

Advantages of Shared Data Structures for Sequences of Balanced Parentheses

Simon Gog¹ Johannes Fischer²

¹Institut of Theoretical Computer Science
Ulm University, Germany

²Center for Bioinformatics (ZBIT)
Universität Tübingen, Germany

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Outline

- Basic Definitions
- Our DS for RMQ
- Geary et al.'s DS for balanced parentheses
- Our result
 - Computing RMQs with $2n + o(n)$ bits
 - Computing LCA with $2n + o(n)$ bits
- Experimental study
 - Comparison of RMQ data structures
 - Comparison of CST implementations
- Conclusion

Range Minimum Queries

Definition

Given an array A of n values. A range minimum query (RMQ) $rmq_A(i, j)$ with $i \leq j$ returns index k and $A[k] = \min\{A[\ell] \mid i \leq \ell \leq j\}$.

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Solution

- Preprocess a RMQ data structure R for A
- R answers a RMQ then in constant time
- Two versions of the problem
 - Systematic: R needs A to answer RMQs
 - Non-systematic: R answers RMQs

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Our solution

- Non-systematic
- $2n + o(n)$ bits ($3n$ bits in practice)
- $3n + o(n)$ bits for construction in linear time

Balanced Parentheses Sequences (BPS)

- Sequence S over the alphabet '(' and ')'
- Each prefix of S contains more '('s than ')s
- Fundamental operations on S :
 - $rank_l(S, i)$
 - $select_l(S, i)$
 - $excess(S, i) = rank_l(i) - rank_r(i)$
 - $find_close(S, i)$ and $find_open(S, i)$
 - $enclose(S, i)$

Example

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$rank_l(S, 5) = 4$

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$select_l(S, 2) = 1$

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$find_close(S, 3) = 20$ and $find_open(S, 20) = 3$

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$enclose(S, 4) = 3$

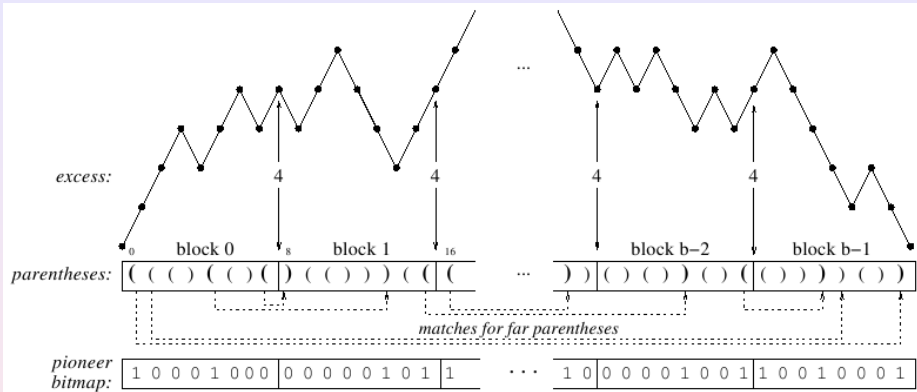
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Time and space

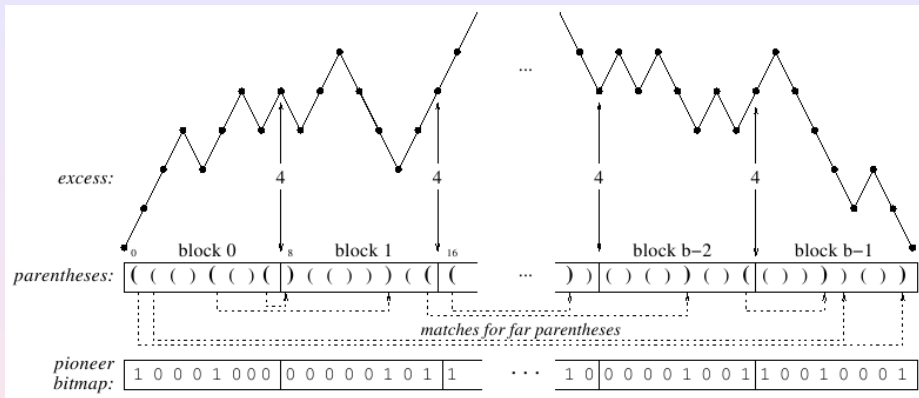
Geary et al.'s data structure of size $o(n)$ supports all operations in constant time

Geary et al.'s support structure for BP



- Partition BPS into b blocks of length $\mathcal{O}(\log n)$
- Calculate far parentheses/pioneers to answer *find_close*, *find_open*, *enclose*
- pioneer bitmap takes $\mathcal{O}\left(\frac{n \log \log n}{\log n}\right)$ bits

Geary et al.'s support structure for BP



- set of pioneers is again a BPS (of length $n_1 < 4b - 6$)
- recursively build data structure for pioneers
- store answers explicitly on the second level

The range restricted enclose method

Construction of the BP of the SCT

Example

i	0	1	2	3	4	5	6	7	8	9	10
A[i]	-1	2	1	3	1	2	0	2	0	1	-1

{
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Construction of the BP of the SCT

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$$\left(\begin{array}{c} (\\ -1 \end{array} \right) \begin{array}{c} (\\ 2 \end{array} \begin{array}{c}) \\ 2 \end{array}$$

