (x^\omega ; v + x^* ; v)$
by (*metis add-left-isotone omega-sub-vector*)
also have $\dots \leq ?rhs$
by (*metis add-left-isotone add-right-isotone mult-right-isotone omega-sub-vector*)
finally show *?thesis*
by *metis*
qed
hence $z ; (x^\omega + x^* ; v) \leq ?rhs$ **using** 2
by (*smt add-associative less-eq-def mult-associative mult-left-dist-add*)
thus $z ; (x * v) \leq y * (z ; v + w ; (x * v))$
by (*metis add-associative mult-associative mult-left-dist-add while-def*)
qed
qed

lemma *while-top*: $T * x = T$

by (*metis add-left-top star.circ-top star-omega-top while-def*)

lemma *while-one-top*: $1 * x = T$

by (*metis add-left-top omega-one while-def*)

lemma *while-finite-associative*: $x^\omega = 0 \rightarrow (x * y) ; z = x * (y ; z)$
by (*metis add-left-zero mult-associative while-def*)

lemma *star-below-while*: $x^* ; y \leq x * y$
by (*metis add-right-upper-bound while-def*)

lemma *while-sub-mult-one*: $x ; (1 * y) \leq 1 * x$
by (*metis top-greatest while-one-top*)

lemma *while-while-one*: $y * (x * 1) = y^\omega + y^* ; x^\omega + y^* ; x^*$
by (*metis add-associative mult-left-dist-add mult-right-one while-def*)

lemma *omega-mult-star-2*: $x^\omega ; y^* = x^\omega$
by (*metis add-right-upper-bound antisym omega-sub-vector star.circ-back-loop-fixpoint*)

lemma *star-omega-absorb*: $y^* ; (y^* ; x)^* ; y^\omega = (y^* ; x)^* ; y^\omega$

proof –

have $y^* ; (y^* ; x)^* ; y^\omega = y^* ; y^* ; x ; (y^* ; x)^* ; y^\omega + y^* ; y^\omega$

by (*metis add-commutative mult-associative mult-right-dist-add star.circ-back-loop-fixpoint star.circ-plus-same*)

thus *?thesis*

by (*metis mult-associative star.circ-loop-fixpoint star.circ-transitive-equal star-mult-omega*)

qed

lemma *star-star-absorb*: $y^* ; (y^* ; x)^* ; y^* = (y^* ; x)^* ; y^*$

by (*metis add-commutative mult-associative star.circ-decompose-4 star.circ-slide star-decompose-1 star-decompose-3*)

lemma *while-simulate-4-plus*: $y ; x \leq x ; (x * (1 + y)) \rightarrow y ; x ; x^* \leq x ; (x * (1 + y))$

proof

have $1: x ; (x * (1 + y)) = x^\omega + x ; x^* + x ; x^* ; y$

by (*metis add-associative mult-associative mult-left-dist-add mult-right-one omega-unfold while-def*)

assume $y ; x \leq x ; (x * (1 + y))$

hence $y ; x ; x^* \leq (x^\omega + x ; x^* + x ; x^* ; y) ; x^*$ **using** 1

by (*metis mult-left-isotone*)

also have $\dots = x^\omega ; x^* + x ; x^* ; x^* + x ; x^* ; y ; x^*$

by (*metis mult-right-dist-add*)

also have $\dots = x ; x^* ; (y ; x ; x^*) + x^\omega + x ; x^* + x ; x^* ; y$

by (*smt add-associative add-commutative mult-associative omega-mult-star-2 star.circ-back-loop-fixpoint star.circ-plus-same star.circ-transitive-equal*)

finally have $y ; x ; x^* \leq x ; x^* ; (y ; x ; x^*) + (x^\omega + x ; x^* + x ; x^* ; y)$

by (*metis add-associative*)

hence $y ; x ; x^* \leq (x ; x^*)^\omega + (x ; x^*)^* ; (x^\omega + x ; x^* + x ; x^* ; y)$

by (*metis add-commutative omega-induct*)

also have $\dots = x^\omega + x^* ; (x^\omega + x ; x^* + x ; x^* ; y)$

by (*metis left-plus-omega star.left-plus-circ*)

finally show $y ; x ; x^* \leq x ; (x * (1 + y))$ **using** 1

by (*metis while-def while-mult-star-exchange while-transitive*)

qed

lemma *while-simulate-4-omega*: $y ; x \leq x ; (x * (1 + y)) \rightarrow y ; x^\omega \leq x^\omega$

