

Dr. Lucia Draque Penso Dr. Jens Maßberg Institut für Optimierung und Operations Research Sommersemester 2014

## Mathematics of Games

**Exercise Session 6** 

Exercise Session due on Thursday, 12.06.2014, by 14:15pm, N24-H15. Total : 20 Points Hand-in IN PAIRS! Let the games begin! :-) Enjoy World Cup 2014! :-)

- 1. Cat-Dog Fight Game: A cat and a dog are fighting for a prize whose current value at any time t = 0, 1, ... is v > 1. Fighting costs 1 unit per period. If only one animal stops fighting at period t, his opponent wins alone the prize without incurring a fighting cost that period, and the choice of the second stopping time is irrelevant. If both animals stop fighting simultaneously, neither one wins the prize. That is, if we consider a per-period discount factor  $\delta$ , the (symmetric) payoff functions are:
  - $L(t) = -(1 + \delta + \ldots + \delta^{t-1}) = -\frac{1-\delta^t}{1-\delta}$ , for the loser(s), and
  - $W(t) = L(t) + \delta^t v$ , for the winner (in case there is one).

Give a symmetric subgame-perfect Nash Equilibrium for the *Cat-Dog Fight Game*.

*Hint:* You may consider a mixed strategy profile.

[5 Points]

2. Three players bargain over the partition of a pie of size 1. A partition is a triple  $(x_1, x_2, x_3)$  of shares for each player, where  $x_i \ge 0$ ,  $\sum_{i=1}^3 x_i = 1$ . At dates  $3k+1, k = 0, 1, \ldots$ , player 1 offers a division, then players 2 and 3 simultaneously decide whether they accept or veto. If players 2 and 3 both accept, the game is over. If one or both of them veto, bargaining goes on. Similarly, at dates 3k+2 (respectively, 3k), player 2 (respectively, player 3) makes the offer. The game stops once an offer by one player has been accepted by the other two players. The players have common discount factor  $\delta$ . Show that, for every  $\delta \in (0, 1)$ , all partitions can be supported as a subgame-perfect Nash Equilibrium.

[ 5 Points]

3. Consider a Cournot duopoly operating in a market with inverse demand P(Q) = a - Q, where  $Q = q_1 + q_2$  is the aggregate quantity on the market. Both firms have total costs  $c_i(q_i) = cq_i$  with a constant c, but demand is uncertain: it is high  $(a = a_H)$  with probability  $\gamma$  and low  $(a = a_L)$  with probability  $1 - \gamma$ . So the payoff depends on a and is  $u_i(q_i, q_j, a) = (P(Q) - c)q_i$  for both firms. Furthermore, information is asymmetric: firm 1 knows whether demand is high or low, but firm 2 does not. All of this is common knowledge. The two firms simultaneously choose quantities. What is the pure-strategy Bayesian Nash equilibrium of this static Bayesian game?

[5 Points]

4. Consider the following model of Bertrand duopoly with differentiated products. Demand for firm *i* is  $q_i(p_i, p_j) = a - p_i - b_i \cdot p_j$ . Costs are zero for both firms. The sensitivity of firm *i*'s demand to firm *j*'s price is either high or low. That is,  $b_i$  is either  $b_H$  or  $b_L$ , where  $b_H > b_L > 0$ . For each firm,  $b_i = b_H$  with probability  $\gamma$  and  $b_i = b_L$  with probability  $1 - \gamma$ , independent of the realization of  $b_j$ . Each firm knows its own  $b_i$  but not its competitor's. All of this is common knowledge. Which (four) conditions define a pure-strategy Bayesian Nash Equilibrium of this game? What is the pure-strategy Bayesian Nash Equilibrium in the specific case  $\gamma = \frac{1}{2}$ ? What about if  $\gamma = 1$  or if  $\gamma = 0$ ?

[5 Points]