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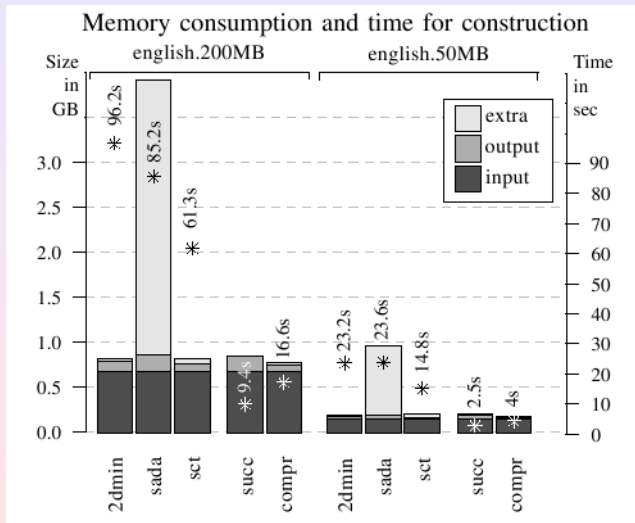
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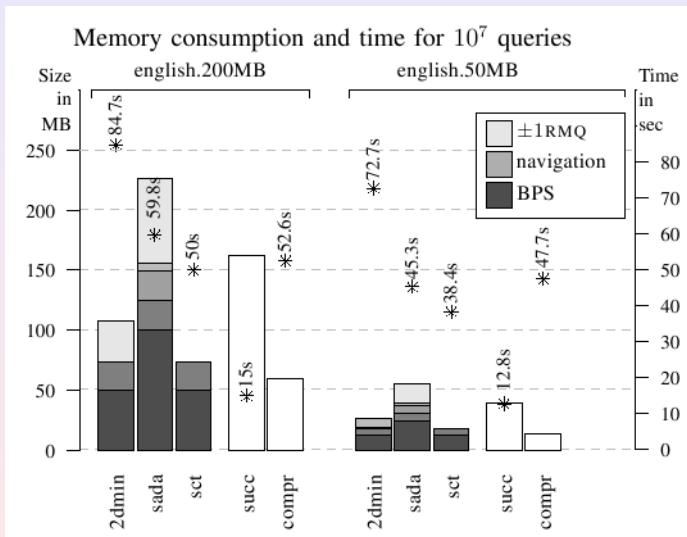
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 $\left(\begin{array}{c} 1 \\ 3 \end{array} \right)$
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 $\left(\begin{array}{c} 3 \\ 5 \end{array} \right)$
 $\left(\begin{array}{c} 1 \\ 6 \end{array} \right)$
 $\left(\begin{array}{c} 2 \\ 2 \end{array} \right)$
 $\left(\begin{array}{c} 1 \\ 1 \end{array} \right)$
 $\left(\begin{array}{c} 0 \\ 7 \end{array} \right)$
 $\left(\begin{array}{c} 2 \\ 8 \end{array} \right)$
 $\left(\begin{array}{c} 2 \\ 9 \end{array} \right)$
 $\left(\begin{array}{c} 0 \\ 10 \end{array} \right)$
 $\left(\begin{array}{c} 1 \\ 1 \end{array} \right)$
 $\left(\begin{array}{c} 0 \\ 11 \end{array} \right)$
 $\left(\begin{array}{c} -1 \\ -1 \end{array} \right)$
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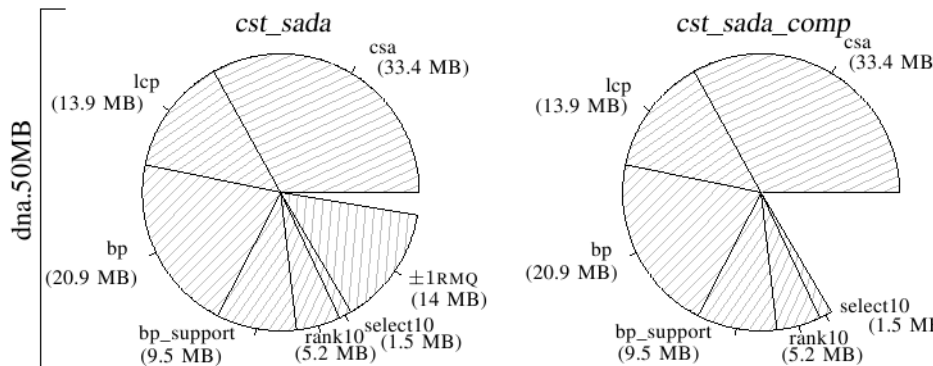
RMQ: Peak memory consumption at construction



RMQ: Final memory consumption and query time



Compressed Suffix Trees: Memory



Other data structures for RMQs

Non-systematic solutions

- sada: BPS of the extended Cartesian Tree ($4n + o(n)$ bits) by Sadakane (JDA 2007)
- 2dmin: BPS of the 2d-Min-Heap ($2n + o(n)$) by Fischer (2009)

Systematic solutions

- succ: Succinct solution ($7n$ bits + size of input array) by Fischer ()
- compr: Compressed solution ($\approx 3n$ bits + size of input array) by Fischer et al. (DCC 2008)

Experimental results

Any Questions?