proof

have 1 : $x ; (x * (1 + y)) = x^\omega + x ; x^* + x ; x^* ; y$

by (*metis add-associative mult-associative mult-left-dist-add mult-right-one omega-unfold while-def*)

assume $y ; x \leq x ; (x * (1 + y))$

hence $y ; x^\omega \leq (x^\omega + x ; x^* + x ; x^* ; y) ; x^\omega$ **using** 1

by (*smt less-eq-def mult-associative mult-right-dist-add omega-unfold*)

also have $\dots = x^\omega ; x^\omega + x ; x^* ; x^\omega + x ; x^* ; y ; x^\omega$

by (*metis mult-right-dist-add*)

also have $\dots = x ; x^* ; (y ; x^\omega) + x^\omega$

by (*metis add-commutative less-eq-def mult-associative omega-sub-vector omega-unfold star-mult-omega*)

finally have $y ; x^\omega \leq x ; x^* ; (y ; x^\omega) + x^\omega$

by *metis*

hence $y ; x^\omega \leq (x ; x^*)^\omega + (x ; x^*)^* ; x^\omega$

by (*metis add-commutative omega-induct*)

thus $y ; x^\omega \leq x^\omega$

by (*metis add-idempotent left-plus-omega star-mult-omega*)

qed

lemma *while-unfold-below*: $x = z + y ; x \rightarrow x \leq y * z$

by (*metis omega-induct-equal while-def*)

lemma *while-unfold-below-1*: $x = y ; x \rightarrow x \leq y * 1$

by (*metis add-right-upper-bound omega-induct while-def*)

lemma *while-square-1*: $x * 1 = (x ; x) * (x + 1)$

by (*metis mult-right-one omega-square star-square-2 while-def*)

lemma *while-absorb-below-one*: $y ; x \leq x \rightarrow y * x \leq 1 * x$

by (*metis top-greatest while-one-top*)

lemma *omega-import*: $p \leq p ; p \wedge p ; x \leq x ; p \rightarrow p ; x^\omega = p ; (p ; x)^\omega$

proof

assume 1 : $p \leq p ; p \wedge p ; x \leq x ; p$

hence $p ; x^\omega \leq p ; (p ; x) ; x^\omega$

by (*metis mult-associative mult-left-isotone omega-unfold*)

also have $\dots \leq p ; x ; p ; x^\omega$ **using** 1

by (*metis mult-associative mult-left-isotone mult-right-isotone*)

finally have $p ; x^\omega \leq (p ; x)^\omega$

by (*metis mult-associative omega-induct-mult*)

hence $p ; x^\omega \leq p ; (p ; x)^\omega$ **using** 1

by (*smt mult-associative mult-left-isotone mult-right-isotone order-trans*)

thus $p ; x^\omega = p ; (p ; x)^\omega$ **using** 1

by (*metis add-left-divisibility antisym mult-right-isotone omega-induct-mult omega-slide omega-sub-dist*)

qed

lemma *star-import*: $p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p \rightarrow p ; x^* = p ; (p ; x)^*$

proof

assume $1: p \leq p ; p \wedge p \leq 1 \wedge p ; x \leq x ; p$
hence $p ; (p ; x)^* ; x \leq p ; (x ; p)^* ; x$
by (*smt add-right-divisibility mult-associative mult-left-isotone mult-left-subdist-add-right star.circ-simulate star.circ-slide*)
hence $p ; (p ; x)^* ; x \leq (p ; x)^*$
by (*metis add-right-divisibility order-trans star.circ-mult*)
hence $2: p ; (p ; x)^* ; x \leq p ; (p ; x)^*$ **using** 1
by (*smt mult-associative mult-left-isotone mult-right-isotone order-trans*)
have $p \leq p ; (p ; x)^*$
by (*metis add-right-divisibility star.circ-back-loop-fixpoint*)
hence $p ; x^* \leq p ; (p ; x)^*$ **using** 2
by (*metis add-least-upper-bound star-right-induct*)
thus $p ; x^* = p ; (p ; x)^*$ **using** 1
by (*metis antisym mult-left-isotone mult-left-one mult-right-isotone star.circ-isotone*)
qed

lemma *while-loop-is-greatest-postfixpoint*: *is-greatest-postfixpoint* $(\lambda x . y ; x + z) (y * z)$

proof –

have $(y * z) \leq (\lambda x . y ; x + z) (y * z)$
by (*metis is-fixpoint-def order-refl while-loop-is-fixpoint*)
thus *?thesis*
by (*smt add-commutative is-greatest-postfixpoint-def omega-induct while-def*)

