#### Evidence for Long-Tails in SLS Algorithms

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#### Motivation and Aim of This Talk

- SLS = successful paradigm for solving SAT
- Recent development [LW20]:
  - GAPSAT solves not the original instance, but a modified, yet logically equivalent one
  - Empirically shown: on average, this improves the performance of state-of-the-art SLS solvers

*Caveat:* Only shallow understanding of how runtimes are affected

**AIM:** Model this modification process and conduct an empirical analysis of the hardness of logically equivalent formulas

- A lognormal distribution perfectly characterizes the hardness
  The hardness is long-tailed
- Restarts are useful for long-tailed algorithms

#### Statistical Hardness Distribution with ALFA

• Model the addition of a set of logically equivalent clauses *L* to a formula *F* and the subsequent solving of the amended formula  $F^{(1)} := F \cup L$  by an SLS solver

Algorithm: Adjusted logical formula algorithm (ALFA)

**Input:** Boolean formula F, **Promise:**  $F \in SAT$ 

Generate **randomly** a set *L* of clauses such that  $F \models L$ Call SLS( $F \cup L$ ) for some SLS solver SLS

Use width-4-bounded resolution to generate L



 $\begin{array}{c|c} \text{ solve 100 times with each solver } S \\ F^{(1)} & \text{flips}_{S}(F^{(1)}, s_{1}), & \dots, & \text{flips}_{S}(F^{(1)}, s_{100}) \\ \vdots \\ F^{(5000)} \end{array}$ 

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$$F = \begin{cases} solve 100 \text{ times with each solver } S \\ F^{(1)} & flips_{S}(F^{(1)}, s_{1}), & \dots, & flips_{S}(F^{(1)}, s_{100}) \\ flips_{S}(F^{(2)}, s_{1}), & \dots, & flips_{S}(F^{(2)}, s_{100}) \\ \vdots \\ F^{(5000)} & flips_{S}(F^{(5000)}, s_{1}), & \dots, & flips_{S}(F^{(5000)}, s_{100}) \\ flips_{S}(F^{(5000)}, s_{10}) \\ flips_{S}(F^{(5000)}, s_{10})$$

Given such a sample  $(x_1, \ldots, x_{5000})$ , we will plot the empirical distribution function

$$\hat{F}_{5000}(t) := \frac{1}{5000} \sum_{i=1}^{5000} \mathbb{1}_{\{x_i \leq t\}}, \quad t \in \mathbb{R}.$$

Instance Types:

- Hidden Solution (with different parameters)
- Oniform Random
- Factoring
- Coloring

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Used Solvers:

- Schöning's Random Walk Algorithm SRWA [Sch02] all instance types
- **PROBSAT solver family**<sup>1</sup> [BS12] -55 hidden solution instances with  $n \in \{50, 100, 150, 200, 300, 800\}$
- **Solution** YalSAT<sup>2</sup> [Bie14] -10 hidden solution instances with n = 300

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Total CPU time: 80 years

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## Experimental Results and Statistical Evaluation (1/3)

Instance A (Factoring solved with SRWA)



Instance B (Hidden Solution solved with PROBSAT)



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#### Experimental Results and Statistical Evaluation (2/3)

Résumé: LogN seems to be well suited

 $\rightarrow$  Concretize this through  $\chi^2$  goodness-of-fit test (and additional bootstrap test)



#### Experimental Results and Statistical Evaluation (3/3)

Table 1: Statistical goodness-of-fit results for ALFA+SRWA									
	hidden different chances uniform factoring coloring    total								
rejected	0	2	1	2	0	5			
# of instances	20	120	25	33	32	230			

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Table 2: Goodness-of-fit results for ALFA+PROBSAT

number of variables	50	100	150	200	300	800	total
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Table 3: Goodness-of-fit results for ALFA+YALSAT

	$\chi^2$	Ĵ
rejected	2	0
# of instances	10	10

#### Conjecture (Strong Conjecture)

The runtime of ALFA with SLS  $\in$  {SRWA, PROBSAT, YALSAT} follows a lognormal distribution.

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#### Definition ([FKZ11])

A positive, real-valued random variable X is long-tailed, if and only if

$$\forall x \in \mathbb{R}^+ : \Pr\left[X > x\right] > 0 \quad \text{and} \quad \forall y \in \mathbb{R}^+ : \lim_{x \to \infty} \frac{\Pr\left[X > x + y\right]}{\Pr\left[X > x\right]} = 1.$$

#### Conjecture (Weak Conjecture)

The runtime of ALFA with SLS  $\in$  {SRWA, PROBSAT, YALSAT} follows a long-tailed distribution.

# Restarts Are Useful For Long-Tailed Distributions

#### Definition

Restarts are *useful* if there is a t > 0 such that

 $\mathbf{E}[X_t] < \mathbf{E}[X].$ 

A condition for the usefulness of restart was proven in [Lor18]:

Theorem ([Lor18])

Strong Conjecture (runtimes are lognormally distributed)  $\implies$  Restarts are useful

Our paper extends this result and mathematically proves that restarts are useful even if only the Weak Conjecture holds (i. e., restarts are useful for long-tailed distributions).

#### Theorem

Consider a positive, long-tailed random variable *X* with continuous pdf *f*, cdf *F*, and hazard rate function  $r(t) := \frac{f(t)}{1 - F(t)}$ .

Also assume that

- either  $E[X] = \infty$  holds;
- or the limits  $\lim_{t\to\infty} r(t)$  and  $\lim_{t\to\infty} t^2 \cdot f(t)$  both exist.

In both cases, restarts are useful for X.

The conditions of this theorem are not restrictive since all naturally occurring long-tail distributions satisfy these conditions (see [NWZ20]).

#### Conjecture (Corollary of the Weak Conjecture)

Restarts are useful for ALFA with SLS  $\in$  {SRWA, PROBSAT, YALSAT}.

# Summary and Outlook

Have provided compelling evidence that the runtime of Alfa follows a long-tailed or lognormal distribution.

According to [SMP11]: super-linear speedups by parallelization  $\implies$  usefulness of restarts

Question 1: Can super-linear speedups be obtained by parallelizing Alfa-type algorithms?

Question 2: Can some of the Conjectures be theoretically proven?

Question 3: What can be said about CDCL solvers?