qed

lemma *while-loop-is-greatest-fixpoint*: *is-greatest-fixpoint* $(\lambda x . y ; x + z) (y * z)$

by (*metis omega-loop-is-greatest-fixpoint while-def*)

lemma $x ; z \leq z \wedge y \leq z \wedge x * 1 \leq z \rightarrow x * y \leq z$

by (*metis add-commutative add-least-upper-bound add-left-zero less-eq-def while-right-isotone while-simulate-absorb*)

end

class *nonstrict-itering-zero* = *nonstrict-itering* +

assumes *mult-right-zero*: $x ; 0 = 0$

begin

lemma *while-finite-associative-2*: $x * 0 = 0 \rightarrow (x * y) ; z = x * (y ; z)$

by (*metis add-left-zero add-right-zero mult-associative mult-right-zero while-def*)

end

class *nonstrict-itering-tarski* = *nonstrict-itering* +

assumes *tarski*: $x \leq x ; T ; x ; T$

begin

lemma *tarski-mult-top-idempotent*: $x ; T = x ; T ; x ; T$

by (*metis add-commutative less-eq-def mult-associative star.circ-back-loop-fixpoint star.circ-left-top tarski top-mult-top*)

lemma *tarski-top-omega-below*: $x ; T \leq (x ; T)^\omega$

by (*metis mult-associative omega-induct-mult order-refl tarski-mult-top-idempotent*)

lemma *tarski-top-omega*: $x ; T = (x ; T)^\omega$

by (*metis antisym mult-top-omega tarski-top-omega-below*)

lemma *tarski-below-top-omega*: $x \leq (x ; T)^\omega$

by (*metis tarski-top-omega top-right-mult-increasing*)

lemma *tarski-mult-omega-omega*: $(x ; y^\omega)^\omega = x ; y^\omega$

by (*metis mult-associative omega-vector tarski-top-omega*)

lemma *tarski-omega-idempotent*: $x^{\omega\omega} = x^\omega$

by (*metis omega-vector tarski-top-omega*)

lemma *while-denest-2a*: $w ; ((x * (y ; w)) * z) = w ; (((x * y) ; w) * z)$

proof –

have $(x^\omega + x^* ; y ; w)^\omega = (x^* ; y ; w)^* ; x^\omega ; (((x^* ; y ; w)^* ; x^\omega)^\omega + ((x^* ; y ; w)^* ; x^\omega)^* ; (x^* ; y ; w)^\omega) + (x^* ; y ; w)^\omega$

by (*metis add-commutative omega-decompose omega-loop-fixpoint*)

also have $\dots \leq (x^* ; y ; w)^* ; x^\omega + (x^* ; y ; w)^\omega$

by (*metis add-left-isotone mult-associative mult-right-isotone omega-sub-vector*)

finally have **1**: $w ; (x^\omega + x^* ; y ; w)^\omega \leq (w ; x^* ; y)^* ; w ; x^\omega + (w ; x^* ; y)^\omega$

by (*smt add-commutative less-eq-def mult-associative mult-left-dist-add while-def while-slide*)

have $(x^\omega + x^* ; y ; w)^* ; z = (x^* ; y ; w)^* ; x^\omega ; ((x^* ; y ; w)^* ; x^\omega)^* ; (x^* ; y ; w)^* ; z + (x^* ; y ; w)^* ; z$

by (*smt add-commutative mult-associative star.circ-add star.circ-loop-fixpoint*)

also have $\dots \leq (x^* ; y ; w)^* ; x^\omega + (x^* ; y ; w)^* ; z$

by (*smt add-commutative add-right-isotone mult-associative mult-right-isotone omega-sub-vector*)

finally have $w ; (x^\omega + x^* ; y ; w)^* ; z \leq (w ; x^* ; y)^* ; w ; x^\omega + (w ; x^* ; y)^* ; w ; z$

by (*metis mult-associative mult-left-dist-add mult-right-isotone star.circ-slide*)

hence $w ; (x^\omega + x^* ; y ; w)^\omega + w ; (x^\omega + x^* ; y ; w)^* ; z \leq (w ; x^* ; y)^* ; (w ; x^\omega)^\omega + (w ; x^* ; y)^\omega + (w ; x^* ; y)^* ; w ; z$ **using** **1**

by (*smt add-associative add-commutative less-eq-def mult-associative tarski-mult-omega-omega*)

also have $\dots \leq (w ; x^\omega + w ; x^* ; y)^* ; (w ; x^\omega + w ; x^* ; y)^\omega + (w ; x^\omega + w ; x^* ; y)^\omega + (w ; x^\omega + w ; x^* ; y)^* ; w ; z$

by (*metis add-isotone add-left-upper-bound add-right-upper-bound mult-isotone mult-left-isotone omega-isotone star.circ-isotone*)

also have $\dots = (w ; x^\omega + w ; x^* ; y)^\omega + (w ; x^\omega + w ; x^* ; y)^* ; w ; z$

by (*metis add-idempotent star-mult-omega*)

finally have $w ; ((x^\omega + x^* ; y ; w)^\omega + (x^\omega + x^* ; y ; w)^* ; z) \leq w ; ((x^\omega + x^* ; y) ; w)^\omega + w ; ((x^\omega + x^* ; y) ; w)^* ; z$

by (*smt mult-associative mult-left-dist-add omega-slide star.circ-slide*)

hence **2**: $w ; ((x * (y ; w)) * z) \leq w ; (((x * y) ; w) * z)$

by (*smt mult-associative mult-left-dist-add while-def while-slide*)

have $w ; (((x * y) ; w) * z) \leq w ; ((x * (y ; w)) * z)$

by (*metis mult-right-isotone while-left-isotone while-sub-associative*)
 thus *?thesis using 2*
 by (*metis antisym*)
 qed

lemma *while-denest-3*: $(x * w) * x^\omega = (x * w)^\omega$

proof –

have 1: $(x * w) * x^\omega = (x * w)^\omega + (x * w)^* ; x^{\omega\omega}$

by (*metis tarski-omega-idempotent while-def*)

also have $\dots \leq (x * w)^\omega + (x * w)^* ; (x^\omega + x^* ; w)^\omega$

by (*metis add-left-upper-bound add-right-isotone mult-right-isotone omega-isotone*)

also have $\dots = (x * w)^\omega$

by (*metis add-idempotent star-mult-omega while-def*)

finally show *?thesis using 1*

by (*metis add-left-upper-bound antisym-conv*)

qed

lemma *while-denest-4a*: $(x * w) * (x * (y ; z)) = (x * w) * ((x * y) ; z)$

proof –

have $(x * w) * (x * (y ; z)) = (x * w)^\omega + ((x * w) * (x^* ; y ; z))$

by (*smt mult-associative while-denest-3 while-def while-left-dist-add*)

also have $\dots \leq (x * w)^\omega + ((x * w) * ((x * y) ; z))$

by (*metis add-right-isotone mult-left-isotone star-below-while while-right-isotone*)

finally have 1: $(x * w) * (x * (y ; z)) \leq (x * w) * ((x * y) ; z)$

by (*smt add-left-upper-bound less-eq-def while-def*)

have $(x * w) * ((x * y) ; z) \leq (x * w) * (x * (y ; z))$

by (*metis while-right-isotone while-sub-associative*)

thus *?thesis using 1*

by (*metis antisym*)

qed

subclass *binary-itering-1-T*

apply *unfold-locales*

by (*smt mult-associative while-denest-2a while-denest-4a while-increasing while-slide*)

lemma *while-mult-top*: $(x ; T) * z = z + x ; T$

proof –

have 1: $z + x ; T \leq (x ; T) * z$

by (*metis add-least-upper-bound while-denest-1 while-increasing while-one-top*)

have $(x ; T) * z = z + x ; T ; ((x ; T) * z)$

by (*metis while-left-unfold*)

also have $\dots \leq z + x ; T$

by (*metis add-right-isotone mult-associative mult-right-isotone top-greatest*)

finally show *?thesis using 1*

by (*metis antisym*)

qed

lemma *tarski-top-omega-below-2*: $x ; T \leq (x ; T) * 0$
by (*metis add-right-divisibility while-mult-top*)

lemma *tarski-top-omega-2*: $x ; T = (x ; T) * 0$
by (*metis add-left-zero while-mult-top*)

lemma *tarski-below-top-omega-2*: $x \leq (x ; T) * 0$
by (*metis tarski-top-omega-2 top-right-mult-increasing*)

end

class *nonstrict-itering-tarski-zero* = *nonstrict-itering-tarski* + *nonstrict-itering-zero*

begin

lemma $1 = (x ; 0) * 1$
by (*metis mult-right-zero while-zero*)

end

context *binary-itering-1-T*

begin

end

end